

Description

The μ PB8289 bus arbiter is used with the μ PB8288 bus controller to interface 8086 and 8088 microprocessors to a multimaster system bus. The μ PB8289 controls the μ PB8288 bus controller and the bus transceivers and address latches, preventing them from accessing the system bus if the processor does not have use of the bus.

An external command sequence will cause the associated microprocessor to enter a wait state until the bus is ready. The processor remains in the wait state until the bus arbiter acquires use of the multimaster system bus. Then, the arbiter allows the bus controller, data transceivers, and address latches to access the system.

Once use of the bus has been acquired and data has been transferred, transfer acknowledge (XACK) is returned to the processor to indicate that the accessed slave device is ready. The processor may then complete its transfer cycle.

Features

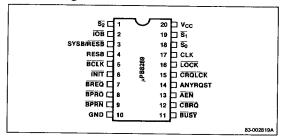
- ☐ Multimaster system bus protocol
- 8086 and 8088 processor synchronization with multimaster bus
- ☐ Simple interface with the 8288 bus controller and 8283/8282 address latches to a system bus
- ☐ Four operating modes for flexible system configuration
- ☐ Simplified interface to Multibus® systems
- ☐ Parallel, serial, and rotating priority resolution
- ☐ Bipolar buffering and drive capability

Ordering Information

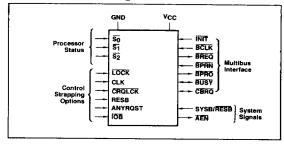
Part	Package	Max Frequency	
Number	Type	of Operation	
μPB8289D	20-Pin Cerdip	8 MHz	

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Pin Configuration



Functional Configuration



Pin Identification

No.	Symbol	Function
1, 18, 19	\bar{S}_0 - \bar{S}_2	Status inputs
2	IOB	I/O bus input
3	SYSB/RESB	System bus/resident bus inputs
4	RESB	Resident bus input
5	BCLK	System bus clock input
6	INIT	Initialize input
7	BREQ	Bus request output
8	BPRO	Bus priority output
9	BPRN	Bus priority input
10	GND	Ground
11	BUSY	Bus interface signal I/O
12	CBRQ	Common bus request I/O
13	AEN	Address enable output
14	ANYRQST	Any request input
15	CROLCK	Common request lock input
16	LOCK	Lock input
17	CLK	Clock input
20	V _{CC}	+5 V power supply



Pin Functions

So-S2 (Status Inputs)

The µPB8289 decodes these status inputs from the 8086 or 8088 processor to begin bus requests and surrenders.

IOB (I/O Bus)

This input signal tells the μ PB8289 that there is an I/O peripheral bus and a multimaster system bus.

SYSB/RESB (System Bus/Resident Bus)

This input determines when bus requests and surrenders are permitted in SR mode.

RESB (Resident Bus Input)

 $\overline{\text{RESB}}$ tells the μPB8289 that there is a multimaster and resident bus. When the signal is high, the SYSB/ $\overline{\text{RESB}}$ pin handles bus arbitration.

BLCK (System Bus Clock)

This clock input synchronize all system bus interface signals.

INIT (Initialize)

This active low-input resets all bus arbiters on the multimaster bus. No arbiters have use of the bus following INIT.

BREQ (Bus Request)

An arbiter uses this output to request use of the multimaster system bus.

BPRO (Bus Priority Output)

In serial priority resolving schemes, this output daisychains to BPRN of the next lower priority arbiter.

BPRN (Bus Priority Input)

This input tells the arbiter it may acquire the bus on the next falling edge at BCLK.

BUSY (Bus Interface Signal)

When the bus is available, this I/O signal notifies all arbiters on the bus. The highest requesting arbiter seizes the bus and pulls BUSY low to keep other arbiters off the bus.

CBRQ (Common Bus Request)

This signal is an input from a lower priority arbiter requesting the bus. It is an output from arbiters that surrender the multimaster bus upon request.

