Z8300 Low Power Z80®L CPU Central Processing Unit

Zilog

Product Specification

April 1985

Features

- The Z80L combines the high pertormance of the Z80 CPU with extremely low power consumption. It has the identical pinout and instruction set of the Z80. The result is increased reliability and lower system power requirements. This dramatic power savings makes the Z80L a natural choice for both hand-held and battery backup applications.
- The Z80L CPU is offered in two versions: Z8300-1—1.0 MHz clock, 15 mA typical current consumption Z8300-3—2.5 MHz clock, 25 mA typical current consumption
- The extensive instruction set contains 158 instructions, resulting in sophisticated data handling capabilities. The 78 instructions of the 8080A are included as a subset; 8080A and Z80 Family software compatibility is maintained.
- The Z80L microprocessors and associated family of peripheral controllers are linked by a vectored interrupt system. This system can be daisy-chained to allow implementation of a priority interrupt scheme. Little, if any, additional logic is required for daisy-chaining.
- Duplicate sets of both general-purpose and flag registers are provided, easing the design and operation of system software. Two 16-bit index registers facilitate program processing of tables and arrays.
- There are three modes of high-speed interrupt processing: 8080 similar, non-Z80 peripheral device, and Z80 Family peripheral with or without daisy chain.
- On-chip dynamic memory refresh counter.

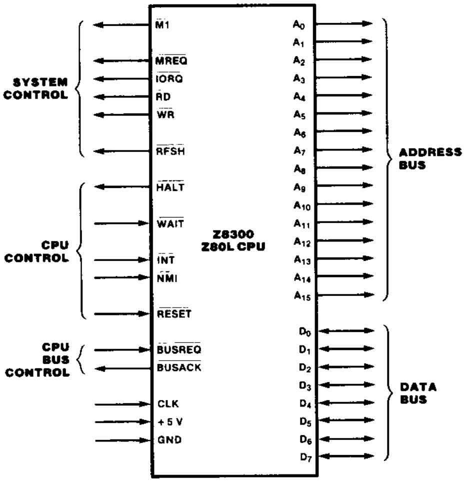


Figure 1. Pin Functions

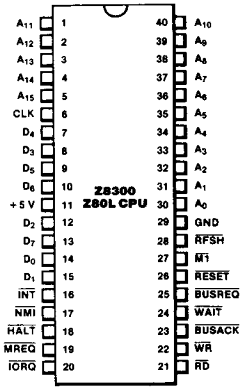


Figure 2. 40-pin Dual-In-Line Package (DIP) Pin Assignments

General Description

The Z80L CPUs are fourth-generation microprocessors with exceptional computational power. They offer high system throughput and efficient memory utilization combined with extremely low power consumption. The internal registers contain 208 bits of read/write memory that are accessible to the programmer. These registers include two sets of six general-purpose registers which may be used individually as either 8-bit registers or as 16-bit register pairs. In addition, there are two sets of accumulator and flag registers. A group of "Exchange" instructions makes either set of main or alternate registers accessible to the programmer. The alternate set allows operation in foreground-background mode or it may be reserved for very fast interrupt response.

The Z80L also contains a Stack Pointer, Program Counter, two index registers, a Refresh register (counter), and an Interrupt register. The CPU is easy to incorporate into a system since it requires only a single +5 V power source, all output signals are fully decoded and timed to control standard memory or peripheral circuits, and it is supported by an extensive family of peripheral controllers. The internal block diagram (Figure 3) shows the primary functions of the Z80L processors. Subsequent text provides more detail on the Z80L I/O controller family, registers, instruction set, interrupts and daisy chaining, CPU timing, and low power requirements.

Z80L Low Power Feature. The Z80L Family offers state-of-the-art microprocessor performance with extremely low power consumption. Its low power requirement rivals comparable CMOS microprocessors. The Z80L Family's lower power consumption provides the ability to reduce system power requirements and enables its use in applications not previously possible. The Z80L is very well suited to battery backup applications or to systems operating primarily on batteries in hand-held or portable systems.

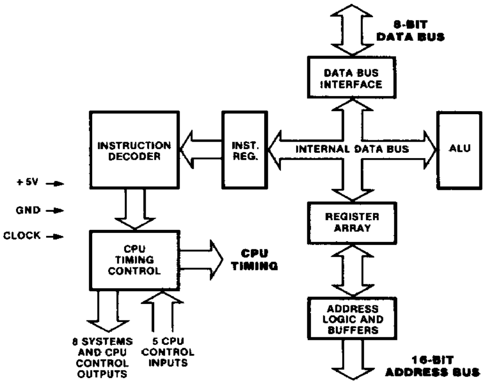


Figure 3. Z80L CPU Block Diagram

Z80L Microprocessor Family

The Zilog Z80L microprocessor is the central element of a comprehensive microprocessor product family. This family works together in most applications with minimum requirements for additional logic, facilitating the design of efficient and cost-effective microcomputer-based systems.

The Z80 Family components provide extensive support for the Z80L microprocessor. These are:

- The PIO (Parallel Input/Output) operates in both data-byte I/O transfer mode (with handshaking) and in bit mode (without handshaking). The PIO may be configured to interface with standard parallel peripheral devices such as printers, tape punches, and keyboards.
- The CTC (Counter/Timer Circuit) features four programmable 8-bit counter/timers, each of which has an 8-bit prescaler. Each of the four channels may be configured to operate in either counter or timer mode.
- The SIO (Serial Input/Output) controller offers two channels. It is capable of operating in a variety of programmable modes for both synchronous and asynchronous communication, including Bi-Synch and SDLC.

Z80L CPU Registers

Figure 4 shows three groups of registers within the Z80L CPU. The first group consists of duplicate sets of 8-bit registers: a principal set and an alternate set (designated by ' [prime], e.g., A'). Both sets consist of the Accumulator Register, the Flag Register, and six general-purpose registers. Transfer of data between these duplicate sets of registers is accomplished by use of "Exchange" instructions. The result is faster response to interrupts and easy, efficient implementation of such versatile programming techniques as background-foreground data processing. The second set of registers consists of six registers with assigned functions. These are the I (Interrupt Register), the R (Refresh Register), the IX and IY (Index Registers), the SP (Stack Pointer), and the PC (Program Counter). The third group consists of two interrupt status flip-flops, plus an additional pair of flip-flops which assists in identifying the interrupt mode at any particular time. Table 1 provides further information on these registers.

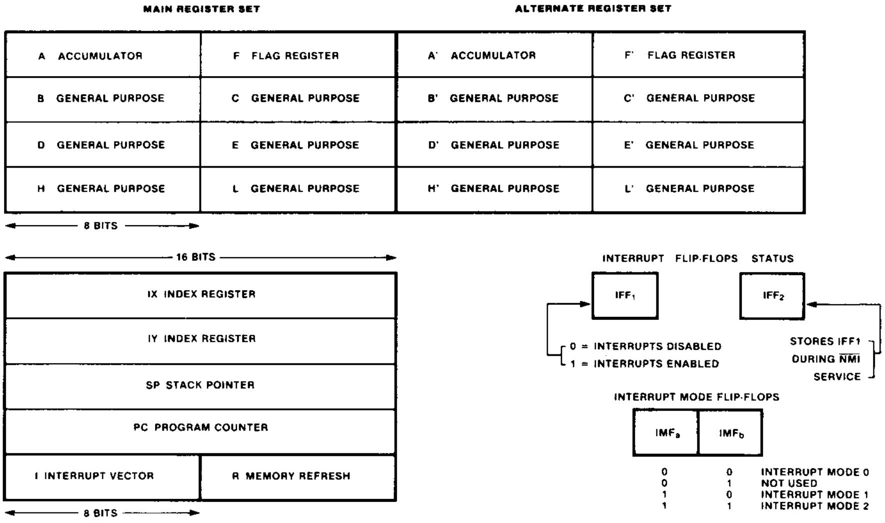


Figure 4. CPU Registers

Z80L CPU Registers (Continued)

Register		Size (Bits)	Remarks
A, A′	Accumulator	8	Stores an operand or the results of an operation.
F, F′	Flags	8	See Instruction Set.
B, B′	General Purpose	8	Can be used separately or as a 16-bit register with C.
C, C′	General Purpose	8	See B, above.
D, D′	General Purpose	8	Can be used separately or as a 16-bit register with E.
E, E′	General Purpose	8	See D, above.
H, H′	General Purpose	8	Can be used separately or as a 16-bit register with L.
L, L′	General Purpose	8	See H, above.
			Note: The (B,C), (D,E), and (H,L) sets are combined as follows: B — High byte C — Low byte D — High byte E — Low byte H — High byte L — Low byte
I	Interrupt Register	8	Stores upper eight bits of memory address for vectored interrupt processing.
R	Refresh Register	8	Provides user-transparent dynamic memory refresh. Lower seven bits are automatically incremented and all eight are placed on the address bus during each instruction fetch cycle refresh time.
IX	Index Register	16	Used for indexed addressing.
IY	Index Register	16	Same as IX, above.
SP	Stack Pointer	16	Holds address of the top of the stack. See Push or Pop in instruction set.
PC	Program Counter	16	Holds address of next instruction.
IFF_1-IFF_2	Interrupt Enable	Flip-Flops	Set or reset to indicate interrupt status (see Figure 4).
IMFa-IMFb	Interrupt Mode	Flip-Flops	Reflect Interrupt mode (see Figure 4).

Table 1. Z80L CPU Registers

Interrupts: General Operation

The CPU accepts two interrupt input signals: $\overline{\text{NMI}}$ and $\overline{\text{INT}}$. The NMI is a non-maskable interrupt and has the highest priority. $\overline{\text{INT}}$ is a lower priority interrupt and it requires that interrupts be enabled in software in order to operate. $\overline{\text{INT}}$ can be connected to multiple peripheral devices in a wired-OR configuration.

The Z80L has a single response mode for interrupt service for the non-maskable interrupt. The maskable interrupt, $\overline{\text{INT}}$, has three programmable response modes available. These are:

- Mode 0 — similar to the 8080 microprocessor.
- Mode 1 — Peripheral Interrupt service, for use with non-8080/Z80 systems.
- Mode 2 — a vectored interrupt scheme, usually daisy-chained, for use with Z80 Family and compatible peripheral devices.

The CPU services interrupts by sampling the $\overline{\text{NMI}}$ and $\overline{\text{INT}}$ signals at the rising edge of the last clock of an instruction. Further interrupt service processing depends upon the type of interrupt that was detected. Details on interrupt responses are shown in the CPU Timing Section.

