

FEATURES

- ☐ Comprehensive MIL-STD-1553B dual-redundant Bus Controller (BC) and Remote Terminal (RT) and Monitor (M) functions
- ☐ Multiple message processing capability in BC, RT, and M modes
- ☐ Automatic polling and intermessage delay in BC mode
- Programmable interrupt scheme and internally generated interrupt history list
- ☐ Register-oriented architecture to enhance programmability

- ☐ DMA or pseudo-dual-port memory interface with 64K addressability
- ☐ Internal wraparound self-test
- ☐ Time tagging and message logging in RT and M modes
- ☐ Packaged in 84-pin pingrid array, 84-lead flatpack, or 84-pad LCC
- ☐ Low-power CMOS technology
- ☐ Full military operating temperature range, -55°C to +125°C, screened to the specific test methods listed in Table I of MIL-STD-883, Method 5004, Class B
- ☐ Architecturally and pin-compatible with UTMC's UT1553B BCRT

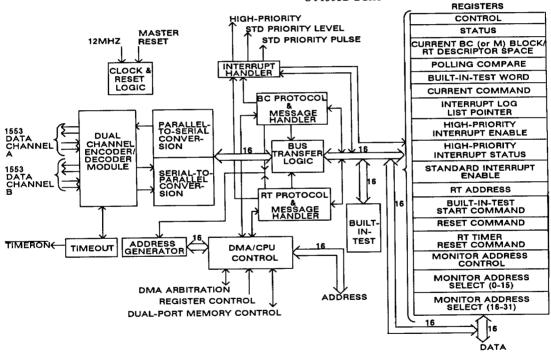


Figure 1. BCRTM Block Diagram

Table of Contents

1.0	INTRODUCTION	
	1.1 Features - Remote Terminal (RT) Mode	. 3
	1.3 Features - Monitor (M) Mode	. 4
2.0	PIN IDENTIFICATION AND DESCRIPTION	. 5
3.0	INTERNAL REGISTERS	12
4.0	SYSTEM OVERVIEW	20
5.0	SYSTEM INTERFACE	20
	5.1 DMA Transfers	20
	5.2 Hardware Interface	21
	5.3 CPU Interconnection	
	5.4 RAM Interface	
	5.5 Transmitter/Receiver Interface	22
6.0	REMOTE TERMINAL ARCHITECTURE	23
0.0		24
	6.1.1 RT Subaddress Descriptor Definitions	24
	6.1.2 Message Status Word	25
		26
	6.2 RT Error Detection	28
	6.3 RT Operational Sequence	28
7.0	BUS CONTROLLER ARCHITECTURE	29
,	7.1 BC Functional Operation	30
	7.2 Polling	
	7.3 BC Error Detection	
	7.4 BC Operational Sequence	32
	7.5 BC Operational Example	
8.0	MONITOR ARCHITECTURE	۰.
0.U	8.1 Monitor Functional Operation	
	8.2 Monitor Error Detection	
	8.3 Monitor Operational Sequence	
	5.5 Montor Operational Sequence	30
9.0	EXCEPTION HANDLING AND INTERRUPT LOGGING	37
10.0	MAXIMUM AND RECOMMENDED OPERATING CONDITIONS RATINGS \dots	40
11.0	DC ELECTRICAL CHARACTERISTICS	41
12.0	AC ELECTRICAL CHARACTERISTICS	42
12.0	DACKACE OUTLINE DRAWINGS	

1.0 Introduction

The monolithic CMOS UT1553B BCRTM provides the system designer with an intelligent solution to MIL-STD-1553B multiplexed serial data bus design problems. The UT1553B BCRTM is a single-chip device that implements all three of the defined MIL-STD-1553B functions - Bus Controller. Remote Terminal, and Monitor. Designed to reduce host CPU overhead, the BCRTM's powerful state machines automatically execute message transfers, provide interrupts, and generate status information. Multiple registers offer many programmable functions as well as extensive information for host use. In the BC mode, the BCRTM uses a linked-list message scheme to provide the host with message chaining capability. The BCRTM enhances memory use by supporting variable-size, relocatable data blocks. In the RT mode, the BCRTM implements time-tagging and message history functions. It also supports multiple (up to 128) message buffering and variable length messages to any subaddress. In the Monitor (M) mode, the BCRTM's powerful linked list command block structure allows it to process a series of monitored 1553 messages without the intervention of the host. The BCRTM can store as much bus traffic as can be contained in its 64K memory space. In addition, the host has the capability of instructing the BCRTM to monitor and store data for only selected remote terminals.

The UT1553B BCRTM is an intelligent, versatile, and easy to implement device -- a powerful asset to system designers.

1.1 Features - Remote Terminal (RT) Mode

Indexing

The BCRTM is programmable to index or buffer messages on a subaddress-by-subaddress basis. The BCRTM, which can index as many as 128 messages, can also assert an interrupt when either the selected number of messages is reached or every time a specified subaddress is accessed.

Variable Space Allocation

The BCRTM can use as little or as much memory (up to 64K) as needed.

Selectable Data Storage

Address programmability within the BCRTM provides flexible data placement and convenient access.

Sequential Data Storage

The BCRTM stores/retrieves, by subaddress, all messages in the order in which they are transacted.

Sequential Message Status Information

The BCRTM provides message validity, time-tag, and word-count information, and stores it sequentially in a separate, cross-referenced list.

Illegalizing Mode Codes and Subaddresses

The host can declare mode codes and subaddresses illegal by setting the appropriate bit(s) in memory.

Programmable Interrupt Selection

The host CPU can select various events to cause an interrupt with provision for high and standard priority interrupts.

Interrupt History List

The BCRTM provides an Interrupt History List that records, in the order of occurrence, the events that caused the interrupts. The list length is programmable.

1.2 Features - Bus Controller (BC) Mode

Multiple Message Processing

The BCRTM autonomously processes any number of messages or lists of messages that may be stored in a 64K memory space.

Automatic Intermessage Delay

When programmed by the host, the BCRTM can delay a host-specified time before executing the next message in sequence.

Automatic Polling

When polling, the BCRTM interrogates the remote terminals and then compares their status word responses to the contents of the Polling Compare Register. The BCRTM can interrupt the host CPU if an erroneous remote terminal status word response occurs.

Automatic Retry

The BCRTM can automatically retry a message on busy, message error, and/or response time-out conditions. The BCRTM can retry up to four times on the same or on the alternate bus.

Programmable Interrupt Selection

The host CPU can select various events to cause an interrupt with provision for high and standard priority interrupts.

Interrupt History List

The BCRTM provides an Interrupt History List that records, in the order of occurrence, the events that caused the interrupts. The list length is programmable.

Variable Space Allocation

The BCRTM uses as little or as much memory (up to 64K) as needed.

Selectable Data Storage

Address programmability within the BCRTM provides flexible data placement and convenient access.

1.3 Features - Monitor (M) Mode

Command History List

The BCRTM's linked list command block structure permits the BCRTM to process a series of monitored messages without host intervention.

Monitor Selected Terminal Addresses

The host can select the remote terminals to be monitored by programming the proper bits in the Terminal Address Select registers (Registers 16 and 17). The BCRTM can monitor any or all remote terminals.

Variable Space Allocation

The BCRTM can use as little or as much memory (up to 64K) as needed.

Selectable Data Storage

Address programmability within the BCRTM provides flexible data placement and convenient access.

Sequential Data Storage

The BCRTM stores, by Terminal Address, all 1553 messages in the order in which they are transacted.

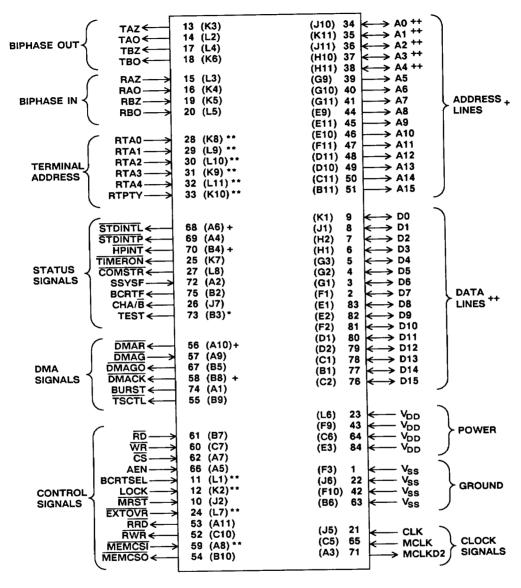
Programmable Interrupt Selection

The host can select a wide variety of events that may cause an interrupting event.

Interrupt History List

The BCRTM stores, chronologically in memory, an Interrupt History List of each event that caused an interrupt.

2.0 PIN IDENTIFICATION AND DESCRIPTION



^{••} Pin internally pulled up.

Pingrid array pin identification in parentheses.
 LCC, flatpack pin number not in parentheses.