AEN (Address Enable)

This output tells the 8288 bus controller, 8284 clock driver, and the processor's address latches when to tri-state their output drivers.

ANYRQST (Any Request)

This signal allows the multimaster bus to be surrendered to a lower priority arbiter.

CRQLCK (Common Request Lock)

This input prevents the μ PB8289 from surrendering the bus in response to a request on the \overline{CBRQ} input.

LOCK (Lock)

This input prevents the arbiter from surrendering the multimaster system bus to any other bus arbiter, regardless of its priority.

CLK (Clock)

This is the clock signal from the 8284 clock generator.

GND (Ground)

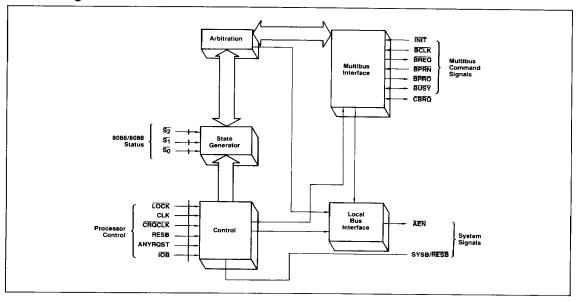
This is the ground.

V_{CC} (Power Supply)

This is the +5 V power supply.



Block Diagram



Functional Description

Bus Master Arbitration

Higher priority masters generally acquire use of the bus when a lower priority master completes its present transfer cycle. Lower priority masters acquire the bus when no higher priority master is accessing the system bus. The ANYRQST strapping option allows the arbiter to surrender the bus to a lower priority master as if it were a higher priority master. The arbiter maintains the bus as long as no other bus masters are requesting the bus and its processor has not entered the halt state. The arbiter does not voluntarily surrender the bus and must be forced off by a request from another bus master, unless the arbiter's processor has entered the halt state. Additional strapping options allow for other sets of conditions.

Priority Resolving Techniques

The μ PB8289 provides several techniques for resolving priority between the many possible bus masters of a multimaster system bus. All of these techniques assume that one bus master will have priority over all others at any given time. You may use parallel, serial, or rotating priority resolving.

Parallel Priority Resolving

This technique (figures 1, 2) uses a bus request line (BREQ) for each arbiter on the multimaster system bus. Each BREQ line goes to a priority encoder that generates the address of the highest priority active BREQ line. This binary address is decoded to select the bus priority in line (BPRN) that is returned to the highest priority arbiter. The arbiter that receives priority (BPRN true) allows its bus master onto the multimaster system bus as soon as the bus becomes available. An arbiter that gets priority over another arbiter cannot immediately seize the bus, but must wait until the current bus transaction is complete. When the transaction is complete, the current occupant of the bus surrenders the bus by releasing BUSY. BUSY is an active low OR tied line which goes to every arbiter on the system bus. When BUSY goes high (inactive), the priority arbiter seizes the bus and brings BUSY low to keep other arbiters off the bus. Note that all multimaster system bus transactions are synchronized to the bus clock (BCLK).



Figure 1. Parallel Priority Resolving

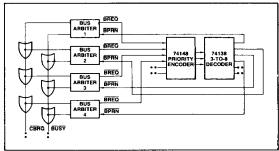
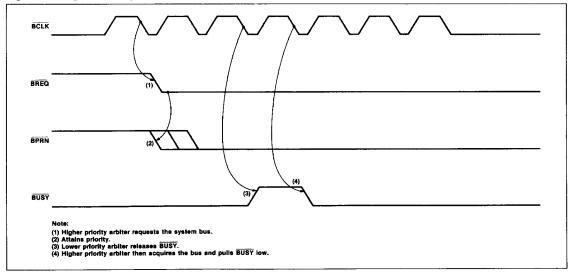


Figure 2. Higher Priority Arbiter Obtaining the Bus from a Lower Priority Arbiter