Interrupts: General Operation (Continued)

Non-Maskable Interrupt ($\overline{\text{NMI}}$). The nonmaskable interrupt cannot be disabled by program control and therefore will be accepted all times by the CPU. $\overline{\text{NMI}}$ is usually reserved for servicing only the highest priority type interrupts, such as that for orderly shutdown after power failure has been detected. After recognition of the $\overline{\text{NMI}}$ signal (providing $\overline{\text{BUSREQ}}$ is not active), the CPU jumps to restart location 0066H. Normally, software starting at this address contains the interrupt service routine.

Maskable Interrupt ($\overline{\text{INT}}$). Regardless of the interrupt mode set by the user, the Z80L response to a maskable interrupt input follows a common timing cycle. After the interrupt has been detected by the CPU (provided that interrupts are enabled and $\overline{\text{BUSREQ}}$ is not active) a special interrupt processing cycle begins. This is a special fetch ($\overline{\text{M1}}$) cycle in which $\overline{\text{IORQ}}$ becomes active rather than $\overline{\text{MREQ}}$, as in a normal $\overline{\text{M1}}$ cycle. In addition, this special $\overline{\text{M1}}$ cycle is automatically extended by two $\overline{\text{WAIT}}$ states, to allow for the time required to acknowledge the interrupt request and to place the interrupt vector on the bus.

Mode 0 Interrupt Operation. This mode is similar to the 8080 microprocessor interrupt service procedures. The interrupting device places an instruction on the data bus. This is normally a Restart instruction, which will initiate a call to the selected one of eight restart locations in page zero of memory. Unlike the 8080, the Z80 CPU responds to the Call instruction with only one interrupt acknowledge cycle followed by two memory read cycles.

Mode 1 Interrupt Operation. Mode 1 operation is very similar to that for the $\overline{\text{NMI}}$. The principal difference is that the Mode 1 interrupt has a vector address of 0038H only.

Mode 2 Interrupt Operation. This interrupt mode has been designed to utilize most effectively the capabilities of the Z80L microprocessor and its associated peripheral family. The interrupting peripheral device selects the starting address of the interrupt service routine. It does this by placing an 8-bit vector on the data bus during the interrupt acknowledge cycle. The CPU forms a pointer using this byte as the lower 8 bits and the contents of the I register as the upper 8 bits. This points to an entry in a table of addresses for interrupt service routines. The CPU then calls the routine at that address. This flexibility in selecting the interrupt service routine address allows the peripheral device to use several different types of service routines. These routines may be located at any available location in memory. Since the interrupting device supplies the low-order byte of the 2-byte vector, bit 0 (A_0) must be a zero.

Interrupt Priority (Daisy Chaining and Nested Interrupts). The interrupt priority of each peripheral device is determined by its physical location within a daisy-chain configuration. Each device in the chain has an interrupt enable input line (IEI) and an interrupt enable output line (IEO), which is fed to the next lower priority device. The first device in the daisy chain has its IEI input hardwired to a High level. The first device has highest priority, while each succeeding device has a corresponding lower priority. This arrangement permits the CPU to select the highest priority interrupt from several simultaneously interrupting peripherals.

The interrupting device disables its IEO line to the next lower priority peripheral until it has been serviced. After servicing, its IEO line is raised, allowing lower priority peripherals to demand interrupt servicing.

The Z80L CPU will nest (queue) any pending interrupts or interrupts received while a selected peripheral is being serviced.

Interrupt Enable/Disable Operation. Two flip-flops, IFF_1 and IFF_2, referred to in the register description are used to signal the CPU interrupt status. Operation of the two flip-flops is described in Table 2. For more details, refer to the *Z80 CPU Technical Manual* and *Z80 Assembly Language Manual.*

Action	IFF_1	IFF_2	Comments
CPU Reset	0	0	Maskable interrupt $\overline{\text{INT}}$ disabled
DI instruction execution	0	0	Maskable interrupt $\overline{\text{INT}}$ disabled
EI instruction execution	1	1	Maskable interrupt $\overline{\text{INT}}$ enabled
LD A,I instruction execution	•	•	IFF_2 → Parity flag
LD A,R instruction execution	•	•	IFF_2 → Parity flag
Accept $\overline{\text{NMI}}$	0	IFF_1	IFF_1 → IFF_2 (Maskable interrupt $\overline{\text{INT}}$ disabled)
RETN instruction execution	IFF_2	•	IFF_2 → IFF_1 at completion of an $\overline{\text{NMI}}$ service routine.

Table 2. State of Flip-Flops

Instruction Set

The Z80L microprocessor has one of the most powerful and versatile instruction sets available in any 8-bit microprocessor and identical to that of the Z80. It includes such unique operations as a block move for fast, efficient data transfers within memory or between memory and I/O. It also allows operations on any bit in any location in memory.

The following is a summary of the Z80L instruction set and shows the assembly language mnemonic, the operation, the flag status, and gives comments on each instruction. The *Z80 CPU Technical Manual* (03-0029-XX) and *Assembly Language Programming Manual* (03-0002-XX) contain significantly more details for programming use.

The instructions in Table 2 are divided into the following categories:

- □ 8-bit loads
- □ 16-bit loads
- □ Exchanges, block transfers, and searches
- □ 8-bit arithmetic and logic operations
- □ General-purpose arithmetic and CPU control
- □ 16-bit arithmetic operations
- □ Rotates and shifts
- □ Bit set, reset, and test operations
- □ Jumps
- □ Calls, returns, and restarts
- □ Input and output operations

A variety of addressing modes are implemented to permit efficient and fast data transfer between various registers, memory locations, and input/output devices. These addressing modes include:

- □ Immediate
- □ Immediate extended
- □ Modified page zero
- □ Relative
- □ Extended
- □ Indexed
- □ Register
- □ Register indirect
- □ Implied
- □ Bit

8-Bit Load Group

Mnemonic	Symbolic Operation	S	Z		Flags H		P/V	N	C	Opcode 76 543 210	Hex	No.of Bytes	No.of M Cycles	No.of T States	Comments
LD r, r'	r ← r'	•	•	X	•	X	•	•	•	01 r r'		1	1	4	r, r' Reg.
LD r, n	r ← n	•	•	X	•	X	•	•	•	00 r 110 ← n →		2	2	7	000 B 001 C
LD r, (HL)	r ← (HL)	•	•	X	•	X	•	•	•	01 r 110		1	2	7	010 D
LD r, (IX+d)	r ← (IX+d)	•	•	X	•	X	•	•	•	11 011 101 01 r 110 ← d →	DD	3	5	19	011 E 100 H 101 L
LD r, (IY+d)	r ← (IY+d)	•	•	X	•	X	•	•	•	11 111 101 01 r 110 ← d →	FD	3	5	19	111 A
LD (HL), r	(HL) ← r	•	•	X	•	X	•	•	•	01 110 r		1	2	7	
LD (IX+d), r	(IX+d) ← r	•	•	X	•	X	•	•	•	11 011 101 01 110 r ← d →	DD	3	5	19	
LD (IY+d), r	(IY+d) ← r	•	•	X	•	X	•	•	•	11 111 101 01 110 r ← d →	FD	3	5	19	
LD (HL), n	(HL) ← n	•	•	X	•	X	•	•	•	00 110 110 ← n →	36	2	3	10	
LD (IX+d), n	(IX+d) ← n	•	•	X	•	X	•	•	•	11 011 101 00 110 110 ← d → ← n →	DD 36	4	5	19	
LD (IY+d), n	(IY+d) ← n	•	•	X	•	X	•	•	•	11 111 101 00 110 110 ← d → ← n →	FD 36	4	5	19	
LD A, (BC)	A ← (BC)	•	•	X	•	X	•	•	•	00 001 010	0A	1	2	7	
LD A, (DE)	A ← (DE)	•	•	X	•	X	•	•	•	00 011 010	1A	1	2	7	
LD A, (nn)	A ← (nn)	•	•	X	•	X	•	•	•	00 111 010 ← n → ← n →	3A	3	4	13	
LD (BC), A	(BC) ← A	•	•	X	•	X	•	•	•	00 000 010	02	1	2	7	
LD (DE), A	(DE) ← A	•	•	X	•	X	•	•	•	00 010 010	12	1	2	7	
LD (nn), A	(nn) ← A	•	•	X	•	X	•	•	•	00 110 010 ← n → ← n →	32	3	4	13	
LD A, I	A ← I	↕	↕	X	0	X	IFF	0	•	11 101 101 01 010 111	ED 57	2	2	9	
LD A, R	A ← R	↕	↕	X	0	X	IFF	0	•	11 101 101 01 011 111	ED 5F	2	2	9	
LD I, A	I ← A	•	•	X	•	X	•	•	•	11 101 101 01 000 111	ED 47	2	2	9	
LD R, A	R ← A	•	•	X	•	X	•	•	•	11 101 101 01 001 111	ED 4F	2	2	9	

NOTES: r, r' means any of the registers A, B, C, D, E, H, L.
IFF the content of the interrupt enable flip-flop, (IFF) is copied into the P/V flag.
For an explanation of flag notation and symbols for mnemonic tables, see Symbolic Notation section following tables.