Figure 2. BCRTM Functional Pin Description

BCRTM-5

⁺ Pin at high impedance when not asserted

^{+ +} Bidirectional pin.

Formerly MEMWIN

Legend for TYPE and ACTIVE fields:

TUI = TTL input (pull-up)

AL = Active low

AH = Active high

ZL = Active low - inactive state is high impedance

TI = TTL input

TO = TTL output

TTO = Three-state TTL output

TTB = Bidirectional

Notes:

1. Address and data buses are in the high-impedance state when idle.

2. Flatpack pin numbers are same as LCC.

ADDRESS BUS

NAME	PIN NI LCC	UMBER PGA	TYPE	ACTIVE	DESCRIPTION
A0	34	J10	ттв		Bit 0 (LSB) of the Address bus.
A1	35	K11	ттв		Bit 1 of the Address bus.
A2	36	J11	ТТВ		Bit 2 of the Address bus.
A3	37	H10	ттв		Bit 3 of the Address bus.
A4	38	H11	ТТВ		Bit 4 of the Address bus.
A5	39	G9	тто		Bit 5 of the Address bus.
A6	40	G10	тто		Bit 6 of the Address bus.
A7	41	G11	тто		Bit 7 of the Address bus.
A8	44	E9	тто		Bit 8 of the Address bus.
A9	45	E11	тто		Bit 9 of the Address bus.
A10	46	E10	тто		Bit 10 of the Address bus.
A11	47	F11	TTO		Bit 11 of the Address bus.
A12	48	D11	тто		Bit 12 of the Address bus.
A13	49	D10	TTO		Bit 13 of the Address bus.
A14	50	C11	тто		Bit 14 of the Address bus.
A15	51	B11	TTO		Bit 15 (MSB) of the Address bus.

DATA BUS

NAME	PIN NU	MBER PGA	TYPE	ACTIVE	DESCRIPTION
D0	9	K 1	ттв		Bit 0 (LSB) of the Data bus.
D1	8	J1	ТТВ		Bit 1 of the Data bus.
D2	7	H2	ттв		Bit 2 of the Data bus.
D3	6	H1	ттв	•	Bit 3 of the Data bus.
D4	5	G3	ттв		Bit 4 of the Data bus.
D5	4	G2	ттв		Bit 5 of the Data bus.
D6	3	G1	ттв		Bit 6 of the Data bus.
D7	2	F1	ттв		Bit 7 of the Data bus.
D8	83	E1	TTB		Bit 8 of the Data bus.
D9	82	E2	ттв		Bit 9 of the Data bus.
D10	81	F2	ттв		Bit 10 of the Data bus.
D11	80	D1	ТТВ		Bit 11 of the Data bus.
D12	79	D2	ттв		Bit 12 of the Data bus.
D13	78	C1	ттв		Bit 13 of the Data bus.
D14	77	B1	TTB		Bit 14 of the Data bus.
D15	76	C2	ттв		Bit 15 (MSB) of the Data bus.

TERMINAL ADDRESS INPUTS

NAME	PIN NU	UMBER PGA	TYPE	ACTIVE	DESCRIPTION
RTA0	28	K8	TUI		Remote Terminal Address Bit 0 (LSB). The entire RT address is strobed in at Master Reset. Verify it by reading the Remote Terminal Address Register. All the Remote Terminal Address bits are internally pulled up.
RTA1	29	L9	TUI		Remote Terminal Address Bit 1. This is bit 1 of the Remote Terminal Address.
RTA2	30	L10	TUI		Remote Terminal Address Bit 2. This is bit 2 of the Remote Terminal Address.
RTA3	31	К9	TUI		Remote Terminal Address Bit 3. This is bit 3 of the Remote Terminal Address.
RTA4	32	L11	TUI		Remote Terminal Address Bit 4. This is bit 4 (MSB) of the Remote Terminal Address.
RTPTY	33	K10	TUI		Remote Terminal (Address) Parity. This is an odd parity input for the Remote Terminal Address.

CONTROL SIGNALS

NAME	PIN NU	JMBER PGA	ТҮРЕ	ACTIVE	DESCRIPTION
RD	61	В7	TI	AL	Read. The host uses this in conjunction with CS to read an internal BCRTM register.
WR	60	C7	TI	AL	Write. The host uses this in conjunction with CS to write to an internal BCRTM register.
CS	62	A7	TI	AL	Chip Select. This selects the BCRTM when accessing the BCRTM's internal register.
AEN	66	A5	TI	АН	Address Enable. The host CPU uses AEN to indicate to the BCRTM that the BCRTM's address lines can be asserted; this is a precautionary signal provided to avoid address bus crash. If not used, it must be tied high.
BCRTSEL	11	L1	TUI		BC/RT Select. This selects between either the Bus Controller or Remote Terminal mode. The BC/RT Mode Select bit in the Control Register overrides this input if the LOCK pin is not high. To operate in Monitor mode, the pin should be held low. This pin is internally pulled high.
LOCK	12	K2	TUI	АН	Lock. When set, this pin prevents internal changes to both the RT address and BC/RT mode select functions. This pin is internally pulled high.
EXTOVR	24	L7	TUI	AL	External Override. Use this in multi-redundant applications. Upon receipt, the BCRTM aborts all current activity. EXTOVR should be connected to COMSTR output of the adjacent BCRTM when used. This pin is internally pulled high.
MRST	10	J2	TI	AL	Master Reset. This resets all internal state machines, encoders, decoders, and registers. The minimum pulse width for a successful Master Reset is 500 ns.
MEMCSO	54	B10	то	AL	Memory Chip Select Out. This is the regenerated MEMCSI input for external RAM during the pseudo-dual-port RAM mode. The BCRTM also uses it to select external memory during memory accesses.
MEMCSI	59	A8	TI	AL	Memory Chip Select In. Used in the pseudodual-port RAM mode only, MEMCSI is received from the host and is propagated through to MEMCSO.
RRD	53	A11	то	AL	RAM Read. In the pseudo-dual-port RAM mode, the host uses this signal in conjunction with MEMSCO to read from external RAM through the BCRTM. It is also the signal the BCRTM uses to read from memory. It is asserted following receipt of DMAG. When the BCRTM performs multiple reads, this signal is pulsed.
RWR	52	C10	то	AL	RAM Write. In the pseudo-dual-port RAM mode, the CPU and BCRTM use this to write to external RAM. This signal is asserted following receipt of DMAG. For multiple writes, this signal is pulsed.

STATUS SIGNALS

NAME	PIN NU	JMBER PGA	TYPE	ACTIVE	DESCRIPTION
STDINTL	68	A 6	тто	ZL	Standard Interrupt Level. This is a level interrupt. It is asserted when one or more events enabled in either the Standard Interrupt Enable Register, RT Descriptor, or BC Command Block occur. Resetting the Standard Interrupt bit in the High-Priority Interrupt Status/Reset Register clears the interrupt.
STDINTP	69	A4	то	AL	Standard Interrupt Pulse. STDINTP pulses when an interrupt is logged.
HPINT	70	В4	тто	ZL	High-Priority Interrupt. The High-Priority Interrupt level is asserted upon occurence of events enabled in the High-Priority Interrupt Enable Register. The corresponding bit(s) in the High-Priority Interrupt Status/Reset Register reset HPINT.
TIMERON	25	K7	то	AL	(RT) Timer On. This is a 760-microsecond fail-safe transmitter enable timer. Started at the beginning of a transmission, TIMERON goes inactive 760 microseconds later or is reset automatically with the receipt of a new command. Use it in conjunction with CHA/B output to provide a fail-safe timer for channel A and B transmitters.
COMSTR	27	L8	то	AL	(RT) Command Strobe. The BCRTM asserts this signal after receiving a valid command. The BCRTM deactivates it after servicing the command.
SSYSF	72	A2	TI	АН	Subsystem Fail. Upon receipt, this signal propagates directly to the RT 1553 status word and the BCRTM Status Register.
BCRTF	75	В2	то	АН	BCRT Fail. This indicates a Built-In-Test (BIT) failure. In the RT mode, the Terminal Flag bit in 1553 status word is also set.
CHA/B	26	J7	то		Channel A/\overline{B} . This indicates the active or last active channel.
TEST	73	В3	то	AL	TEST. This pin is used as a factory test pin. (Formerly MEMWIN.)

BIPHASE INPUTS

NAME	PIN NU	JMBER PGA	TYPE	ACTIVE	DESCRIPTION
RAO	16	K4	TI		Receive Channel A One. This is the Manchester- encoded true signal input from Channel A of the bus receiver.
RAZ	15	L3	TI		Receive Channel A Zero. This is the Manchester- encoded complementary signal input from Channel A of the bus receiver.
RBO	20	L5	TI		Receive Channel B One. This is the Manchester- encoded true signal input from Channel B of the bus receiver.
RBZ	19	K5	TI		Receive Channel B Zero. This is the Manchester- encoded complementary signal input from Channel B of the bus receiver.