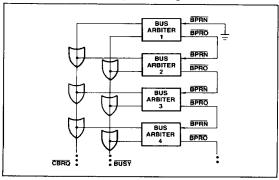




Serial Priority Resolving

The serial priority resolving technique (figure 3) daisychains the bus arbiters together by connecting the higher priority arbiter's \overline{BPRQ} output to the \overline{BPRN} of the next lowest priority arbiter. This eliminates the need for the priority encoder-decoder arrangement. The number of arbiters that may be daisy-chained together is a function of \overline{BCLK} and the propagation delay from arbiter to arbiter. At 10 MHz, only 3 arbiters may be daisy-chained.

Figure 3. Serial Priority Resolving



Rotating Priority Resolving

This technique resembles the parallel priority resolving technique except that priority is dynamically reassigned. The priority encoder is replaced by a circuit that rotates priority between arbiters to allow each arbiter an equal chance to use the system bus.

Modes of Operation

The μ PB8289 has two basic operating modes: I/O peripheral bus mode (\overline{IOB} mode), and resident bus mode (RESB mode). The \overline{IOB} strapping option configures the μ PB8289 into \overline{IOB} mode and the RESB strapping option configures it to RESB mode. If both options are strapped false, the arbiter interfaces the processor to a multimaster system bus only. If both options are strapped true, the arbiter interfaces the processor to a multimaster system bus, a resident bus, and an I/O bus. Figure 4 shows the μ PB8289 in a typical CPU system.

IOB Mode

 $\overline{\text{IOB}}$ mode allows the processor to access both an I/O peripheral bus and a multimaster system bus. On an I/O peripheral bus, all devices on the bus, including memory, are treated as I/O devices and addressed by I/O commands. All memory commands are directed to the multimaster system bus. In $\overline{\text{IOB}}$ mode, the processor communicates with and controls peripherals over the peripheral bus and communicates with system memory over the system memory bus. Figure 5 shows the μPB8289 in this mode.

RESB Mode

RESB mode allows the processor to communicate over both a resident bus and a multimaster system bus. A resident bus can issue memory and I/O commands, but it is separate from the multimaster system bus. The resident bus has one master and is dedicated to only that master. The 8086 and 8088 can communicate with a resident bus and a multimaster system bus. The processor can access the memory and peripherals of both buses. Memory mapping selects which bus is accessed. The SYSB/RESB input on the arbiter instructs the arbiter on which bus to access. The signal connected to SYSB/RESB also enables and disables commands from one of the bus controllers. Figure 6 shows the $\mu PB8289$ in this mode.



Mode Summary

	Status Lines From 8886 or 8088 or 8089		IOB Mode RESB (Mode) Only <u>IOB</u> Mode RESB M Only <u>IOB</u> = High RESB = High IOB = Low RESB =						
	<u>\$2</u>	<u>\$1</u>	šō	IOB = Low	SYSB/RESB = High	SYSB/RESB = Low	SYSB/RESB = High	SYSB/RESS = Low	
	0	0	0	х	~	x	x	×	
1/0 commands	0	0	1	X	✓	X	x	x	✓
	0	11	0	X		×	х	x	-
Halt	0	1	1	X	X	x	x	x	×
Memory commands	1	0	0	~	∠	x	_	x	_
	1	0	1	~	~	x	_	x	· /
	11	1	0	~		x		x	-
ldie	1	1	1	х	x	×	х	х	x

Notes:

(1) x = Multimaster system bus is allowed to be surrendered.

(2) = Multimaster system bus is requested.