16-Bit Load Group

Mnemonic	Symbolic Operation				Flags					Opcode		No.of	No.of M	No.of T	Comments
		S	Z		H		P/V	N	C	76 543 210	Hex	Bytes	Cycles	States	
LD dd, nn	dd ← nn	•	•	X	•	X	•	•	•	00 dd0 001 ← n → ← n →		3	3	10	dd Pair 00 BC 01 DE 10 HL 11 SP
LD IX, nn	IX ← nn	•	•	X	•	X	•	•	•	11 011 101 00 100 001 ← n → ← n →	DD 21	4	4	14	
LD IY, nn	IY ← nn	•	•	X	•	X	•	•	•	11 111 101 00 100 001 ← n → ← n →	FD 21	4	4	14	
LD HL, (nn)	H ← (nn + 1) L ← (nn)	•	•	X	•	X	•	•	•	00 101 010 ← n → ← n →	2A	3	5	16	
LD dd, (nn)	dd_H ← (nn + 1) dd_L ← (nn)	•	•	X	•	X	•	•	•	11 101 101 01 dd1 011 ← n → ← n →	ED	4	6	20	
LD IX, (nn)	IX_H ← (nn + 1) IX_L ← (nn)	•	•	X	•	X	•	•	•	11 011 101 00 101 010 ← n → ← n →	DD 2A	4	6	20	
LD IY, (nn)	IY_H ← (nn + 1) IY_L ← (nn)	•	•	X	•	X	•	•	•	11 111 101 00 101 010 ← n → ← n →	FD 2A	4	6	20	
LD (nn), HL	(nn + 1) ← H (nn) ← L	•	•	X	•	X	•	•	•	00 100 010 ← n → ← n →	22	3	5	16	
LD (nn), dd	(nn + 1) ← dd_H (nn) ← dd_L	•	•	X	•	X	•	•	•	11 101 101 01 dd0 011 ← n → ← n →	ED	4	6	20	
LD (nn), IX	(nn + 1) ← IX_H (nn) ← IX_L	•	•	X	•	X	•	•	•	11 011 101 00 100 010 ← n → ← n →	DD 22	4	6	20	
LD (nn), IY	(nn + 1) ← IY_H (nn) ← IY_L	•	•	X	•	X	•	•	•	11 111 101 00 100 010 ← n → ← n →	FD 22	4	6	20	
LD SP, HL	SP ← HL	•	•	X	•	X	•	•	•	11 111 001	F9	1	1	6	
LD SP, IX	SP ← IX	•	•	X	•	X	•	•	•	11 011 101 11 111 001	DD F9	2	2	10	
LD SP, IY	SP ← IY	•	•	X	•	X	•	•	•	11 111 101 11 111 001	FD F9	2	2	10	
PUSH qq	(SP − 2) ← qq_L (SP − 1) ← qq_H SP → SP − 2	•	•	X	•	X	•	•	•	11 qq0 101		1	3	11	qq Pair 00 BC 01 DE 10 HL 11 AF
PUSH IX	(SP − 2) ← IX_L (SP − 1) ← IX_H SP → SP − 2	•	•	X	•	X	•	•	•	11 011 101 11 100 101	DD E5	2	4	15	
PUSH IY	(SP − 2) ← IY_L (SP − 1) ← IY_H SP → SP − 2	•	•	X	•	X	•	•	•	11 111 101 11 100 101	FD E5	2	4	15	
POP qq	qq_H ← (SP + 1) qq_L ← (SP) SP → SP + 2	•	•	X	•	X	•	•	•	11 qq0 001		1	3	10	
POP IX	IX_H ← (SP + 1) IX_L ← (SP) SP → SP + 2	•	•	X	•	X	•	•	•	11 011 101 11 100 001	DD E1	2	4	14	
POP IY	IY_H ← (SP + 1) IY_L ← (SP) SP → SP + 2	•	•	X	•	X	•	•	•	11 111 101 11 100 001	FD E1	2	4	14	

NOTES: dd is any of the register pairs BC, DE, HL, SP.
qq is any of the register pairs AF, BC, DE, HL.
$(PAIR)_H$, $(PAIR)_L$ refer to high order and low order eight bits of the register pair respectively,
e.g., BC_L = C, AF_H = A.

Exchange, Block Transfer, Block Search Groups

EX DE, HL	DE ↔ HL	•	•	X	•	X	•	•	•	11 101 011	EB	1	1	4	
EX AF, AF'	AF ↔ AF'	•	•	X	•	X	•	•	•	00 001 000	08	1	1	4	
EXX	BC ↔ BC' DE ↔ DE' HL ↔ HL'	•	•	X	•	X	•	•	•	11 011 001	D9	1	1	4	Register bank and auxiliary register bank exchange
EX (SP), HL	H ↔ (SP + 1) L ↔ (SP)	•	•	X	•	X	•	•	•	11 100 011	E3	1	5	19	
EX (SP), IX	IX_H ↔ (SP + 1) IX_L ↔ (SP)	•	•	X	•	X	•	•	•	11 011 101 11 100 011	DD E3	2	6	23	
EX (SP), IY	IY_H ↔ (SP + 1) IY_L ↔ (SP)	•	•	X	•	X	•	•	•	11 111 101 11 100 011	FD E3	2	6	23	
LDI	(DE) ← (HL) DE ← DE + 1 HL ← HL + 1 BC ← BC − 1	•	•	X	0	X	① ↕	0	•	11 101 101 10 100 000	ED A0	2	4	16	Load (HL) into (DE), increment the pointers and decrement the byte counter (BC)
LDIR	(DE) ← (HL) DE ← DE + 1 HL ← HL + 1 BC ← BC − 1 Repeat until BC = 0	•	•	X	0	X	① 0	0	•	11 101 101 10 110 000	ED B0	2 2	5 4	21 16	If BC ≠ 0 If BC = 0

NOTE: ① P/V flag is 0 if the result of BC − 1 = 0, otherwise P/V = 1.

Exchange, Block Transfer, Block Search Groups (Continued)

Mnemonic	Symbolic Operation	S	Z		Flags H		P/V	N	C	Opcode 76 543 210	Hex	No.of Bytes	No.of M Cycles	No.of T States	Comments
LDD	(DE) ← (HL) DE ← DE − 1 HL ← HL − 1 BC ← BC − 1	•	•	X	0	X	① ↕	0	•	11 101 101 10 101 000	ED A8	2	4	16	
LDDR	(DE) ← (HL) DE ← DE − 1 HL ← HL − 1 BC ← BC − 1 Repeat until BC = 0	•	•	X	0	X	② 0	0	•	11 101 101 10 111 000	ED B8	2 2	5 4	21 16	If BC ≠ 0 If BC = 0
CPI	A − (HL) HL ← HL + 1 BC ← BC − 1	↕	③ ↕	X	↕	X	① ↕	1	•	11 101 101 10 100 001	ED A1	2	4	16	
CPIR	A − (HL) HL ← HL + 1 BC ← BC − 1 Repeat until A = (HL) or BC = 0	↕	③ ↕	X	↕	X	① ↕	1	•	11 101 101 10 110 001	ED B1	2 2	5 4	21 16	If BC ≠ 0 and A ≠ (HL) If BC = 0 or A = (HL)
CPD	A − (HL) HL ← HL − 1 BC ← BC − 1	↕	③ ↕	X	↕	X	① ↕	1	•	11 101 101 10 101 001	ED A9	2	4	16	
CPDR	A − (HL) HL ← HL − 1 BC ← BC − 1 Repeat until A = (HL) or BC = 0	↕	③ ↕	X	↕	X	① ↕	1	•	11 101 101 10 111 001	ED B9	2 2	5 4	21 16	If BC ≠ 0 and A ≠ (HL) If BC = 0 or A = (HL)

NOTES: ① P/V flag is 0 if the result of BC − 1 = 0, otherwise P/V = 1.
② P/V flag is 0 at completion of instruction only.
③ Z flag is 1 if A = (HL), otherwise Z = 0.

8-Bit Arithmetic and Logical Group

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	Opcode 76 543 210	Hex	No.of Bytes	No.of M Cycles	No.of T States	Comments
ADD A, r	A ← A + r	↕	↕	X	↕	X	V	0	↕	10 $\boxed{000}$ r		1	1	4	r Reg.
ADD A, n	A ← A + n	↕	↕	X	↕	X	V	0	↕	11 $\boxed{000}$ 110 ← n →		2	2	7	000 B 001 C 010 D
ADD A, (HL)	A ← A + (HL)	↕	↕	X	↕	X	V	0	↕	10 $\boxed{000}$ 110		1	2	7	011 E
ADD A, (IX + d)	A ← A + (IX + d)	↕	↕	X	↕	X	V	0	↕	11 011 101 10 $\boxed{000}$ 110 ← d →	DD	3	5	19	100 H 101 L 111 A
ADD A, (IY + d)	A ← A + (IY + d)	↕	↕	X	↕	X	V	0	↕	11 111 101 10 $\boxed{000}$ 110 ← d →	FD	3	5	19	
ADC A, s	A ← A + s + CY	↕	↕	X	↕	X	V	0	↕	$\boxed{001}$					s is any of r, n, (HL), (IX + d), (IY + d) as shown for ADD instruction. The indicated bits replace the $\boxed{000}$ in the ADD set above.
SUB s	A ← A − s	↕	↕	X	↕	X	V	1	↕	$\boxed{010}$					
SBC A, s	A ← A − s − CY	↕	↕	X	↕	X	V	1	↕	$\boxed{011}$					
AND s	A ← A ∧ s	↕	↕	X	1	X	P	0	0	$\boxed{100}$					
OR s	A ← A ∨ s	↕	↕	X	0	X	P	0	0	$\boxed{110}$					
XOR s	A ← A ⊕ s	↕	↕	X	0	X	P	0	0	$\boxed{101}$					
CP s	A − s	↕	↕	X	↕	X	V	1	↕	$\boxed{111}$					
INC r	r ← r + 1	↕	↕	X	↕	X	V	0	•	00 r $\boxed{100}$		1	1	4	
INC (HL)	(HL) ← (HL) + 1	↕	↕	X	↕	X	V	0	•	00 110 $\boxed{100}$		1	3	11	
INC (IX + d)	(IX + d) ← (IX + d) + 1	↕	↕	X	↕	X	V	0	•	11 011 101 00 110 $\boxed{100}$ ← d →	DD	3	6	23	
INC (IY + d)	(IY + d) ← (IY + d) + 1	↕	↕	X	↕	X	V	0	•	11 111 101 00 110 $\boxed{100}$ ← d →	FD	3	6	23	
DEC m	m ← m − 1	↕	↕	X	↕	X	V	1	•	$\boxed{101}$					m is any of r, (HL), (IX + d), (IY + d) as shown for INC. DEC same format and states as INC. Replace $\boxed{100}$ with $\boxed{101}$ in opcode.

General-Purpose Arithmetic and CPU Control Groups

Mnemonic	Symbolic Operation	S	Z		Flags H		P/V	N	C	Opcode 76 543 210	Hex	No.of Bytes	No.of M Cycles	No.of T States	Comments
DAA	Converts acc. content into packed BCD following add or subtract with packed BCD operands.	↕	↕	X	↕	X	P	•	↕	00 100 111	27	1	1	4	Decimal adjust accumulator.
CPL	A ← $\overline{A}$	•	•	X	1	X	•	1	•	00 101 111	2F	1	1	4	Complement accumulator (one's complement).
NEG	A ← 0 − A	↕	↕	X	↕	X	V	1	↕	11 101 101 01 000 100	ED 44	2	2	8	Negate acc. (two's complement).
CCF	CY ← $\overline{CY}$	•	•	X	X	X	•	0	↕	00 111 111	3F	1	1	4	Complement carry flag.
SCF	CY ← 1	•	•	X	0	X	•	0	1	00 110 111	37	1	1	4	Set carry flag.
NOP	No operation	•	•	X	•	X	•	•	•	00 000 000	00	1	1	4	
HALT	CPU halted	•	•	X	•	X	•	•	•	01 110 110	76	1	1	4	
DI ★	IFF ← 0	•	•	X	•	X	•	•	•	11 110 011	F3	1	1	4	
EI ★	IFF ← 1	•	•	X	•	X	•	•	•	11 111 011	FB	1	1	4	
IM 0	Set interrupt mode 0	•	•	X	•	X	•	•	•	11 101 101 01 000 110	ED 46	2	2	8	
IM 1	Set interrupt mode 1	•	•	X	•	X	•	•	•	11 101 101 01 010 110	ED 56	2	2	8	
IM 2	Set interrupt mode 2	•	•	X	•	X	•	•	•	11 101 101 01 011 110	ED 5E	2	2	8	

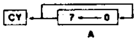
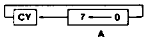
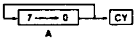
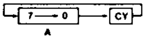
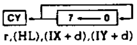
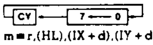
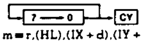
NOTES: IFF indicates the interrupt enable flip-flop.
CY indicates the carry flip-flop.
★ indicates interrupts are not sampled at the end of EI or DI.