BCRTM-9

BIPHASE OUTPUTS

NAME	PIN N LCC	UMBER PGA	TYPE	ACTIVE	DESCRIPTION
TAO	14	L2	то	-	Transmit Channel A One. This is the Manchester- encoded true output to be connected to the Chan- nel A bus transmitter input. This signal is idle low.
TAZ	13	К3	то		Transmit Channel A Zero. This is the Manchester- encoded complementary output to be connected to the Channel A bus transmitter input. This signal is idle low.
ТВО	18	К6	то		Transmit Channel B One. This is the Manchester- encoded true output to be connected to the Chan- nel B bus transmitter input. This signal is idle low.
TBZ	17	L4	то		Transmit Channel B Zero. This is the Manchester- encoded complementary output to be connected to the Channel B bus transmitter input. This signal is idle low.

DMA SIGNALS

NAME	PIN NU LCC	JMBER PGA	ТҮРЕ	ACTIVE	DESCRIPTION
DMAR	56	A10	тто	ZL	DMA Request. The BCRTM issues this signal when access to RAM is required. It goes inactive after receiving a DMAG signal.
DMAG	57	A9	TI	AL	DMA Grant. This input to the BCRTM allows the BCRT to access RAM. It is recognized 45 ns before the rising edge of MCLKD2.
DMAGO	67	B 5	то	AL	DMA Grant Out. If DMAG is received but not needed, it passes through to this output.
DMACK	58	В8	тто	ZL	DMA Acknowledge. The BCRTM asserts this signal to confirm receipt of DMAG; it stays low until memory access is complete.
BURST	74	A 1	то	АН	Burst (DMA Cycle). This indicates that the current DMA cycle transfers at least two words; worst case is five words plus a "dummy" word.
TSCTL	55	В9	то	AL	Three-State Control. This signal indicates when the BCRTM is actually accessing memory. The host subsystem's address and data lines must be in the high-impedance state when the signal is active. This signal assists in placing the external data and address buffers into the high-impedance state.

CLOCK SIGNALS

NAME	PIN N LCC	UMBER PGA	TYPE	ACTIVE	DESCRIPTION
CLK	21	J 5	TI		Clock. The 12 MHz input clock requires a 50% +/-10% duty cycle with an accuracy of +/-0.01%. The accuracy is required in order to meet the Manchester encoding/decoding requirements of MIL-STD-1553B.
MCLK	65	C5	TI		Memory Clock. This is the input clock frequency the BCRTM uses for memory accesses. The memory cycle time is equal to two MCLK cycles. Therefore, RAM access time is dependent upon the chosen MCLK frequency (6 MHz minimum, 12 MHz maximum). Please see the BCRTM DMA timing diagrams in this data sheet.
MCLKD2	71	А3	то		Memory Clock Divided by Two. This signal is the Memory Clock input divided by two. It assists the host subsystem in synchronizing DMA events.

POWER AND GROUND

NAME	PIN N LCC	UMBER PGA	ТҮРЕ	ACTIVE	DESCRIPTION
V_{DD}	23	L6	PWR		+5V
V_{DD}	43	F9	PWR		+5V
V_{DD}	64	C6	PWR		+5V
V_{DD}	84	E3	PWR		+5V
Vss	1	F3	GND		Ground
V _{SS}	22	J6	GND		Ground
V _{SS}	42	F10	GND		Ground
V _{SS}	63	В6	GND		Ground

3.0 Internal Registers

Description

The BCRTM's internal registers (see table 1 on pages 18-19) enable the CPU to control the actions of the BCRTM while maintaining low DMA overhead by the BCRTM. All functions are active high and ignored

when low unless stated otherwise. Functions and parameters are used in both RT and BC modes except where indicated. Registers are addressed by the binary equivalent of their decimal number. For example, Register 1 is addressed as 0001B. Register usage is defined as follows:

#0 Control Register

D:4

Number	Description
BIT 15	Reserved.
BIT 14	RT Address 31. When RT31 = 0, the BCRTM recognizes RT Address 31 as a Broadcast command. When RT31 = 1, the BCRTM treats RT Address 31 as a normal terminal address.
BIT 13	Subaddress 31. When SA31 = 0, the BCRTM recognizes a command word with either subaddress 0 or 31 as being a valid mode code. When SA31 = 1, the BCRTM only recognizes a command word with a subaddress of 0 as a valid mode code.
BIT 12	Bus Controller Time-Out. When the BCRTM is a BC and BCTO = 0, the BCRTM allows an RT up to 16 μ s to respond with a status word before it declares a bus time-out. If BCTO = 1, the BCRTM allows an RT up to 32 μ s to respond with a status word before it declares a bus time-out.
BIT 11	Enable External Override. For use in multi-redundant systems. This bit enables the EXTOVR pin.
BIT 10	BC/RT Select. This function selects between the Bus Controller and Remote Terminal/Monitor operation modes. It overrides the external BCRTSEL input setting if the Change Lock-Out function is not used. A reset operation must be performed when changing between BC and RT/M modes. For Monitor operation this bit must be "0". This bit is write-only.

- BIT 9 (BC) Retry on Alternate Bus. This bit enables an automatic retry to operate on alternate buses. For example, if on bus A, with two automatic retries programmed, the automatic retries occur on bus B.
- BIT 8 (RT,M) Channel B Enable. When set, this bit enables Channel B operation. (BC) No significance.
- BIT 7 (RT,M) Channel A Enable. When set, this bit enables Channel A operation. (BC) Channel Select A/B. When set, this bit selects Channel A.
- BITs 6-5 (BC) Retry Count. These bits program the number (1-4) of retries to attempt. (00 = 1 retry, 11 = 4 retries)
- BIT 4 (BC) Retry on Bus Controller Message Error. This bit enables automatic retries on an error the Bus Controller detects (see the Bus Controller Architecture section, page 29).
- BIT 3 (BC) Retry on Time-out. This bit enables an automatic retry on a response time-out condition.
- BIT 2 (BC) Retry on Message Error. This bit enables an automatic retry when the Message Error bit is set in the RT's status word response.
- BIT 1 (BC) Retry on Busy. This bit enables automatic retry on a received Busy bit in an RT status word response.
- BIT 0 Start Enable. In the BC mode, this bit starts/restarts Command Block execution. In the RT or M mode, it enables the BCRTM to receive a valid command. RT or M operation does not start until a valid command is received. When using this function:
 - Restart the BCRTM after each Master Reset or programmed reset.
 - This bit is not readable; verify operation by reading bit 0 of the BCRTM's Status Register.

#1 Status Register (Read Only)

These bits indicate the BCRTM's current status.

Description Rit

Number

- TEST. This bit reflects the inverse of the TEST output. It changes state simultaneously with the **BIT 15** TEST output.
- (RT,M) Remote Terminal (or Monitor) Active. Indicates that the BCRTM, in the Remote Terminal **BIT 14** (or Monitor) mode, is presently servicing a command. This bit reflects the inverse of the COMSTR pin.
- (RT) Dynamic Bus Control Acceptance. This bit reflects the state of the Dynamic Bus Control **BIT 13** Acceptance bit in the RT status word (see Register 10 on page 16).
- (RT) Terminal Flag bit is set in RT status word. This bit reflects the result of writing to Register 10, **BIT 12** bit 11.
- (RT) Service Request bit is set in RT status word. This bit reflects the result of writing to **BIT 11** Register 10, bit 10.
- (RT) Busy bit is set in RT status word. This bit reflects the result of writing to Register 10, **BIT 10** bits 9 or 14.
- BIT 9 BIT is in progress.
- Reset is in progress. This bit indicates that either a write to Register 12 has just occurred or the BIT 8 BCRTM has just received a Reset Remote Terminal (#01000) Mode Code. This bit remains set less than one microsecond.
- $BC/(\overline{RT} \text{ or } \overline{M})$ Mode. Indicates the current mode of operation. A reset operation must be performed BIT 7 when changing between BC and RT modes.
- Channel A/\overline{B} . Indicates either the channel presently in use or the last channel used. BIT 6
- Subsystem Fail Indicator. Indicates receiving a subsystem fail signal from the host subsystem on the BIT 5 SSYSF input.
- Reserved. BITs 4-1
- (BC) Command Block Execution is in progress. (RT) Remote Terminal is in operation. (M) Monitor BIT 0 is in operation. This bit reflects bit 0 of Register 0.
- #2 Current Command Block Register (BC,M)/Remote Terminal Descriptor Space Address Register (RT) (BC) This register contains the address of the head pointer of the Command Block being executed. Accessing a new Command Block updates it.
- (RT) The host CPU initializes this register to indicate the starting location of the RT Descriptor Space. The host must allocate 320 sequential locations following this starting address. For proper operation, this location must start on an I x 512 decimal address boundary, where I is an integer multiple. (I = 0 is valid boundary conditions.)
- (M) This register contains the address of the control/status word of the current Monitor Command Block. Accessing a new Command Block updates it.