Multimaster System Bus

Mode	Pin Strapping	Requested (1)	Surrendered (2)
Single bus multimaster mode	IOB = high RESB = low	When the pro- cessor's status lines go active	HLT + Ti • HPBRQ†
RESB mode only	10B = high RESB = high	SYSB/RESB ≠ High 2 active	(SYSB/RESB = low + TI) CBRQ + HLT + HPBRQ
IOB mode only	IOB = low RESB = low	Memory commands	(I/O status + TI) • CBRQ + HLT + HPBRQ
IOB mode • RESB mode	10B = low RESB = high	(Memory command) • (SYSB/RESB = high)	(I/O status commands) + (TI) (SYSB/RESB = low) • CBRQ + HPBRQ† + HLT)

Note:

- (1) Except for HALT and idle status.
- (2) LOCK prevents surrender of bus to any other arbiter. CRQLCK prevents surrender of bus to a lower priority arbiter.
- (3) HLT = processor halt; $\overline{S}_2 \overline{S}_0 = 011$.
- (4) TI = processor idle; $\overline{S}_2 \overline{S}_0 = 111$.
- (5) + means OR.
- (6) means AND.
 - † HPBRQ = higher priority bus request or BPRN = 1.

Absolute Maximum Ratings

 $T_A = 25^{\circ}$

Operating temperature	0°C to 70°C
Storage temperature	-65°C to +150°C
Voltage on any pin	-0.5 V to +7 V
All input voltages	-1.0 V to +5.5 V
Power dissipation	1.5 W

Comment: Exposing the device to stresses above those listed in Absolute Maximum Ratings could cause permanent damage. The device is not meant to be operated under conditions outside the limits described in the operational sections of this specification. Exposure to absolute maximum rating conditions for extended periods may affect device reliability.

DC Characteristics

 $T_A = 0$ °C to +70°C; $V_{CC} = 5 \text{ V} \pm 10\%$

	Limits					Test	
Parameter	Symbol	Min	Тур	Max	Unit	Conditions	
Input low voltage	VIL			0.8	٧		
Input high voltage	VIH	2.0			٧		
Input clamp voltage	V _C			-1.0	٧	$V_{CC} = 4.50 \text{ V},$ $I_{C} = -5 \text{ mA}$	
Input forward current	ΙĖ			-0.5	mA	$V_{CC} = 5.50 \text{ V},$ $V_F = 0.45 \text{ V}$	
Reverse input leakage current	IR			60	μΑ	$V_{CC} = 5.50 \text{ V}$ $V_{R} = 5.50 \text{ V}$	
Output low voltage	v_{0L}			*			
BUSÝ, CBRQ				0.45	٧	$I_{OL} = 20 \text{ mA}$	
AEN DEEC				0.45	٧	$1_{OL} = 16 \text{ mA}$	
BPRO, BREQ				0.45	٧	$1_{OL} = 10 \text{ mA}$	
Output high_	v_{oh}	0pe	n colle	ctor			
voltage BUSY, CBRQ		2.4			٧	$I_{OH} = 400 \mu A$	
All other outputs							
Power supply current	lcc			165	mA		



Capacitance

			Limits			Test
Parameter	Symbol	Min Typ		Max	Unit	Conditions
Input capacitance	C _{IN} status 25			pF		
Input capacitance	C _{IN} (others)			12	pF	
Note:						
1) BUSY, CBRQ	2) AEN			3)	BPRO,	BREQ
2.3 V		2.3	v			2.3 V
≸ 92.5 Ω		≱ 11	ΩΩ			≱ 170 Ω
	·	-		-		
☆ 250 pF		十10) pF			±60 pF
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Figure 4. Typical CPU System Using the μPD8289 Bus Arbiter