16-Bit Arithmetic Group

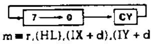
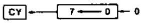
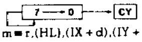
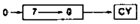
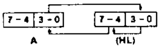
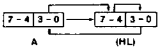
Mnemonic	Symbolic Operation	S	Z		Flags H		P/V	N	C	Opcode 76 543 210	Hex	No.of Bytes	No.of M Cycles	No.of T States	Comments
ADD HL, ss	HL ← HL + ss	•	•	X	X	X	•	0	↕	00 ss1 001		1	3	11	ss Reg. 00 BC 01 DE 10 HL 11 SP
ADC HL, ss	HL ← HL + ss + CY	↕	↕	X	X	X	V	0	↕	11 101 101 01 ss1 010	ED	2	4	15	
SBC HL, ss	HL ← HL − ss − CY	↕	↕	X	X	X	V	1	↕	11 101 101 01 ss0 010	ED	2	4	15	
ADD IX, pp	IX ← IX + pp	•	•	X	X	X	•	0	↕	11 011 101 01 pp1 001	DD	2	4	15	pp Reg. 00 BC 01 DE 10 IX 11 SP
ADD IY, rr	IY ← IY + rr	•	•	X	X	X	•	0	↕	11 111 101 00 rr1 001	FD	2	4	15	rr Reg. 00 BC 01 DE 10 IY 11 SP
INC ss	ss ← ss + 1	•	•	X	•	X	•	•	•	00 ss0 011		1	1	6	
INC IX	IX ← IX + 1	•	•	X	•	X	•	•	•	11 011 101 00 100 011	DD 23	2	2	10	
INC IY	IY ← IY + 1	•	•	X	•	X	•	•	•	11 111 101 00 100 011	FD 23	2	2	10	
DEC ss	ss ← ss − 1	•	•	X	•	X	•	•	•	00 ss1 011		1	1	6	
DEC IX	IX ← IX − 1	•	•	X	•	X	•	•	•	11 011 101 00 101 011	DD 2B	2	2	10	
DEC IY	IY ← IY − 1	•	•	X	•	X	•	•	•	11 111 101 00 101 011	FD 2B	2	2	10	

NOTES: ss is any of the register pairs BC, DE, HL, SP.
pp is any of the register pairs BC, DE, IX, SP.
rr is any of the register pairs BC, DE, IY, SP.

Rotate and Shift Group

Mnemonic	Symbolic Operation	S	Z		Flags H		P/V	N	C	Opcode 76 543 210	Hex	No.of Bytes	No.of M Cycles	No.of T States	Comments
RLCA		•	•	X	0	X	•	0	↕	00 000 111	07	1	1	4	Rotate left circular accumulator.
RLA		•	•	X	0	X	•	0	↕	00 010 111	17	1	1	4	Rotate left accumulator.
RRCA		•	•	X	0	X	•	0	↕	00 001 111	0F	1	1	4	Rotate right circular accumulator.
RRA		•	•	X	0	X	•	0	↕	00 011 111	1F	1	1	4	Rotate right accumulator.
RLC r		↕	↕	X	0	X	P	0	↕	11 001 011 00 $\boxed{000}$ r	CB	2	2	8	Rotate left circular register r.
RLC (HL)		↕	↕	X	0	X	P	0	↕	11 001 011 00 $\boxed{000}$ 110	CB	2	4	15	r Reg. 000 B 001 C 010 D 011 E 100 H 101 L 111 A
RLC (IX + d)	 r,(HL),(IX + d),(IY + d)	↕	↕	X	0	X	P	0	↕	11 011 101 11 001 011 ← d → 00 $\boxed{000}$ 110	DD CB	4	6	23	
RLC (IY + d)		↕	↕	X	0	X	P	0	↕	11 111 101 11 001 011 ← d → 00 $\boxed{000}$ 110	FD CB	4	6	23	
RL m	 m ≡ r,(HL),(IX + d),(IY + d)	↕	↕	X	0	X	P	0	↕	$\boxed{010}$					Instruction format and states are as shown for RLC's. To form new opcode replace $\boxed{000}$ or RLC's with shown code.
RRC m	 m ≡ r,(HL),(IX + d),(IY + d)	↕	↕	X	0	X	P	0	↕	$\boxed{001}$					

Rotate and Shift Group (Continued)

	Symbolic				Flags					Opcode		No.of	No.of M	No.of T	
Mnemonic	Operation	S	Z		H		P/V	N	C	76 543 210	Hex	Bytes	Cycles	States	Comments
RR m	 m ≡ r,(HL),(IX + d),(IY + d)	↕	↕	X	0	X	P	0	↕	$\boxed{011}$					
SLA m	 m ≡ r,(HL),(IX + d),(IY + d)	↕	↕	X	0	X	P	0	↕	$\boxed{100}$					
SRA m	 m ≡ r,(HL),(IX + d),(IY + d)	↕	↕	X	0	X	P	0	↕	$\boxed{101}$					
SRL m	 m ≡ r,(HL),(IX + d),(IY + d)	↕	↕	X	0	X	P	0	↕	$\boxed{111}$					
RLD		↕	↕	X	0	X	P	0	•	11 101 101 01 101 111	ED 6F	2	5	18	Rotate digit left and right between the accumulator and location (HL).
RRD		↕	↕	X	0	X	P	0	•	11 101 101 01 100 111	ED 67	2	5	18	The content of the upper half of the accumulator is unaffected.

Bit Set, Reset and Test Group

	Symbolic				Flags					Opcode		No.of	No.of M	No.of T	
Mnemonic	Operation	S	Z		H		P/V	N	C	76 543 210	Hex	Bytes	Cycles	States	Comments
BIT b, r	$Z \leftarrow \overline{r}_b$	X	↕	X	1	X	X	0	•	11 001 011 01 b r	CB	2	2	8	r Reg. 000 B 001 C 010 D 011 E 100 H 101 L 111 A
BIT b, (HL)	$Z \leftarrow (\overline{HL})_b$	X	↕	X	1	X	X	0	•	11 001 011 01 b 110	CB	2	3	12	
BIT b, $(IX+d)_b$	$Z \leftarrow (\overline{IX+d})_b$	X	↕	X	1	X	X	0	•	11 011 101 11 001 011 ← d → 01 b 110	DD CB	4	5	20	
BIT b, $(IY+d)_b$	$Z \leftarrow (\overline{IY+d})_b$	X	↕	X	1	X	X	0	•	11 111 101 11 001 011 ← d → 01 b 110	FD CB	4	5	20	b Bit Tested 000 0 001 1 010 2 011 3 100 4 101 5 110 6 111 7
SET b, r	$r_b \leftarrow 1$	•	•	X	•	X	•	•	•	11 001 011 $\boxed{11}$ b r	CB	2	2	8	
SET b, (HL)	$(HL)_b \leftarrow 1$	•	•	X	•	X	•	•	•	11 001 011 $\boxed{11}$ b 110	CB	2	4	15	
SET b, (IX+d)	$(IX+d)_b \leftarrow 1$	•	•	X	•	X	•	•	•	11 011 101 11 001 011 ← d → $\boxed{11}$ b 110	DD CB	4	6	23	
SET b, (IY+d)	$(IY+d)_b \leftarrow 1$	•	•	X	•	X	•	•	•	11 111 101 11 001 011 ← d → $\boxed{11}$ b 110	FD CB	4	6	23	
RES b, m	$m_b \leftarrow 0$ m ≡ r, (HL), (IX+d), (IY+d)	•	•	X	•	X	•	•	•	$\boxed{10}$					To form new opcode replace $\boxed{11}$ of SET b, s with $\boxed{10}$. Flags and time states for SET instruction.

NOTES: The notation m_b indicates bit b (0 to 7) or location m.

Jump Group

	Symbolic				Flags					Opcode		No.of	No.of M	No.of T	
Mnemonic	Operation	S	Z		H		P/V	N	C	76 543 210	Hex	Bytes	Cycles	States	Comments
JP nn	PC ← nn	•	•	X	•	X	•	•	•	11 000 011 ← n → ← n →	C3	3	3	10	
JP cc, nn	If condition cc is true PC ← nn, otherwise continue	•	•	X	•	X	•	•	•	11 cc 010 ← n → ← n →		3	3	10	cc Condition 000 NZ non-zero 001 Z zero 010 NC non-carry 011 C carry 100 PO parity odd 101 PE parity even 110 P sign positive 111 M sign negative
JR e	PC ← PC+e	•	•	X	•	X	•	•	•	00 011 000 ← e−2 →	18	2	3	12	
JR C, e	If C = 0, continue	•	•	X	•	X	•	•	•	00 111 000 ← e−2 →	38	2	2	7	If condition not met.
	If C = 1, PC ← PC+e											2	3	12	If condition is met.
JR NC, e	If C = 1, continue	•	•	X	•	X	•	•	•	00 110 000 ← e−2 →	30	2	2	7	If condition not met.
	If C = 0, PC ← PC+e											2	3	12	If condition is met.
JP Z, e	If Z = 0 continue	•	•	X	•	X	•	•	•	00 101 000 ← e−2 →	28	2	2	7	If condition not met.
	If Z = 1, PC ← PC+e											2	3	12	If condition is met.
JR NZ, e	If Z = 1, continue	•	•	X	•	X	•	•	•	00 100 000 ← e−2 →	20	2	2	7	If condition not met.
	If Z = 0, PC ← PC+e											2	3	12	If condition is met.
JP (HL)	PC ← HL	•	•	X	•	X	•	•	•	11 101 001	E9	1	1	4	
JP (IX)	PC ← IX	•	•	X	•	X	•	•	•	11 011 101 11 101 001	DD E9	2	2	8	

Jump Group (Continued)

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	Opcode 76 543 210	Hex	No.of Bytes	No.of M Cycles	No.of T States	Comments
JP (IY)	PC ← IY	•	•	X	•	X	•	•	•	11 111 101 11 101 001	FD E9	2	2	8	
DJNZ, e	B ← B - 1 If B = 0, continue	•	•	X	•	X	•	•	•	00 010 000 ← e-2 →	10	2	2	8	If B = 0.
	If B ≠ 0, PC ← PC+e											2	3	13	If B ≠ 0.