#3 Polling Compare Register

In the polling mode, the CPU sets the Polling Compare Register to indicate the RT response word on which the BCRTM should interrupt. This register is 11 bits wide, corresponding to bit times 9 through 19 of the RT's 1553 status word response. The sync, Remote Terminal Address, and parity bits are not included (see the section on Polling, page 32).

#4 BIT (Built-In-Test) Word Register

The BCRTM uses the contents of this register when it responds to the Transmit BIT Word Mode Code (#10011). In addition, the BCRTM writes to the two most significant bits of the BIT Word Register in response to either an Initiate Self-Test Mode Code (RT mode) or a write to Register 11 (BIT Start Command). If the BIT Word needs to be modified, it can be read out, modified, then rewritten to this register. Note that if the processor writes a "1" to either bit 14 or 15 of this register, it effectively induces a BIT failure.

Bit Description

Number

BIT 15 Channel B failure.

BIT 14 Channel A failure.

BITs 13-0 BIT Word. The least significant fourteen bits of the BIT Word are user programmable.

#5 Current Command Register (Read Only)

In the RT or Monitor mode, this register contains the command currently being processed. When not processing a command, the BCRTM stores the last command or status word transmitted on the 1553B bus. This register is updated only when bit 0 of Register 0 is set. In the BC mode, this register contains the most current command sent out on the 1553B bus.

#6 Interrupt Log List Pointer Register

Initialized by the CPU, the Interrupt Log List Pointer Register indicates the start of the Interrupt Log List. After each list entry, the BCRTM updates this register with the address of the next entry in the list. (See page 37.)

#7 High-Priority Interrupt Enable Register (Read/Write)

Setting the bits in this register causes a High-Priority Interrupt when the enabled event occurs. To service the High-Priority Interrupt, the user reads Register 8 to determine the cause of the interrupt, then writes to Register 8 to clear the appropriate bits. The BCRTM also provides a Standard Priority Interrupt Scheme that does not require host intervention. If High-Priority Interrupt service is not possible in a given application, it is advisable to use the Standard Priority features.

Bit Description

Number

BITs 15-9 Reserved.

- BIT 8 Data Overrun Enable. When set, this bit enables an interrupt when DMAG was not received by the BCRTM within the allotted time needed for a successful data transfer to memory.
- BIT 7 (BC) Illogical Command Error Enable. This bit enables a High-Priority Interrupt to be asserted upon the occurrence of an Illogical Command. Illogical commands include incorrectly formated RT-RT Command Blocks.
- BIT 6 (RT) Dynamic Bus Control Mode Code Interrupt Enable. When set, the BCRTM asserts an interrupt when the Dynamic Bus Control Mode Code is received.
- BIT 5 Subsystem Fail Enable. When set, a High-Priority Interrupt is asserted after receiving a Subsystem Fail (SSYSF) input pin.
- BIT 4 End of BIT Enable. This bit indicates the end of the internal BIT routine.
- BIT 3 BIT Word Fail Enable. This bit enables an interrupt indicating that the BCRTM detected a BIT failure.
- BIT 2 (BC) End of Command Block List Enable (see Command Block Control Word, page 31.) This interrupt can be superseded by other high-priority interrupts.
- BIT 1 Message Error Enable. If enabled, a High-Priority Interrupt is asserted at the occurrence of a message error. If a High-Priority Interrupt condition occurs, as the result of an enabled message error, the device will halt operation until the user clears the interrupt by writing a "1" to Bit 1 of the High-Priority Interrupt Status/Reset Register (Reg. #8). If this interrupt is not cleared, the BCRTM remains in the HALTED state (appearing to be "locked-up"), even if it receives a valid message. This High-Priority Interrupt scheme is necessary in order to maintain the BCRTM's state of operation so that the host CPU has this information available at the time of interrupt service.
- BIT 0 Standard Interrupt Enable. Setting this bit enables the STDINTL pin, but does not cause a high-priority interrupt. If low, only the STDINTP pin is asserted when a Standard Interrupt occurs.

#8 High-Priority Interrupt Status/Reset Register

When a High-Priority Interrupt is asserted, this register indicates the event that caused it. To clear the interrupt signal and reset the bit, write a "1" to the appropriate bit. See the corresponding bit definitions of Register 7, High-Priority Interrupt Enable Register.

Description Bit

Number

BITs 15-9 Reserved.

- Data Overrun. BIT 8
- Illogical Command. BIT 7
- Dynamic Bus Control Mode Code Received. BIT 6
- BIT 5 Subsystem Fail.
- End of BIT. BIT 4
- BIT Word Fail. BIT 3
- End of Command Block. BIT 2
- Message Error. BIT 1
- Standard Interrupt. The BCRTM sets this bit when any Standard Interrupt occurs, providing bit 0 of BIT 0 Register 7 is enabled. (Reset STDINTL output.)

#9 Standard Interrupt Enable Register

This register enables Standard Interrupt logging for any of the following enabled events (Standard Interrupt logging can also occur for events enabled in the BC Command Block or RT Subaddress/Mode Code Descriptor):

Description

Number

BITs 15-6 Reserved.

- (RT) Illegal Broadcast Command. When set, this bit enables an interrupt indicating that an Illegal BIT 5 Broadcast Command has been received.
- (RT) Illegal Command. When set, this bit enables an interrupt indicating that an illegal command has BIT 4 been received.
- (BC) Polling Comparison Match. This enables an interrupt indicating that a polling event has BIT 3 occurred. The user must also set bit 12 in the BC Command Block Control Word for this interrupt to occur.
- (BC) Retry Fail. This bit enables an interrupt indicating that all the programmed number of retries BIT 2 have failed.
- (BC,RT,M) Message Error Event. This bit enables a standard interrupt for message errors. BIT 1
- (BC,M) Command Block Interrupt and Continue. This bit enables an interrupt indicating that a BIT 0 Command Block, with the Interrupt and Continue Function enabled, has been executed.

#10 Remote Terminal Address Register

This register sets the Remote Terminal Address via software. The Change Lock-Out Enable feature, when set, prevents the Remote Terminal Address or the BCRTM Mode Selection from changing.

Bit Description

- BIT 15 (RT) Instrumentation. Setting this bit sets the RT status word Instrumentation bit.
- BIT 14 (RT) Busy. Setting this bit sets the RT status word Busy bit. It does not inhibit data transfers to the subsystem.
- BIT 13 (RT) Subsystem Fail. Setting this bit sets the RT status word Subsystem Flag bit. In the RT mode, the Subsystem Fail is also logged into the Message Status Word.
- BIT 12 (RT) Dynamic Bus Control Acceptance. Setting this bit sets the RT status word Dynamic Bus Control Acceptance bit when the BCRTM receives the Dynamic Bus Control Mode Code from the currently active Bus Controller. Host intervention is required for the BCRTM to take over as the active Bus Controller.
- BIT 11 (RT) Terminal Flag. Setting this bit sets the RT status word Terminal Flag bit; the Terminal Flag bit in the RT status word is also internally set if the BIT fails.
- BIT 10 (RT) Service Request. Setting this bit sets the RT status word Service Request bit.
- BIT 9 (RT) Busy Mode Enable. Setting this bit sets the RT status word Busy bit and inhibits all data transfers to the subsystem.
- BIT 8 BC/RT Mode Select. This bit's state reflects the external pin BCRTSEL. It does not necessarily reflect the state of the chip, since the BC/RT Mode Select is software-programmable via bit 10 of Register 0. This bit is read-only.
- BIT 7 Change Lock-Out. This bit's state reflects the external pin LOCK. When set, this bit indicates that changes to the RT address or the BC/RT Mode Select are not allowed using internal registers. This bit is read-only.
- BIT 6 Remote Terminal Address Parity Error. This bit indicates a Remote Terminal Address Parity error. It appears after the Remote Terminal Address is latched if a parity error exists.
- BIT 5 Remote Terminal Address Parity. This is an odd parity input bit used with the Remote Terminal Address. It ensures accurate recognition of the Remote Terminal Address.
- BITs 4-0 Remote Terminal Address (Bit 0 is the LSB). This reflects the RTA4-0 inputs at Master Reset.

 Modify the Remote Terminal Address by writing to these bits.

#11 BIT Start Register (Write Only)

Any write (i.e., data = don't care) to this register's address location initiates the internal BIT routine, which lasts 100 µs. Verify using the BIT-in-Progress bit in the Status Register. A programmed reset (write to Register 12) must precede a write to this register to initiate the internal BIT. A failure of the BIT will be indicated in Register 4 and the BCRTF pin.