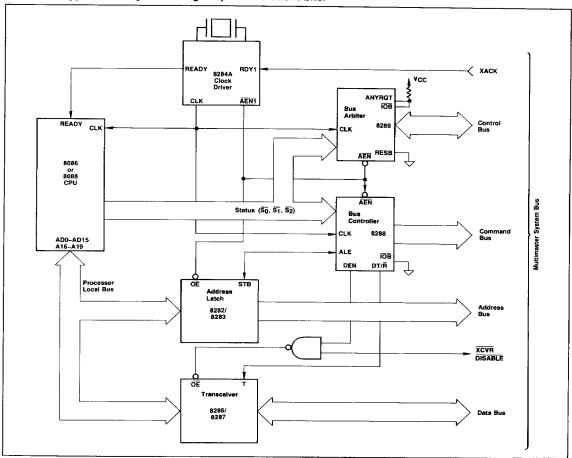




Figure 5. Typical Medium-Complexity IOB System

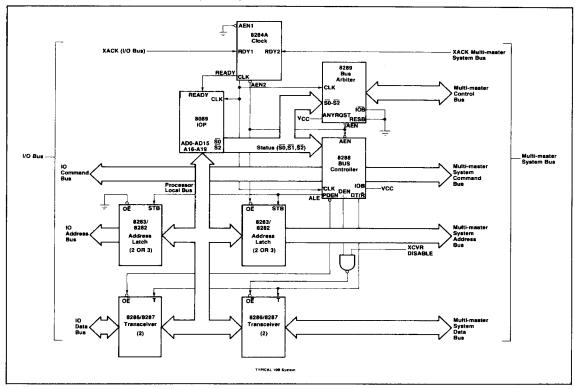
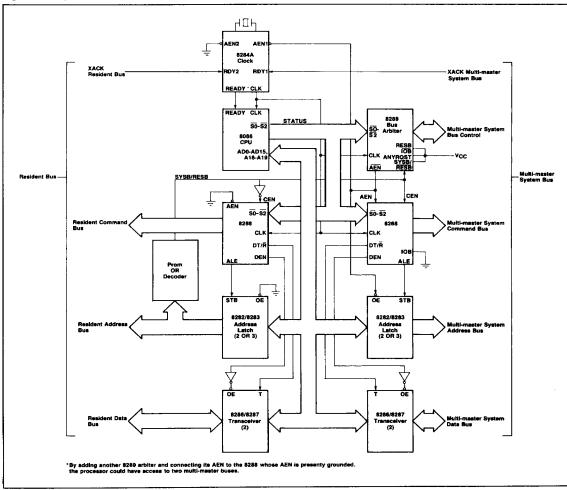




Figure 6. Typical System, Resident Bus Configuration





AC Characteristics

 $T_A = 0^{\circ}C \text{ to } + 70^{\circ}C; V_{CC} = 5V \pm 10\%$

		Test				
Parameter	Symbo	l Min	Тур	Max	Unit	
CLK cycle period	t _{CLCL}	125			ns	
CLK low time	tCLCH	65			ns	
CLK high time	t _{CHCL}	35			ns	
Status active setup	tsvch	65	t _C	LCL - 10	ns	
Status inactive setup	tshcl	50	t _C	LCL - 10	ns	
Status active hold	thvch	10			ns	
Status inactive hold	tHVCL	10			ns	
BUSY∱↓ setup to BCLK↓	†BYSBL	20			ns	
CBRQ†↓ setup to BCLK↓	tcbsbl	20			ns	
BCLK cycle time	tBLBL	100			ns	
BCLK high time	t _{BHCL}	30	0.6	5 (t _{BLBL}) ns	
LOCK inactive hold	t _{CLLL1}	10			ns	
LOCK active setup	t _{CLLL2}	40			ns	
BPRN+1 to BCLK setup time	t _{PNBL}	15			ns	
SYSB/RESB setup	[†] CLSR1	0			ns	
SYSB/RESB hold	t _{CLSR2}	20			ns	
nitialization oulse width	t _{IVIH}	3 t _{BLBL} 3 t _{CLCL}	+		ns	
nput rise time	t _{ILIH}			20	ns	From 0.8 V to 2.0 V
nput fall time	t _{IHIL}			12	ns	From 2.0 V to 0.8 V