NOTES. e represents the extension in the relative addressing mode.
e is a signed two's complement number in the range < -126, 129 >.
e-2 in the opcode provides an effective address of pc+e as PC is incremented by 2 prior to the addition of e.

Call and Return Group

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	Opcode 76 543 210	Hex	No.of Bytes	No.of M Cycles	No.of T States	Comments
CALL nn	(SP-1) ← PC_H (SP-2) ← PC_L PC ← nn	•	•	X	•	X	•	•	•	11 001 101 ← n → ← n →	CD	3	5	17	
CALL cc, nn	If condition cc is false continue, otherwise same as CALL nn	•	•	X	•	X	•	•	•	11 cc 100 ← n →		3	3	10	If cc is false.
										← n →		3	5	17	If cc is true.
RET	PC_L ← (SP) PC_H ← (SP+1)	•	•	X	•	X	•	•	•	11 001 001	C9	1	3	10	
RET cc	If condition cc is false continue, otherwise same as RET	•	•	X	•	X	•	•	•	11 cc 000		1	1	5	If cc is false.
												1	3	11	If cc is true. cc Condition 000 NZ non-zero 001 Z zero 010 NC non-carry 011 C carry 100 PO parity odd 101 PE parity even 110 P sign positive 111 M sign negative
RETI	Return from interrupt	•	•	X	•	X	•	•	•	11 101 101 01 001 101	ED 4D	2	4	14	
RETN[1]	Return from non-maskable interrupt	•	•	X	•	X	•	•	•	11 101 101 01 000 101	ED 45	2	4	14	
RST p	(SP-1) ← PC_H (SP-2) ← PC_L PC_H ← 0 PC_L ← p	•	•	X	•	X	•	•	•	11 t 111		1	3	·11	t p 000 00H 001 08H 010 10H 011 18H 100 20H 101 28H 110 30H 111 38H

NOTE: [1]RETN loads IFF_2 → IFF_1

Input and Output Group

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	Opcode 76 543 210	Hex	No.of Bytes	No.of M Cycles	No.of T States	Comments
IN A, (n)	A ← (n)	•	•	X	•	X	•	•	•	11 011 011 ← n →	DB	2	3	11	n to $A_0 \sim A_7$ Acc. to $A_8 \sim A_{15}$
IN r, (C)	r ← (C) if r = 110 only the flags will be affected	↕	↕	X	↕	X	P	0	•	11 101 101 01 r 000	ED	2	3	12	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
INI	(HL) ← (C) B ← B-1 HL ← HL + 1	X	① ↕	X	X	X	X	1	X	11 101 101 10 100 010	ED A2	2	4	16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
INIR	(HL) ← (C) B ← B-1 HL ← HL + 1 Repeat until B = 0	X	② 1	X	X	X	X	1	X	11 101 101 10 110 010	ED B2	2 2	5 (If B≠0) 4 (If B=0)	21 16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
IND	(HL) ← (C) B ← B -1 HL ← HL-1	X	① ↕	X	X	X	X	1	X	11 101 101 10 101 010	ED AA	2	4	16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
INDR	(HL) ← (C) B ← B-1 HL ← HL-1 Repeat until B = 0	X	② 1	X	X	X	X	1	X	11 101 101 10 111 010	ED BA	2 2	5 (If B≠0) 4 (If B=0)	21 16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
OUT (n), A	(n) ← A	•	•	X	•	X	•	•	•	11 010 011 ← n →	D3	2	3	11	n to $A_0 \sim A_7$ Acc. to $A_8 \sim A_{15}$
OUT (C), r	(C) ← r	•	•	X	•	X	•	•	•	11 101 101 01 r 001	ED	2	3	12	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
OUTI	(C) ← (HL) B ← B-1 HL ← HL + 1	X	① ↕	X	,X	X	X	1	X	11 101 101 10 100 011	ED A3	2	4	16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
OTIR	(C) ← (HL) B ← B-1 HL ← HL + 1 Repeat until B = 0	X	② 1	X	X	X	X	1	X	11 101 101 10 110 011	ED B3	2 2	5 (If B≠0) 4 (If B=0)	21 16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$
OUTD	(C) ← (HL) B ← B-1 HL ← HL-1	X	① ↕	X	X	X	X	1	X	11 101 101 10 101 011	ED AB	2	4	16	C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$

NOTE: ① If the result of B-1 is zero the Z flag is set, otherwise it is reset.
② Z flag is set upon instruction completion only.

Input and Output Group (Continued)

Mnemonic	Symbolic Operation	S	Z		H		P/V	N	C	Opcode 76 543 210 Hex	No.of Bytes	No.of M Cycles	No.of T States	Comments
OTDR	(C) ← (HL) B ← B − 1 HL ← HL − 1 Repeat until B = 0	X	1	X	X	X	X	↕	X	11 101 101 ED 10 111 011	2 2	5 (If B≠0) 4 (If B=0)	21 16	C to A_0 ~ A_7 B to A_8 ~ A_{15}

Summary of Flag Operation

Instruction	D_7 S	Z		H		P/V	N	D_0 C	Comments
ADD A, s; ADC A, s	↕	↕	X	↕	X	V	0	↕	8-bit add or add with carry.
SUB s; SBC A, s; CP s; NEG	↕	↕	X	↕	X	V	1	↕	8-bit subtract, subtract with carry, compare and negate accumulator.
AND s	↕	↕	X	1	X	P	0	0	Logical operations.
OR s, XOR s	↕	↕	X	0	X	P	0	0	
INC s	↕	↕	X	↕	X	V	0	•	8-bit increment.
DEC s	↕	↕	X	↕	X	V	1	•	8-bit decrement.
ADD DD, ss	•	•	X	X	X	•	0	↕	16-bit add.
ADC HL, ss	↕	↕	X	X	X	V	0	↕	16-bit add with carry.
SBC HL, ss	↕	↕	X	X	X	V	1	↕	16-bit subtract with carry.
RLA, RLCA, RRA; RRCA	•	•	X	0	X	•	0	↕	Rotate accumulator.

RL m; RLC m; RR m;
RRC m; SLA m;
SRA m; SRL m
	↕	↕	X	0	X	P	0	↕	Rotate and shift locations.

RLD; RRD	↕	↕	X	0	X	P	0	•	Rotate digit left and right.
DAA	↕	↕	X	↕	X	P	•	↕	Decimal adjust accumulator.
CPL	•	•	X	1	X	•	1	•	Complement accumulator.
SCF	•	•	X	0	X	•	0	1	Set carry.
CCF	•	•	X	X	X	•	0	↕	Complement carry.
IN r (C)	↕	↕	X	0	X	P	0	•	Input register indirect.
INI, IND, OUTI; OUTD	X	↕	X	X	X	X	1	•	Block input and output. Z = 0 if B ≠ 0 otherwise Z = 0.
INIR; INDR; OTIR; OTDR	X	1	X	X	X	X	1	•	
LDI; LDD	X	X	X	0	X	↕	0	•	Block transfer instructions. P/V = 1 if BC ≠ 0, otherwise P/V = 0.
LDIR; LDDR	X	X	X	0	X	0	0	•	
CPI; CPIR; CPD; CPDR	X	↕	X	X	X	↕	1	•	Block search instructions. Z = 1 if A = (HL), otherwise Z = 0. P/V = 1 if BC ≠ 0, otherwise P/V = 0.
LD A, I, LD A, R	↕	↕	X	0	X	IFF	0	•	The content of the interrupt enable flip-flop (IFF) is copied into the P/V flag.
BIT b, s	X	↕	X	1	X	X	0	•	The state of bit b of location s is copied into the Z flag.

Symbolic Notation

Symbol	Operation
S	Sign flag. S = 1 if the MSB of the result is 1.
Z	Zero flag. Z = 1 if the result of the operation is 0.
P/V	Parity or overflow flag. Parity (P) and overflow (V) share the same flag. Logical operations affect this flag with the parity of the result while arithmetic operations affect this flag with the overflow of the result. If P/V holds parity, P/V = 1 if the result of the operation is even, P/V = 0 if result is odd. If P/V holds overflow, P/V = 1 if the result of the operation produced an overflow.
H	Half-carry flag. H = 1 if the add or subtract operation produced a carry into or borrow from bit 4 of the accumulator.
N	Add/Subtract flag. N = 1 if the previous operation was a subtract.
H & N	H and N flags are used in conjunction with the decimal adjust instruction (DAA) to properly correct the result into packed BCD format following addition or subtraction using operands with packed BCD format.
C	Carry/Link flag. C = 1 if the operation produced a carry from the MSB of the operand or result.

Symbol	Operation
↕	The flag is affected according to the result of the operation.
•	The flag is unchanged by the operation.
0	The flag is reset by the operation.
1	The flag is set by the operation.
X	The flag is indeterminate.
V	P/V flag affected according to the overflow result of the operation.
P	P/V flag affected according to the parity result of the operation.
r	Any one of the CPU registers A, B, C, D, E, H, L.
s	Any 8-bit location for all the addressing modes allowed for the particular instruction.
ss	Any 16-bit location for all the addressing modes allowed for that instruction.
ii	Any one of the two index registers IX or IY.
R	Refresh counter.
n	8-bit value in range < 0, 255 >.
nn	16-bit value in range < 0, 65535 >.

Pin Descriptions

A_0-A_{15}. *Address Bus* (output, active High, 3-state). A_0-A_{15} form a 16-bit address bus. The Address Bus provides the address for memory data bus exchanges (up to 64K bytes) and for I/O device exchanges.

$\overline{\text{BUSACK}}$. *Bus Acknowledge* (output, active Low). Bus Acknowledge indicates to the requesting device that the CPU address bus, data bus, and control signals $\overline{\text{MREQ}}$, $\overline{\text{IORQ}}$, $\overline{\text{RD}}$, and $\overline{\text{WR}}$ have entered their high-impedance states. The external circuitry can now control these lines.

$\overline{\text{BUSREQ}}$. *Bus Request* (input, active Low). Bus Request has a higher priority than $\overline{\text{NMI}}$ and is always recognized at the end of the current machine cycle. $\overline{\text{BUSREQ}}$ forces the CPU address bus, data bus, and control signals $\overline{\text{MREQ}}$, $\overline{\text{IORQ}}$, $\overline{\text{RD}}$, and $\overline{\text{WR}}$ to go to a high-impedance state so that other devices can control these lines. $\overline{\text{BUSREQ}}$ is normally wire-ORed and requires an external pullup for these applications. Extended $\overline{\text{BUSREQ}}$ periods due to extensive DMA operations can prevent the CPU from properly refreshing dynamic RAMs.