The BCRTM's self-test performs an internal wrap around test between its Manchester encoder and its two Manchester decoders. If the BCRTM detects a failure on either the primary or the secondary channel, it flags this failure by setting bit 14 of Register 4 (BIT Word Register) for Channel A and/or bit 15 for Channel B. When in the Remote Terminal mode, while the BCRTM is performing its self-test, it ignores any commands on the 1553 bus until it has completed the self-test.

#12 Programmed Reset Register (Write Only)

Any write (i.e., data = don't care) to this register's address location initiates a reset sequence of the encoder/decoder and protocol sections of the BCRTM which lasts less than 1 microsecond. This is identical to the reset used for the Reset Remote Terminal Mode Code except that command processing halts. For a total reset (i.e., including registers), see the MRST signal description.

#13 RT Timer Reset Register (Write Only)

Any write (i.e., data = don't care) to this register's address location resets the RT Time Tag timer to zero.

The BCRTM's Remote Terminal Timer time-tags message transactions. The time tag is generated from a free-running eight-bit timer of 64 microseconds resolution. This timer can be reset to zero simply by writing to Register 13. When the timer is reset, it immediately starts running.

#14 Bus Monitor Control Register

- Bus Monitor Select. This bit should be cleared for RT mode operation. The host sets this bit to **BIT 15** enable the BCRTM's Monitor mode of operation. Bit 10 of Register 0 must also be "0" to enable the
- Monitor All Terminals. When this bit is set, the BCRTM monitors all remote terminal bus activity. If BIT 14 this bit is not set, then bit 13 must be set. This bit should be cleared for RT mode operation.
- Monitor Declared Terminals. When this bit is set, the BCRTM monitors only the remote terminals **BIT 13** that are selected in Registers 16 and 17. If this bit is not set, then bit 14 must be set. This bit should be cleared for RT Mode operation.

BITs 12-0 Reserved.

#15 Reserved Register

This register is reserved for BCRTM use only and the host should not access it.

#16 Monitor Selected Remote Terminal Addresses 15 - 0

BITs 15-0 Monitor Selected Remote Terminal Addresses 15-0. By setting the appropriate bit in this register, the host can determine which of the remote terminals, from RT 0 through RT 15, the BCRTM will monitor. For example, by setting bit 5 in this register, the host instructs the BCRTM to only monitor the bus activity for remote terminal number 5. These bits are not mutually exclusive, therefore, the host can monitor any number of different remote terminals by selecting the proper combination of

#17 Monitor Selected Remote Terminal Addresses 31 - 16

BITs 15-0 Monitor Selected Remote Terminal Addresses 31-16. By setting the appropriate bit in this register, the host can determine which of the remote terminals, from RT 16 through RT 31, the BCRTM will monitor. For example, by setting bit 5 in this register, the host instructs the BCRTM to only monitor the bus activity for remote terminal number 21. These bits are not mutually exclusive, therefore, the host can monitor any number of different remote terminals by selecting the proper combination of bits on this register and Register 16.

	15	14	13	12	11	10	9	8		
	UNUSED	RT31	SA31	всто	EXTOVR	BC/RT	RTYALTB	BUSBEN		
	7 CHNSEL	6	5	4	3	2	1	0		
	BUSAEN	RTY	CNT	RTYBCME	RTYTO	RTYME	RTYBSY	STEN		
‡1	BC/RT STA	TUS REGIS	TER 13	12	11	10	9	8		
	TEST	RTACT	DYNBUS	RT FLAG	SRQ	BUSY	BIT	RESET		
	7	6	5	4	3	2	1	0		
	BC/RT	BUSA/B	SSFAIL	UNUSED	UNUSED	UNUSED	UNUSED	CMBKPG		
7 2	(BC) CURRENT COMMAND BLOCK REGISTER (RT) REMOTE TERMINAL DESCRIPTOR SPACE ADDRESS REGISTER									
	15	14	13	12	11	10	9	8		
	A15	A14	A13	A12	A11	A10	A9	A8		
	7	6	5	4	3	2	1	0		
	A7	A6	A5	A4	A3	A2	A1	A0		
13	POLLING C	OMPARE R	EGISTER 13	12	11	10	9	8		
	X	×	×	×	×	MSGERR	INSTR	SRQ		
	7	6	5	4	3	2	1	0		
	SWBT12	SWBT13	SWBT14	BRDCST	BUSY	SS FLAG	DBC	TF		
4	BIT WORD	REGISTER 14	13	12	11	10	9	8		
	CHBFAIL	CHAFAIL	D13	D12	D11	D10	D9			
	7	6	5	4	3	2		D8		
	D7	D6	D5	D4			1	0		
		1 00		D4	D3	D2	D1	D0		
5	CURRENT C	OMMAND 14	REGISTER 13	12	11	10	9	8		
	D15	D14	D13	D12	D11	D10	D9	D8		
	7	6	5	4	3	2	1	0		
	D7	D6	D5	D4	D3	D2	D1	D0		
6	INTERRUPT	LOG LIST I	POINTER RE	GISTER 12	44	40				
	A15				11	10	9	. 8		
	7	A14 6	A13 5	A12	A11	A10	A9	A8		
	A7	A6	A5	A4	3	2	1	0		
			_ A3	A4	А3	A2	A1	A0		
7	BCRTM HIG	H-PRIORIT	Y INTERRUP	T ENABLE	REGISTER	10	9	8		
	UNUSED	UNUSED	UNUSED	UNUSED	UNUSED	UNUSED	UNUSED	DATOVR		
	7	6	5	4	3	2	1	0		
	ILLCMD	DYNBUS	SSFAIL	ENDBIT	BITFAIL	EOL	MSGERR	STDINT		
						l				
#8	BCRTM HIG	H-PRIORIT	Y INTERRUP 13	T STATUS/F	RESET REGI:	STER 10	9	8		
18							9 UNUSED	8 DATOVR		
8	15	14	13	12	11	10	·			

Table 1. BCRTM Registers (continued)