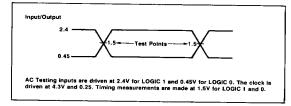
Timing Response

			Limits			Test	
Parameter	Symbol	Min	Тур	Max	Unit	Conditions	
BCLK to BREQ delay (1)	t _{BLBRL}			35	ns		
BCLK to BPRO (1)(2)	t _{BLPOH}			40	ns		
BPRN↓↑ to BPRO↓↑ delay ₍₁₎₍₂₎	†PNP0			25	ns		
BCLK to BUSY low	t _{BLBYL}			60	ns		
BCLK to BUSY float (3)	t _{BLBYH}			35	ns		
CLK to AEN high	tCLAEH			65	ns		
BCLK to AEN low	t _{BLAEL}			40	ns		
BCLK to CBRQ low	t _{BLCBL}			60	ns		
BCLK to CBRQ float (3)	t _{RLCRH}			35	ns		
Output rise time	toLoH			20	ns	From 0.8 V to 2.0 V	
Output fall time	toHOL			12	ns	From 2.0 V to 0.8 V	

Note:

- Denotes that the spec applies to both transitions of the signal.
- (2) BCLK generates the first BPRO. Subsequent changes of BPRO are generated through BPRON.
- (3) Measured at 0.5 V above GND.

AC Test Condition





Timing Waveforms

The signals related to CLK are typical processor signals and do not relate to the depicted sequence of events of the signals referenced to BCLK. The signals shown related to the BCLK represent a hypothetical sequence of events for illustration. Assume three bus arbiters of priorities 1, 2, and 3 configured in the serial priority resolving scheme.

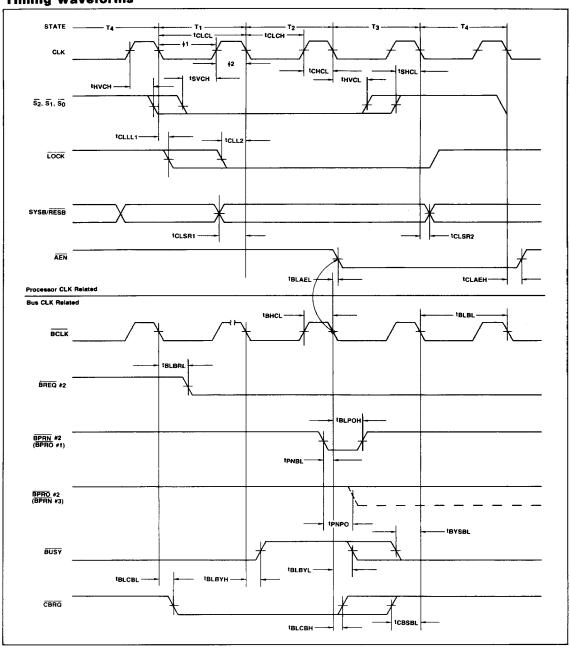
Assume arbiter 1 has the bus and is holding BUSY low. Arbiter 2 detects its processor wants the bus and pulls BREQ #2 low. If BPRN #2 is high (as shown), arbiter 2 pulls CBRQ low. CBRQ signals to higher-priority arbiter 1 that a lower-priority arbiter wants the bus. A higher-priority arbiter would be given BPRN when it makes the bus request rather than having to wait for another arbiter to release the bus through CBRQ.

Arbiter 1 relinquishes the multimaster system bus when it enters a state of not requiring it, by lowering its BPRO #1 (tied to BPRN #2) and releasing BUSY. Arbiter 2 now sees that it has priority from BPRN #2 being low and releases CBRQ. As soon as BUSY signifies the bus is available (high), arbiter 2 pulls BUSY low on the next falling edge of BCLK.

Note that if arbiter 2 didn't want the bus at the time it received priority, it would pass priority to the next lower priority by lowering its BPRO #2 (TPNPO). Note also that even a higher-priority arbiter which is aquiring the bus through BPRN will momentarily drop CBRQ until it has acquired the bus.







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