D_0-D_7. *Data Bus* (input/output, active High, 3-state). D_0-D_7 constitute an 8-bit bidirectional data bus, used for data exchanges with memory and I/O.

$\overline{\text{HALT}}$. *Halt State* (output, active Low). $\overline{\text{HALT}}$ indicates that the CPU has executed a Halt instruction and is awaiting either a non-maskable or a maskable interrupt (with the mask enabled) before operation can resume. While halted, the CPU executes NOPs to maintain memory refresh.

$\overline{\text{INT}}$. *Interrupt Request* (input, active Low). Interrupt Request is generated by I/O devices. The CPU honors a request at the end of the current instruction if the internal software-controlled interrupt enable flip-flop (IFF) is enabled. $\overline{\text{INT}}$ is normally wire-ORed and requires an external pullup for these applications.

$\overline{\text{IORQ}}$. *Input/Output Request* (output, active Low, 3-state). $\overline{\text{IORQ}}$ indicates that the lower half of the address bus holds a valid I/O address for an I/O read or write operation. $\overline{\text{IORQ}}$ is also generated concurrently with $\overline{\text{M1}}$ during an interrupt acknowledge cycle to indicate that an interrupt response vector can be placed on the data bus.

$\overline{\text{M1}}$. *Machine Cycle One* (output, active Low). $\overline{\text{M1}}$, together with $\overline{\text{MREQ}}$, indicates that the current machine cycle is the opcode fetch cycle of an instruction execution. $\overline{\text{M1}}$, together with $\overline{\text{IORQ}}$, indicates an interrupt acknowledge cycle.

$\overline{\text{MREQ}}$. *Memory Request* (output, active Low, 3-state). $\overline{\text{MREQ}}$ indicates that the address bus holds a valid address for a memory read or memory write operation.

$\overline{\text{NMI}}$. *Non-Maskable Interrupt* (input, negative edge-triggered). $\overline{\text{NMI}}$ has a higher priority than $\overline{\text{INT}}$. $\overline{\text{NMI}}$ is always recognized at the end of the current instruction, independent of the status of the interrupt enable flip-flop, and automatically forces the CPU to restart at location 0066H.

$\overline{\text{RD}}$. *Read* (output, active Low, 3-state). $\overline{\text{RD}}$ indicates that the CPU wants to read data from memory or an I/O device. The addressed I/O device or memory should use this signal to gate data onto the CPU data bus.

$\overline{\text{RESET}}$. *Reset* (input, active Low). $\overline{\text{RESET}}$ initializes the CPU as follows: it resets the interrupt enable flip-flop, clears the PC and Registers I and R, and sets the interrupt status to Mode 0. During reset time, the address and data bus go to a high-impedance state, and all control output signals go to the inactive state. Note that $\overline{\text{RESET}}$ must be active for a minimum of three full clock cycles before the reset operation is complete.

$\overline{\text{RFSH}}$. *Refresh* (output, active Low). $\overline{\text{RFSH}}$, together with $\overline{\text{MREQ}}$, indicates that the lower seven bits of the system's address bus can be used as a refresh address to the system's dynamic memories.

$\overline{\text{WAIT}}$. *Wait* (input, active Low). $\overline{\text{WAIT}}$ indicates to the CPU that the addressed memory or I/O devices are not ready for a data transfer. The CPU continues to enter a Wait state as long as this signal is active. Extended $\overline{\text{WAIT}}$ periods can prevent the CPU from refreshing dynamic memory properly.

$\overline{\text{WR}}$. *Write* (output, active Low, 3-state). $\overline{\text{WR}}$ indicates that the CPU data bus holds valid data to be stored at the addressed memory or I/O location.

CPU Timing

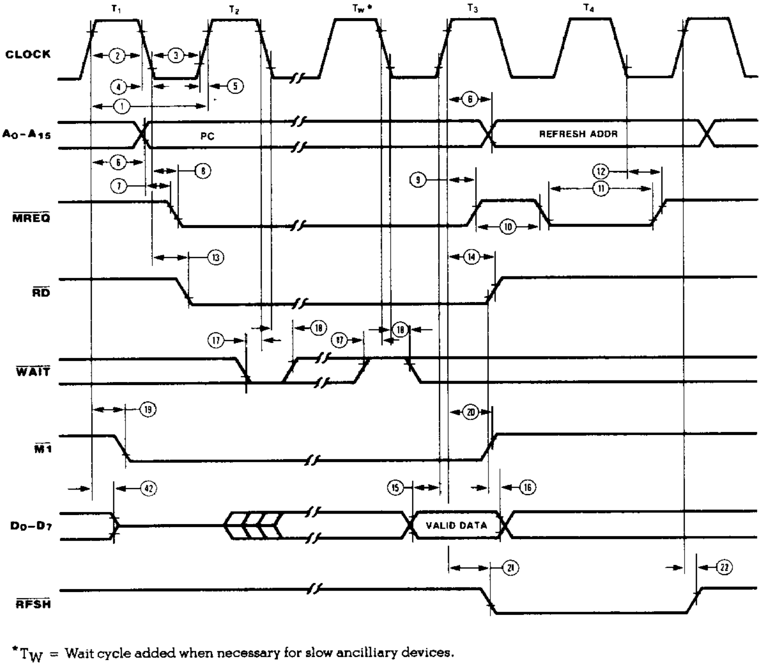
The CPU executes instructions by proceeding through a specific sequence of operations:

- Memory read or write
- I/O device read or write
- Interrupt acknowledge

The basic clock period is referred to as a T time or cycle, and three or more T cycles make up a machine cycle (M1, M2 or M3 for instance). Machine cycles can be extended either by the CPU automatically inserting one or more Wait states or by the insertion of one or more Wait states by the user.

Instruction Opcode Fetch. The CPU places the contents of the Program Counter (PC) on the address bus at the start of the cycle (Figure 5). Approximately one-half clock cycle later, $\overline{MREQ}$ goes active. When active, $\overline{RD}$ indicates that the memory data can be enabled onto the CPU data bus.

The CPU samples the $\overline{WAIT}$ input with the falling edge of clock state T_2. During clock states T_3 and T_4 of an M1 cycle dynamic RAM refresh can occur while the CPU starts decoding and executing the instruction. When the Refresh Control signal becomes active, refreshing of dynamic memory can take place.



*T_W = Wait cycle added when necessary for slow ancilliary devices.

Figure 5. Instruction Opcode Fetch

CPU Timing (Continued)

Memory Read or Write Cycles. Figure 6 shows the timing of memory read or write cycles other than an opcode fetch ($\overline{\text{M1}}$) cycle. The $\overline{\text{MREQ}}$ and $\overline{\text{RD}}$ signals function exactly as in the fetch cycle. In a memory write cycle, $\overline{\text{MREQ}}$ also becomes active when the address bus is stable. The $\overline{\text{WR}}$ line is active when the data bus is stable, so that it can be used directly as an R/$\overline{\text{W}}$ pulse to most semiconductor memories.

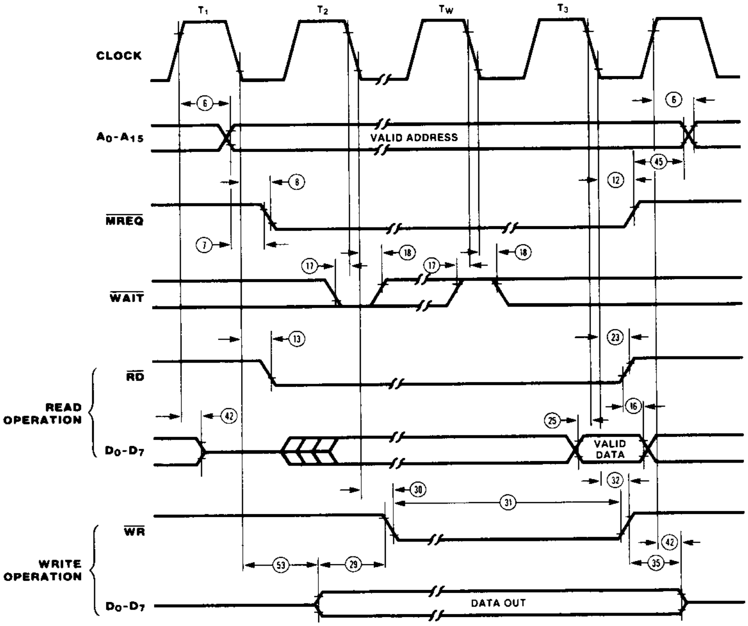
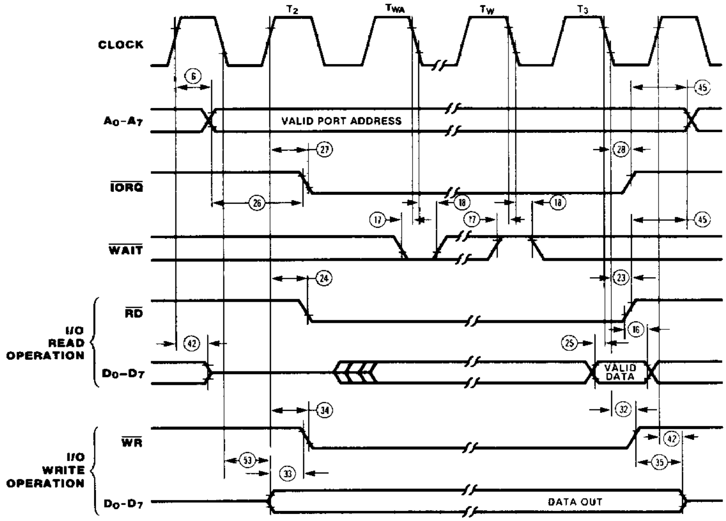


Figure 6. Memory Read or Write Cycles

CPU Timing (Continued)

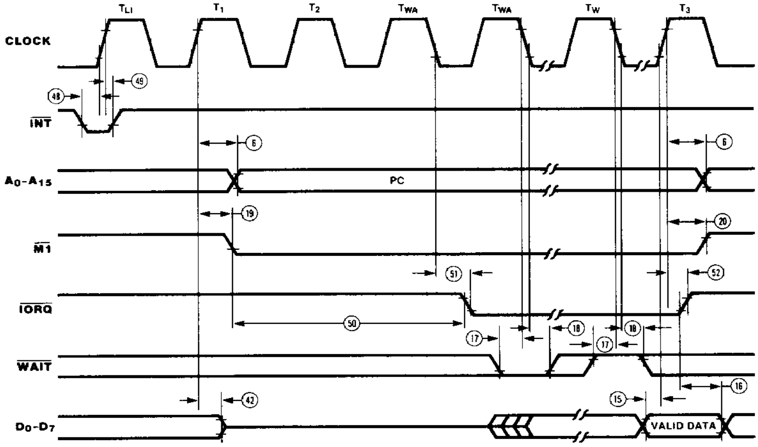
Input or Output Cycles. Figure 7 shows the timing for an I/O read or I/O write operation. During I/O operations, the CPU automatically inserts a single Wait state (T_{WA}). This extra Wait state allows sufficient time for an I/O port to decode the address from the port address lines.