	15	14	13	12	11	10 UNUSED	9 UNUSED	8 UNUSED
	UNUSED	UNUSED	UNUSED	UNUSED	UNUSED	2	1	0
	77	6	5	4	3		MSGERR	CMDBLK
	UNUSED	UNUSED	ILLBCMD	ILLCMD	POLMTCH	RTYFAIL	WISGERN	CINIDBER
	REMOTE TE	DRAINIAI AI	NODESS RE	GISTER				
10	15	14	13	12	11	10	9	88
	INSTR	BUSY2	SS FLAG	DBC	RT FLAG	SRQ	BUSY1	BC/RT
	7	6	5	4	3	2	11	0
	LOCK	PARERR	RTAPAR	RTA4	RTA3	RTA2	RTA1	RTA0
11	BUILT-IN-TE		REGISTER 13	12	11	10	9	8
	15	14		X	1 x	T x	T x	×
	<u>×</u>	X	5 ×	4	3	2	1	0
	7		ı ×	×	T x	T x	×	×
	X	X	L^_		· · ·	1	_ 	
ŧ12	PROGRAMI	MED RESET	REGISTER				•	8
	15	14	13	12	11	10	9	T ×
	X	×	x	X	X	x	<u> </u>	1^
	7	6	5	4	3	2	1 1	T x 1
	×	X	X	X	×	X	x	<u> </u>
¥13			IMER RESE	40	`	4.0	9	•
	15	14 T v	13 X	12 T x	11 x	10 X	X	8 X
	Х	х	х	X 4	X 3			
	X 7	X 6		Х	Х	X	х	X
#14	X 7 X BUS MONI	X 6 X TOR CONT	X 5	X 4 X	X 3	X 2	1 1	X 0
#14	7 X	6 X	X 5 X ROL REGIST	X 4 X	X 3 X	X 2 X	1 X	0 X
#14	X 7 X BUS MONI	X 6 X TOR CONT	X 5 X ROL REGIST	X 4 X X ER 12	3 X	X 2 X	1 X	0 X
#14	X 7 X BUS MONI 15 BMS 7	X 6 X TOR CONTI	X 5 X ROL REGIST 13 MDT	X 4 X X ER 12 X	3 X X X X X X X X X X X X X X X X X X X	2 X X 10 X	1 X 9 X	0 X B
	X 7 X BUS MONII 15 BMS 7 X RESERVED	X 6 X TOR CONTI 14 14 MAT 6 X D REGISTER	X 5 X ROL REGIST 13 MDT 5 X (Host CPU	X 4 X ER 12 X 4 X Should No	X 3 X 11 X 3 X t Access)	X	X	X
#15	X 7 X BUS MONII 15 BMS 7 X RESERVED	X 6 X TOR CONTI 14 MAT 6 X REGISTER SELECTED	X 5 X ROL REGIST 13 MDT 5 X (Host CPU	X 4 X ER 12 X 4 X Should No	X 3 X 11 X 3 X t Access)	X 2 X	X	X
#14 #15 #16	X 7 X BUS MONII 15 BMS 7 X RESERVED MONITOR	X 6 X TOR CONTI 14 MAT 6 X D REGISTER SELECTED 14	X 5 X ROL REGIST 13 MDT 5 X (Host CPU REMOTE TI 13	X 4 X EER 12 X 4 X Should No	X 3 X 11 X 3 X 12 X ACCESS ADDRESSES 11	X 2 X	X	X 0
#15	X 7 X BUS MONII 15 BMS 7 X RESERVED MONITOR 15 TA15	X 6 X TOR CONTI 14 MAT 6 X D REGISTER SELECTED 14 TA14	X 5 X ROL REGIST 13 MDT 5 X (Host CPU REMOTE TI 13 TA13	X 4 X EER 12 X 4 X Should No ERMINAL A 12 TA12	X 3 X 11 X 3 X 12 X Access) ADDRESSES 11 TA11	X 2 X	X	X
#15	X 7 X BUS MONII 15 BMS 7 X RESERVED MONITOR 15 TA15 7	X 6 X TOR CONTI 14 MAT 6 X D REGISTER SELECTED 14 TA14 6	X 5 X ROL REGIST 13 MDT 5 X (Host CPU REMOTE TI 13 TA13 5	X 4 X EER 12 X 4 X Should No ERMINAL A 12 TA12 4	X 3 X	X 2 X	X	X
#15	X 7 X BUS MONII 15 BMS 7 X RESERVED MONITOR 15 TA15	X 6 X TOR CONTI 14 MAT 6 X D REGISTER SELECTED 14 TA14	X 5 X ROL REGIST 13 MDT 5 X (Host CPU REMOTE TI 13 TA13	X 4 X EER 12 X 4 X Should No ERMINAL A 12 TA12	X 3 X 11 X 3 X 12 X Access) ADDRESSES 11 TA11	X 2 X	X	X
#15	X 7 X BUS MONI 15 BMS 7 X RESERVED MONITOR 15 TA15 7 TA7	X 6 X TOR CONTI 14 MAT 6 X D REGISTER 14 TA14 6 TA6	X 5 X ROL REGIST 13 MDT 5 X (Host CPU REMOTE TI 13 TA13 5 TA5	X 4 X EER 12 X 4 X Should No ERMINAL A 12 TA12 4 TA4	X 3 X	X 2 X 10 X 2 X X 2 X X 10 TA10 TA2 TA2	X	X
#15 #16	X 7 X BUS MONI 15 BMS 7 X RESERVED MONITOR 15 TA15 7 TA7 MONITOR 15	X 6 X TOR CONTI 14 MAT 6 X D REGISTER SELECTED 14 TA14 6 TA6 SELECTED	X 5 X ROL REGIST 13 MDT 5 X (Host CPU REMOTE TI 13 TA13 5 TA5 REMOTE T	X 4 X EER 12 X 4 X Should No ERMINAL A 12 TA12 4 TA4 ERMINAL A	X 3 X 11 X 3 X 4 Access) ADDRESSES 11 TA11 3 TA3 ADDRESSES	X 2 X 10 X 2 X X 2 X X 10 TA10 TA2 TA2 16 - 31	X	X
#15 #16	X 7 X BUS MONI 15 BMS 7 X RESERVED MONITOR 15 TA15 7 TA7 MONITOR 15 TA31	X 6 X TOR CONTIL 14 MAT 6 X D REGISTER SELECTED 14 TA14 6 TA6 SELECTED 14	X 5 X ROL REGIST 13 MDT 5 X (Host CPU REMOTE TI 13 TA13 5 TA5 REMOTE TI 13	X 4 X EER 12 X 4 X Should No ERMINAL A 12 TA12 4 TA4 ERMINAL A	X 3 X 11 X 3 X ACCESS 11 TA11 3 TA3 ADDRESSES 11	X 2 X 10 X 2 X X 2 X X 10 TA10 TA10 2 TA2	X	X
#15 #16	X 7 X BUS MONI 15 BMS 7 X RESERVED MONITOR 15 TA15 7 TA7 MONITOR 15	X 6 X TOR CONTI 14 MAT 6 X CONTI 14 MAT 6 TA14 6 TA6 TA6 SELECTED 14 TA30 TA30	X 5 X ROL REGIST 13 MDT 5 X (Host CPU REMOTE TI 13 TA13 TA5 REMOTE TI 13 TA5	X 4 X EER 12 X 4 X Should No ERMINAL A 12 TA4 ERMINAL A 12 TA4	X 3 X 11 X 3 X 4 Access) ADDRESSES 11 TA11 3 TA3 ADDRESSES 11 TA27	X 2 X 10 X 2 X X 2 X X X 10 TA10 TA2 TA2 TA2 TA26 TA26	X	X

Table 1. BCRTM Registers (continued from page 18)

4.0 System Overview

The BCRTM can be configured for a variety of processor and memory environments. The host processor and the BCRTM communicate via a flexible, programmable interrupt structure, internal registers, and a user-definable shared memory area. The shared memory area (up to 64K) is completely user-programmable and communicates BCRTM control information -- message data, and status/error information.

Built-in memory management functions designed specifically for MIL-STD-1553B applications aid processor off-loading. The host needs only to establish the parameters within memory so the BCRTM can access this information as required. For example, in the RT mode, the BCRTM can store data associated with individual subaddresses anywhere within its 64K address space. The BCRTM then can automatically buffer up to 128 incoming messages of the same subaddress, thus preventing the previous messages from being overwritten by subsequent messages. This buffering also extends the intervals required by the host processor to service the data. Selecting an appropriate MCLK frequency to meet system memory access time requirements controls the memory access rate. The completion of a user-defined task or the occurrence of a user-selected event is indicated by using the extensive set of interrupts provided.

In the BC mode, the BCRTM can process multiple messages, assist in scheduling message lists, and provide host-programmable functions such as auto retry. The BCRTM is incorporated in systems with a variety of interrupt latencies by using the Interrupt

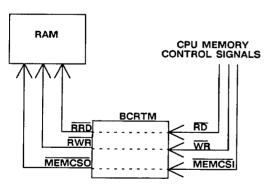


Figure 3a. Pseudo-Dual-Port RAM Control Signals

History List feature (see Exception Handling and Interrupt Logging, page 37). The Interrupt History List sequentially stores the events that caused the interrupt in memory without losing information if a host processor does not respond immediately to an interrupt.

In the Monitor (M) mode, the BCRTM's powerful linked list command block structure allows it to process a series of monitored 1553 messages without the intervention of the host. The BCRTM can store as much bus traffic as can be contained in its 64K memory space. In addition, the host has the capability of instructing the BCRTM to monitor and store data for only selected remote terminals. The host system is responsible for initializing an area in memory that tells the BCRTM where to store command word information and data for each command that the BCRTM receives on the 1553 bus. This area of memory consists of "Bus Monitor Command Blocks." An M Command Block is very similar to the BC Command Block in the BCRTM. The only real differences are the direction of information flow, and that there is no Head Pointer in the M Command Block.

5.0 System Interface

5.1 DMA Transfers

The BCRTM initiates DMA transfers whenever it executes command blocks (BC mode) or services commands (RT mode). \overline{DMAR} initiates the transfer and is terminated by the inactive edge of \overline{DMACK} . The Address Enable (AEN) input enables the BCRTM to output an address onto the Address bus.

The BCRTM requests transfer cycles by asserting the DMAR output, and initiates them when a DMAG input is received. A DMACK output indicates that the BCRTM has control of the Data and Address buses. The TSCTL output is asserted when the BCRTM is actually asserting the Address and Data buses.

To support using multiple bus masters in a system, the BCRTM outputs the \overline{DMAGO} signal that results from the \overline{DMAG} signal passing through the chip when a BCRTM bus request was not generated (\overline{DMAR}) inactive). You can use \overline{DMAGO} in daisy-chained multimaster systems.

5.2 Hardware Interface

The BCRTM provides a simple subsystem interface and facilitates DMA arbitration. The user can configure the BCRTM to operate in a variety of memory-processor environments including pseudo-dual-port RAM and standard DMA configurations.

For complete circuit description, such as arbitration logic and I/O, please refer to the appropriate application note.

5.3 CPU Interconnection

Pseudo-Dual-Port RAM Configuration
The BCRTM's Address and Data buses connect directly to RAM, with buffers isolating the BCRTM's buses from those of the host CPU (figures 3a and 3b). The CPU's memory control signals (RD, WR, and MEMCSI) pass through the BCRTM and connect to memory as RRD, RWR, and MEMCSO.

Standard DMA Configuration

The BCRTM's and CPU's data, address, and control signals are connected to each other as shown in figures 3c and 3d. The RWR RRD, and MEMCSO are activated after DMAG is asserted.