T_{WA} = Wait cycle automatically inserted by CPU.

Figure 7. Input or Output Cycles

Interrupt Request/Acknowledge Cycle. The CPU samples the interrupt signal with the rising edge of the last clock cycle at the end of any instruction (Figure 8). When an interrupt is accepted, a special $\overline{\text{M1}}$ cycle is generated. During this $\overline{\text{M1}}$ cycle, $\overline{\text{IORQ}}$ becomes active (instead of $\overline{\text{MREQ}}$) to indicate that the interrupting device can place an 8-bit vector on the data bus. The CPU automatically adds two Wait states to this cycle.

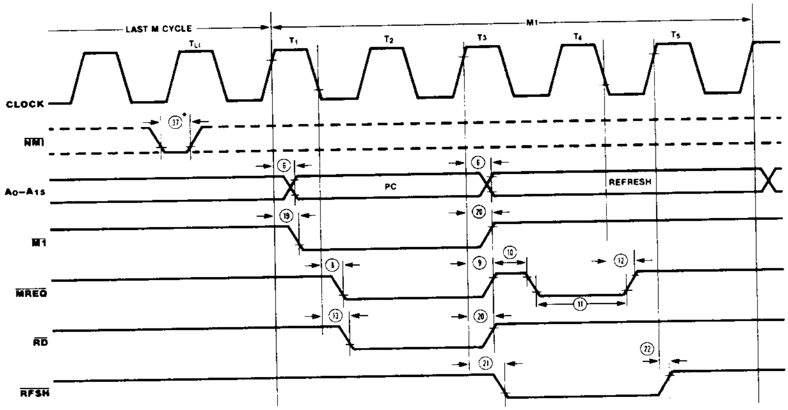


T_{LI} = Last state of any instruction cycle.
T_{WA} = Wait cycle automatically inserted by CPU.

Figure 8. Interrupt Request/Acknowledge Cycle

CPU Timing (Continued)

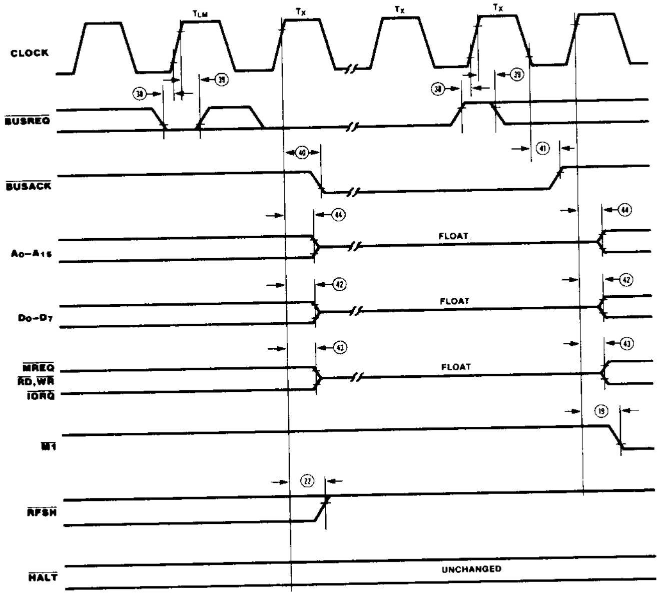
Non-Maskable Interrupt Request Cycle. $\overline{NMI}$ is sampled at the same time as the maskable interrupt input $\overline{INT}$ but has higher priority and cannot be disabled under software control. The subsequent timing is similar to that of a normal memory read operation except that data put on the bus by the memory is ignored. The CPU instead executes a restart (RST) operation and jumps to the $\overline{NMI}$ service routine located at address 0066H (Figure 9).



* Although $\overline{NMI}$ is an asynchronous input, to guarantee its being recognized on the following machine cycle, $\overline{NMI}$'s falling edge must occur no later than the rising edge of the clock cycle preceding the last state of any instruction cycle (T_{LI}).

Figure 9. Non-Maskable Interrupt Request Operation

Bus Request/Acknowledge Cycle. The CPU samples $\overline{BUSREQ}$ with the rising edge of the last clock period of any machine cycle (Figure 10). If $\overline{BUSREQ}$ is active, the CPU sets its address, data, and $\overline{MREQ}$, $\overline{IORQ}$, $\overline{RD}$, and $\overline{WR}$ lines to a high-impedance state with the rising edge of the next clock pulse. At that time, any external device can take control of these lines, usually to transfer data between memory and I/O devices.



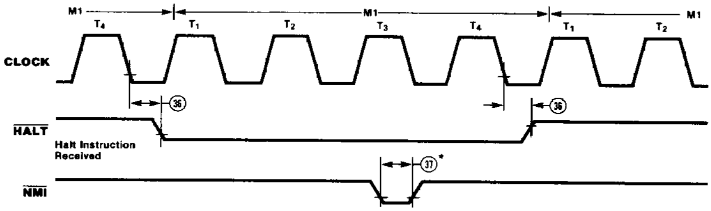
NOTES: 1) T_{LM} = Last state of any M cycle.

2) T_X = An arbitrary clock cycle used by requesting device.

Figure 10. Z-BUS Request/Acknowledge Cycle

CPU Timing (Continued)

Halt Acknowledge Cycle. When the CPU receives a $\overline{\text{HALT}}$ instruction, it executes NOP states until either an $\overline{\text{INT}}$ or $\overline{\text{NMI}}$ input is received. When in the Halt state, the $\overline{\text{HALT}}$ output is active and remains so until an interrupt is processed (Figure 11). $\overline{\text{INT}}$ will also force a Halt exit.



* Although $\overline{\text{NMI}}$ is an asynchronous input, to guarantee its being recognized on the following machine cycle, $\overline{\text{NMI}}$'s falling edge must occur no later than the rising edge of the clock cycle preceding the last state of any instruction cycle (T_{LI}).

Figure 11. Halt Acknowledge Cycle

Reset Cycle. $\overline{\text{RESET}}$ must be active for at least three clock cycles for the CPU to properly accept it. As long as $\overline{\text{RESET}}$ remains active, the address and data buses float, and the control outputs are inactive. Once $\overline{\text{RESET}}$ goes inactive, two internal T cycles are consumed before the CPU resumes normal processing operation. $\overline{\text{RESET}}$ clears the PC register, so the first opcode fetch will be to location 0000 (Figure 12).

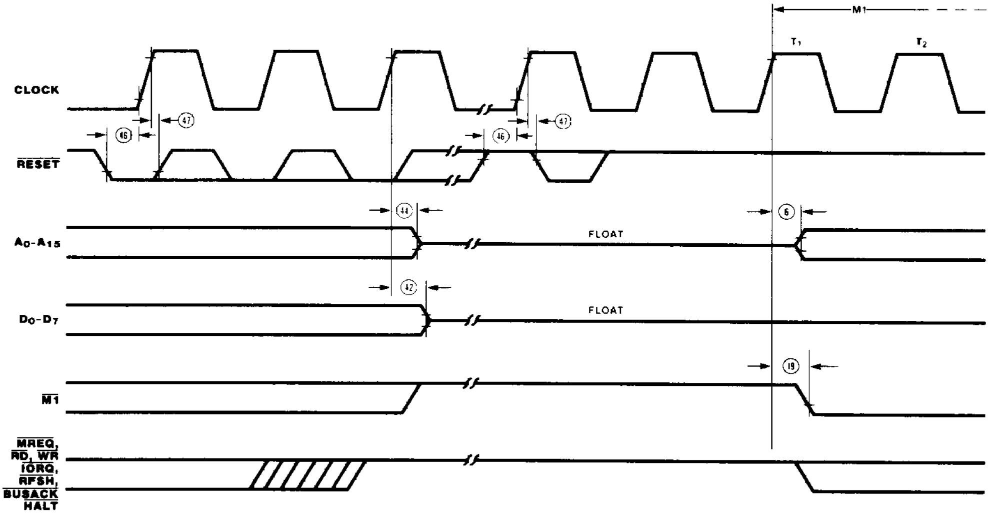

M1
T1
T2
CLOCK
46
47
46
47
$\overline{\text{RESET}}$
44
6
A_0-A_{15}
FLOAT
42
D_0-D_7
FLOAT
19
$\overline{\text{M1}}$
$\overline{\text{MREQ}}$, $\overline{\text{RD}}$, $\overline{\text{WR}}$
$\overline{\text{IORQ}}$, $\overline{\text{RFSH}}$, $\overline{\text{BUSACK}}$
$\overline{\text{HALT}}$

Figure 12. Reset Cycle

AC Characteristics†

Number	Symbol	Parameter	Z8300-1 (1.0 MHz) Min (ns)	Z8300-1 (1.0 MHz) Max (ns)	Z8300-3 (2.5 MHz) Min (ns)	Z8300-3 (2.5 MHz) Max (ns)
1	TcC	Clock Cycle Time	1000*		400*	
2	TwCh	Clock Pulse Width (High)	470	2000	180	2000
3	TwCl	Clock Pulse Width (Low)	470	2000	180	2000
4	TfC	Clock Fall Time	—	30	—	30
5	TrC	Clock Rise Time		30		30
6	TdCr(A)	Clock ↑ to Address Valid Delay	—	380	—	145
7	TdA(MREQf)	Address Valid to $\overline{\text{MREQ}}$ ↓ Delay	370*	—	125*	—
8	TdCf(MREQf)	Clock ↓ to $\overline{\text{MREQ}}$ ↓ Delay	—	260	—	100
9	TdCr(MREQr)	Clock ↑ to $\overline{\text{MREQ}}$ ↑ Delay	—	260	—	100
10	TwMREQh	$\overline{\text{MREQ}}$ Pulse Width (High)	410*		170*	
11	TwMREQl	$\overline{\text{MREQ}}$ Pulse Width (Low)	890*	—	360*	—
12	TdCf(MREQr)	Clock ↓ to $\overline{\text{MREQ}}$ ↑ Delay	—	260	—	100
13	TdCf(RDf)	Clock ↓ to $\overline{\text{RD}}$ ↓ Delay	—	340	—	130
14	TdCr(RDr)	Clock ↑ to $\overline{\text{RD}}$ ↑ Delay	—	260	—	100
15	TsD(Cr)	Data Setup Time to Clock ↑	140		50	
16	ThD(RDr)	Data Hold Time to $\overline{\text{RD}}$ ↑	—	0	—	0
17	TsWAIT(Cf)	$\overline{\text{WAIT}}$ Setup Time to Clock ↓	190	—	70	—
18	ThWAIT(Cf)	$\overline{\text{WAIT}}$ Hold Time after Clock ↓	—	0	—	0
19	TdCr(M1f)	Clock ↑ to $\overline{\text{M1}}$ ↓ Delay	—	340	—	130
20	TdCr(M1r)	Clock ↑ to $\overline{\text{M1}}$ ↑ Delay		340		130
21	TdCr(RFSHf)	Clock ↑ to $\overline{\text{RFSH}}$ ↓ Delay	—	460	—	180
22	TdCr(RFSHr)	Clock ↑ to $\overline{\text{RFSH}}$ ↑ Delay	—	390	—	150
23	TdCf(RDr)	Clock ↓ to $\overline{\text{RD}}$ ↑ Delay	—	290	—	110
24	TdCr(RDf)	Clock ↑ to $\overline{\text{RD}}$ ↓ Delay	—	260	—	100
25	TsD(Cf)	Data Setup to Clock ↓ during M_2, M_3, M_4 or M_5 Cycles	160		60	
26	TdA(IORQf)	Address Stable prior to $\overline{\text{IORQ}}$ ↓	790'	—	320*	—
27	TdCr(IORQf)	Clock ↑ to $\overline{\text{IORQ}}$ ↓ Delay	—	240	—	90
28	TdCf(IORQr)	Clock ↓ to $\overline{\text{IORQ}}$ ↑ Delay	—	290	—	110
29	TdD(WRf)	Data Stable prior to $\overline{\text{WR}}$ ↓	470*	—	190*	—
30	TdCf(WRf)	Clock ↓ to $\overline{\text{WR}}$ ↓ Delay		240		90
31	TwWR	$\overline{\text{WR}}$ Pulse Width	890*	—	360*	—
32	TdCf(WRr)	Clock ↓ to $\overline{\text{WR}}$ ↑ Delay	—	260	—	100
33	TdD(WRf)	Data Stable prior to $\overline{\text{WR}}$ ↓	-30*	—	30*	—
34	TdCr(WRf)	Clock ↑ to $\overline{\text{WR}}$ ↓ Delay	—	210	—	80
35	TdWRr(D)	Data Stable from $\overline{\text{WR}}$ ↑	290*		130*	
36	TdCf(HALT)	Clock ↓ to $\overline{\text{HALT}}$ ↑ or ↓	—	760	—	300
37	TwNMI	$\overline{\text{NMI}}$ Pulse Width	210	—	80	—
38	TsBUSREQ(Cr)	$\overline{\text{BUSREQ}}$ Setup Time to Clock ↑	210	—	80	—

*For clock periods other than the minimums shown in the table, calculate parameters using the expressions in the table on the following page.

Calculated values above assumed.

TrC = TfC = 20 ns.

†All timings assume equal loading on pins within 50 pf.

AC Characteristics† (Continued)

Number	Symbol	Parameter	Z8300-1 Min (ns)	Z8300-1 Max (ns)	Z8300-3 Min (ns)	Z8300-3 Max (ns)
39	ThBUSREQ(Cr)	$\overline{BUSREQ}$ Hold Time after Clock ↑	0	—	0	—
40	TdCr(BUSACKf)	Clock ↑ to $\overline{BUSACK}$ ↓ Delay		310		120
41	TdCf(BUSACKr)	Clock ↓ to $\overline{BUSACK}$ ↑ Delay	—	290	—	110
42	TdCr(Dz)	Clock ↑ to Data Float Delay	—	240	—	90
43	TdCr(CTz)	Clock ↑ to Control Outputs Float Delay ($\overline{MREQ}$, $\overline{IORQ}$, $\overline{RD}$, and $\overline{WR}$)	—	290	—	110
44	TdCr(Az)	Clock ↑ to Address Float Delay	—	290	—	110
45	TdCTr(A)	$\overline{MREQ}$ ↑, $\overline{IORQ}$ ↑, $\overline{RD}$ ↑, and $\overline{WR}$ ↑ to Address Hold Time	400*		160*	
46	TsRESET(Cr)	$\overline{RESET}$ to Clock ↑ Setup Time	240	—	90	—
47	ThRESET(Cr)	$\overline{RESET}$ to Clock ↑ Hold Time	—	0	—	0
48	TsINTf(Cr)	$\overline{INT}$ to Clock ↑ Setup Time	210	—	80	—
49	ThINTr(Cr)	$\overline{INT}$ to Clock ↑ Hold Time	—	0	—	0
50	TdM1f(IORQf)	M1 ↓ to $\overline{IORQ}$ ↓ Delay	2300*		920*	
51	TdCf(IORQf)	Clock ↓ to $\overline{IORQ}$ ↓ Delay	—	290	—	110
52	TdCf(IORQr)	Clock ↑ to $\overline{IORQ}$ ↑ Delay	—	260	—	100
53	TdCf(D)	Clock ↓ to Data Valid Delay	—	290	—	230

*For clock periods other than the minimums shown in the table, calculate parameters using the following expressions. Calculated values above assumed TrC = TfC = 20 ns.

† All timings assume equal loading on pins with 50 pF.

Footnotes to AC Characteristics

Number	Symbol	Z8300-1	Z8300-3
1	TcC	TwCh + TwCl + TrC + TfC	TwCh + TwCl + TrC + TfC
7	TdA(MREQf)	TwCh + TfC − 200	TwCh + TfC − 75
10	TwMREQh	TwCh + TfC − 90	TwCh + TfC − 30
11	TwMREQl	TcC − 110	TcC − 30
26	TdA(IORQf)	TcC − 210	TcC − 80
29	TdD(WRf)	TcC − 540	TcC − 210
31	TwWR	TcC − 110	TcC − 40
33	TdD(WRf)	TwCl + TrC − 470	TwCl + TrC − 180
35	TdWRr(D)	TwCl + TrC − 210	TwCl + TrC − 80
45	TdCTr(A)	TwCl + TrC − 110	TwCl + TrC − 40
50	TdM1f(IORQf)	2TcC + TwCh + TfC − 210	2TcC + TwCh + TfC − 80

AC Test Conditions:
V_{IH} = 2.0 V
V_{IL} = 0.8 V
V_{IHC} = V_{CC} −0.6 V
V_{ILC} = 0.45 V
V_{OH} = 2.0 V
V_{OL} = 0.8 V
FLOAT = ±0.5 V

Absolute Maximum Ratings

Voltages on all pins with respect
to ground . . . . −0.3V to +7V
Operating Ambient
Temperature . . . . . . .See Ordering Information
Storage Temperature. . . . . −65°C to +150°C

Stresses greater than those listed under Absolute Maximum Ratings may cause permanent damage to the device. This is a stress rating only; operation of the device at any condition above those indicated in the operational sections of these specifications is not implied. Exposure to absolute maximum rating conditions for extended periods may affect device reliability.

Standard Test Conditions

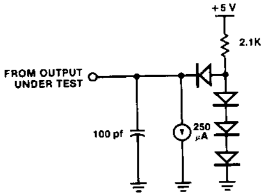
The DC characteristics and capacitance sections listed below apply for the following standard test conditions, unless otherwise noted. All voltages are referenced to GND (0V). Positive current flows into the referenced pin.

Available operating temperature is:

- S = 0°C to +70°C, +4.75V ≤ V_{CC} ≤ +5.25V

All ac parameters assume a load capacitance of 100 pf. Add 10 ns delay for each 50 pf increase in load up to a maximum of 200 pf for the data bus and 100 pf for address and control lines.

The Ordering Information section lists package temperature ranges and product numbers. Package drawings are in the Package Information section. Refer to the Literature List for additional documentation.



DC Characteristics

Symbol	Parameter	Min	Max	Unit	Test Condition
V_{ILC}	Clock Input Low Voltage	-0.3	0.45	V	
V_{IHC}	Clock Input High Voltage	V_{CC}-.6	V_{CC}+.3	V	
V_{IL}	Input Low Voltage	-0.3	0.8	V	
V_{IH}	Input High Voltage	2.0	V_{CC}	V	
V_{OL}	Output Low Voltage		0.4	V	I_{OL} = 2.0 mA
V_{OH}	Output High Voltage	2.4		V	I_{OH} = -250 μA
I_{LI}	Input Leakage Current		10	μA	V_{IN} = 0 to V_{CC}
I_{LO}	3-State Output Leakage		±10[1]	μA	V_{OUT} = 0.4 to V_{CC}
I_{CC}	Power Supply Current				

	Temperature				
Frequency	0°C Max	25°C Max	25°C Typical	70°C Max	Unit
Z8300-1 (1.0 MHz)	30	25	15	20	mA
Z8300-3 (2.5 MHz)	45	40	25	35	mA

[1] A_{15}-A_0, D_7-D_0, $\overline{MREQ}$, $\overline{IORQ}$, $\overline{RD}$, and $\overline{WR}$.

Capacitance

Symbol	Parameter	Min	Max	Unit	Note
C_{CLOCK}	Clock Capacitance		35	pF	
C_{IN}	Input Capacitance		5	pF	
C_{OUT}	Output Capacitance		15	pF	

T_A = 25°C, f = 1 MHz.
Unmeasured pins returned to ground.

ORDERING INFORMATION

Z80L CPU, 1.0 MHz
40-pin DIP
Z8300-1 PS

Z80L CPU, 2.5 MHz
40-pin DIP
Z8300-3 PS

Codes

First letter is for package; second letter is for temperature.

C = Ceramic DIP
P = Plastic DIP
L = Ceramic LCC
V = Plastic PCC

R = Protopack
T = Low Profile Protopack
DIP = Dual-In-Line Package
LCC = Leadless Chip Carrier
PCC = Plastic Chip Carrier (Leaded)

TEMPERATURE
S = 0°C to +70°C
E = −40°C to +85°C
M*= −55°C to +125°C

FLOW
B = 883 Class B

Example: PS is a plastic DIP, 0°C to +70°C.

†Available soon.

*For Military Orders, contact your local Zilog Sales Office for Military Electrical Specifications.