In either case, the BCRTM's Address and Data busses remain in a high-impedance state unless the $\overline{\text{CS}}$ and $\overline{\text{RD}}$ signals are active, indicating a host register access; or $\overline{\text{TSCTL}}$ is asserted, indicating a memory access by the BCRTM. CPU attempts to access BCRTM registers are ignored during BCRTM memory access. Inhibit DMA transfers by using the Busy function in the Remote Terminal Address Register while operating in the Remote Terminal mode.

The designer can use TSCTL to indicate when the BCRTM is accessing memory or when the CPU can access memory. AEN is also available (use is optional), giving the CPU control over the BCRTM's Address bus. A DMA Burst (BURST) signal indicates multiple DMA accesses.

Register Access

Registers 0 through 13 are accessed with the decode of the four LSBs of the Address bus (A0-A3) and asserting \overline{CS} . Pulse either \overline{RD} or \overline{WR} for multiple register accesses.

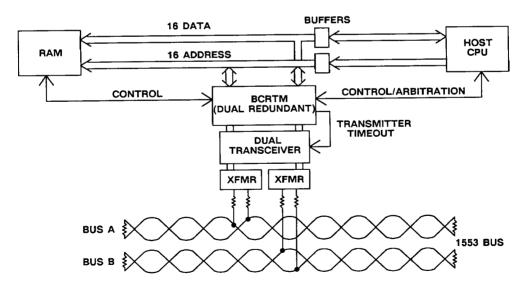


Figure 3b. CPU/BCRTM Interface -- Pseudo-Dual-Post RAM Configuration

BCRTM-21

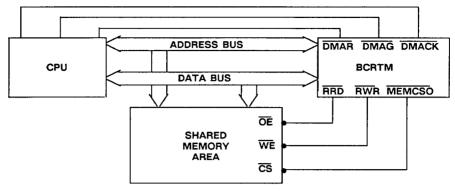


Figure 3c. DMA Signals

5.4 RAM Interface

The BCRTM's \overline{RRD} , \overline{RWR} , and \overline{MEMCSO} signals serve as read and write controls during BCRTM memory accesses. The host subsystem signals \overline{RD} , \overline{WR} , and \overline{MEMCSI} propagate through the BCRTM to become \overline{RRD} , \overline{RWR} , and \overline{MEMCSO} outputs to support a pseudo-dual-port. During BCRTM-RAM

data transfers, the host subsystem's memory signals are ignored until the BCRTM access is complete.

5.5 Transmitter/Receiver Interface

The BCRTM's Manchester II encoder/decoder interfaces directly with the 1553 bus transceiver, using the TAO-TAZ and RAZ-RAO signals for Channel A, and TBO-TBZ and RBZ-RBO signals for Channel B.

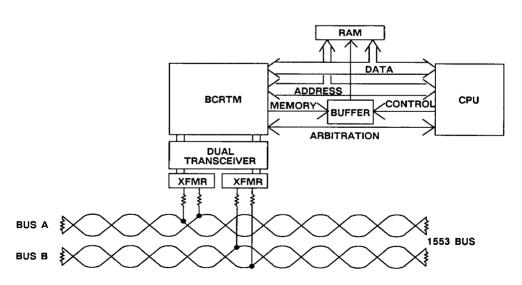


Figure 3d. CPU/BCRTM Interface -- DMA Configuration

BCRTM-22

The BCRTM also provides a $\overline{\text{TIMERON}}$ signal output and an active channel output indicator (CHA/ \overline{B}) to assist in meeting the MIL-STD-1553B fail-safe timer requirements (see figure 4).

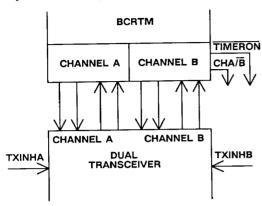


Figure 4. Dual-Channel Transceiver

6.0 REMOTE TERMINAL ARCHITECTURE

The Remote Terminal architecture is a descriptor-based configuration of relevant parameters. It is composed of an RT Descriptor Space (see figure 5) and internal, host-programmable registers. The Descriptor Space contains only descriptors. Descriptors contain programmable subaddress parameters relating to handling message transfers. Each descriptor consists of four words: (1) a Control Word, (2) a Message Status List Pointer, (3) a Data List Pointer, and (4) an unused fourth word (see figure 6.) These words indicate how to perform the data transfers associated with the designated subaddress.

A receive descriptor and a transmit descriptor are associated with each subaddress. The descriptors reside in memory and are listed sequentially by subaddress. By using the index within the descriptor, the BCRTM can buffer incoming and outgoing messages, which reduces host CPU overhead. This message buffering also reduces the risk of incoming messages being overwritten by subsequent incoming messages.

Each descriptor contains a programmable interrupt structure for subsystem notification of user-selected message transfers and indicates when the message buffers are full. Illegalizing subaddresses, in normal and broadcast modes, is accomplished by using programmable bits within the descriptor (see the RT Functional Operation section on next page).

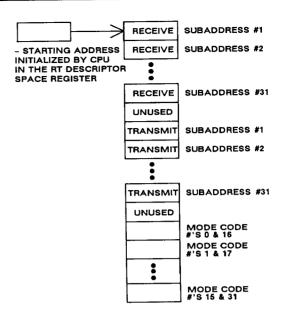


Figure 5. Descriptor Space

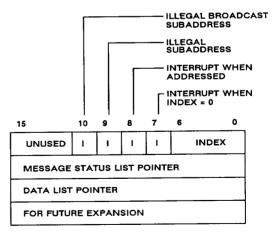


Figure 6. Remote Terminal Subaddress Descriptor

Message Status information -- including word count, an internally generated time tag, and broadcast and message validity information -- is provided for each message. The Message Status Words are stored in a separate Message Status Word list according to subaddress. The list's starting locations are programmable within the descriptor.

Message data, received or transmitted, is also stored in lists. The message capacity of the lists and the lists' locations are user selectable within the descriptor.

6.1 RT Functional Operation

The RT off-loads the host computer of all routine data transfers involved with message transfers over the 1553B bus by providing a wide range of user-programmable functions. These functions make the BCRTM's operation flexible for a variety of applications. The following paragraphs give each function's operational descriptions.

6.1.1 RT Subaddress Descriptor Definition
The host sets words within the descriptor (see figure 6, page 23). The BCRTM then reads the descriptor words when servicing a command corresponding to the specified descriptor. All bit-selectable functions are active high and inhibited when low.

A. Control Word. The first word in the descriptor, the Control Word, selects or disables message transfers and selects an index.

Bit Description Number

- 15-11 Reserved.
- BIT 10 Illegal Broadcast Subaddress. Indicates to the BCRTM not to access this subaddress using broadcast commands. The Message Error bit in the status word is set if the illegal broadcast subaddress is addressed. Since transmit commands do not apply to broadcast, this bit applies only to receive commands.
- BIT 9 Illegal Subaddress. Set by the host CPU, it indicates to the BCRTM that a command with this subaddress is illegal. If a command uses an illegal subaddress the Message Error bit in the 1553 status word is set. The Illegal Command Interrupt is also asserted if enabled.
- BIT 8 Interrupt Upon Valid Command Received. Indicates that the BCRTM is to assert an interrupt every time a command addresses this descriptor. The interrupt occurs just prior to post-command descriptor updating.
- BIT 7 Interrupt When Index = 0. Indicates that the BCRTM initiates an interrupt when the index is decremented to zero.
- BITs 6-0 Index. These bits are for indexed message buffering. Indexing means transacting a pre-specified number of messages before notifying the host CPU. After each message transaction, the BCRTM decrements the index by one until index = 0. Note that the index is decremented for messages that contain message errors.
- B. Message Status List Pointer. The host sets the Message Status List Pointer, the second word within the descriptor, and the BCRTM uses it as a starting address for the Message Status List. It is incremented by one with each Message Status Word write. If the Control Word Index is already equal to zero, the Message Status List Pointer is not incremented and the previous Message Status Word is overwritten.

Note: A Message Status Word is also written and the pointer is incremented when the BCRTM detects a message error.

- C. Data List Pointer. The Data List Pointer is the third word within the descriptor. The BCRTM stores data in RAM beginning at the address indicated by the Data List Pointer. The Data List Pointer is updated at the end of each successful message with the next message's starting address with the following exceptions:
 - If the message is erroneous, the Data List Pointer is not updated. The next message overwrites any data corresponding to the erroneous message.
 - Upon receiving a message, if the index is already equal to zero, the Data List Pointer is not incremented and data from the previous message is overwritten.
- D. Reserved. The fourth descriptor word is reserved for future use.

6.1.2 Message Status Word

Each message the BCRTM transacts has a corresponding Message Status Word, which is pointed to by the Message Status List Pointer of the Descriptor. This word allows the host CPU to evaluate the message's validity, determine the word count, and calculate the approximate time frame in which the message was transacted (figures 7 and 8).

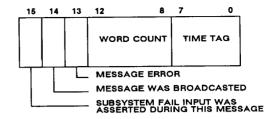


Figure 7. Message Status Word

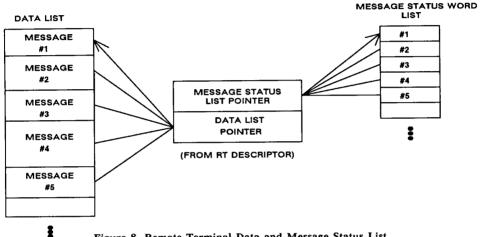


Figure 8. Remote Terminal Data and Message Status List

Message Status Word Definition

Bit Description

Number

Subsystem Failed. Indicates SSYSF was asserted before the Message Status Word transfer to memory. **BIT 15** This bit is also set when the user sets bit 13 of Register 10.

Broadcast Message. Indicates that the corresponding message was received in the broadcast mode. BIT 14

Message Error. Indicates a message is invalid due to improper synchronization, bit count, word BIT 13 count. Manchester error.

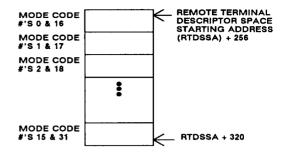
BITs 12-8 Word Count. Indicates the number of words in the message and reflects the Word Count field in the command word. Should the message contain a different number of words than the Word Count field, the Message Error flag is triggered. If there are too many words, they are withheld from RAM. If the actual word count is less than or grater than it should be, the Message Error bit in the 1553 status word is set.

BITs 7-0 Time Tag. The BCRTM writes the internally generated Time Tag to this location after message completion. The resolution is 64 microseconds. (See Register 13). If the timer reads 2, it indicates the message was completed 128 to 191 microseconds after the timer started.

BCRTM-25

6.1.3 Mode Code Descriptor Definition Mode codes are handled similarly to subaddress transactions. Both use the four-word descriptors residing in the RT descriptor space to allow the host to program their operational mode. Corresponding to each mode code is a descriptor (see figure 9a). Of the 32 address combinations for mode codes in MIL-STD-1553B, some are clearly defined functions while others are reserved for future use. Sixteen descriptors are used for mode code operations with each descriptor handling two mode codes: one mode code with an associated data word and one mode code without an associated data word. All mode codes are handled in accordance with MIL-STD-1553B. The function of the first word of the Mode Code Descriptor is similar to that of the Subaddress Descriptor and is defined below. The remaining three words serve the same purpose as in the Subaddress Descriptor.

Interrupt if Index = 0.



Note:

Mode code descriptor blocks are also provided for reserved mode codes but have no associated predefined BCRTM operation.

Figure 9a. (RT) Mode Code Descriptor Space

Control Word

BIT 7

Description Number Interrupt on Reception of Mode Code (without Data Word). **RIT 15 BIT 14** Illegalize Broadcast Mode Code (without Data Word). **BIT 13** Illegalize Mode Code (without Data Word). **BIT 12** Reserved. RIT 11 Illegalize Broadcast Mode Code (with Data Word). **BIT 10** Illegalize Transmit Mode Code (with Data Word). BIT 9 Illegalize Receive Mode Code (with Data Word). BIT 8 Interrupt on Reception of Mode Code (with Data Word).

BITs 6-0 Index. Functionally equivalent to the index described in the Subaddress Descriptor. It applies to mode codes with data words only.

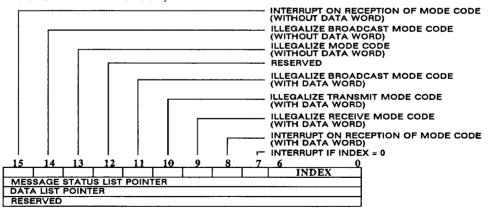


Figure 9b. (RT) Mode Code Descriptor

The descriptors, numbered sequentially from 0 to 15, correspond to mode codes 0 to 15 without data words and mode codes 16 to 31 with data words. For example, mode codes 0 and 16 correspond to descriptor 0 and mode codes 1 and 17 correspond to descriptor 1. The Mode Code Descriptor Space is appended to the Subaddress Descriptor Space starting at 0100H (256D) of the 320-word RT Descriptor Space (see figure 5).

The BCRTM autonomously supports all mode codes without data words by executing the specific function and transmitting the 1553 status word. The subsystem provides the data word for mode codes with data words (see the Data List Pointer section). For all mode codes, an interrupt can be asserted upon successful completion of the mode command by setting the appropriate bit in the control word (see figure 9b).

Dynamic Bus Control #00000

This mode code is accepted automatically if the Dynamic Bus Control Enable bit in the Remote Terminal Address Register is set. Setting the Dynamic Bus Control Acceptance bit in the 1553 status word and BCRTM Status Register confirms the mode code acceptance. A High-Priority Interrupt is also asserted if enabled. If the Dynamic Bus Control Enable bit is not set, the BCRTM does not accept Dynamic Bus Control.

Synchronize (Without Data Word) #00001
If enabled in the Mode Code #00001 Descriptor
Control Word, the BCRTM asserts an interrupt when
this mode code is received.

Transmit Status Word #00010

The BCRTM automatically transmits the 1553 status word corresponding to the last message transacted.

Initiate Self-Test #00011

The BCRTM automatically starts its BIT routine. An interrupt, if enabled, is asserted when the test is completed. The BIT Word Register and external pin BCRTF are updated when the test is completed. A failure in BIT will also set the TF status word bit.

Transmitter Shutdown #00100

The BCRTM disables the channel opposite the channel on which the command was received.

Override Transmitter Shutdown #00101
The BCRTM enables the channel previously disabled.

Inhibit Terminal Flag Bit #00110 The BCRTM inhibits the Terminal Flag from being set in the status word.

Override Inhibit Terminal Flag Bit #00111
The BCRTM disables the Terminal Flag inhibit.

Reset Remote Terminal #01000

The BCRTM automatically resets the encoder, decoders, and protocol logic.

Transmit Vector Word #10000

The BCRTM transmits the vector word from the location addressed by the Data List Pointer in the Mode Code Descriptor Block.

Synchronize (with Data Word) #10001
On receiving this mode code, the BCRTM simply stores the associated data word.

Transmit Last Command #10010

The BCRTM transmits the last command executed and the corresponding 1553 status word.

Transmit BIT Word #10011

The BCRTM transmits BIT information from the BIT Register.

Selected Transmitter Shutdown #10100 On receiving this mode code, the BCRTM simply stores the associated data word.

Override Selected Transmitter Shutdown #10101 On receiving this mode code, the BCRTM simply stores the associated data word.

Mode codes 9-15 and 22-31 are reserved for future expansion of MIL-STD-1553B.

6.2 RT Error Detection

In accordance with MIL-STD-1553B, the remote terminal handles superseding commands on the same or opposite bus. When receiving, the remote terminal performs a response time-out function of 56 microseconds for RT-RT transfers. If the response time-out condition occurs, a Message Error bit is set in the 1553 status word and in the Message Status Word. Error checking occurs on both of the Manchester logic and the word formats. Detectable errors include word count errors, long words, short words, Manchester errors (including zero crossing deviation), parity errors, and data contiguity.

6.3 RT Operational Sequence

The following is a general description of the typical behavior of the BCRTM as it processes a message in the RT mode. It is assumed that the user has already written a "1" to Register 0, bit 0, enabling RT operation.

Valid Command Received.

COMSTR goes active

 DMA Descriptor Read. After receiving a valid command, the BCRTM initiates a burst DMA:

DMA arbitration (BURST)
Control Word read
Message Status List Pointer read
Data List Pointer read

Data Transmitted/Received.

Data Word DMA.

If the BCRTM needs to transmit data from memory, it initiates a DMA cycle for each Data Word shortly before the Data Word is needed on the 1553B bus:

DMA arbitration
Data Word read (starting at Data List
Pointer address, incremented for each
successive word)

If the BCRTM receives data, it writes each Data Word to memory after the Data Word is received:

DMA arbitration
Data Word write (starting at Data List
Pointer address, incremented for each
successive word)

Status Word Transmission.

The BCRTM automatically transmits the Status Word as defined in MIL-STD-1553B. The

Message Error and Broadcast Command Received bits are generated internally. Writing to Register 10 enables the other predefined bits. For illegalized commands, the BCRTM sets the Message Error Bit in the 1553 Status Word.

Exception Handling.

If an interrupting condition occurs during the message, the following occurs:

For High-Priority Interrupts:

HPINT is asserted (if enabled in Register 7). For message errors, the BCRTM is put in a hold state until the interrupt is acknowledged (by writing a "1" to the appropriate bit in Register 8).

For Standard Interrupts:

DMA arbitration (BURST)
Interrupt Status Word write
RT Descriptor Block Pointer write
Tail Pointer read (into Register 6)
STDINTP pulses low
STDINTL asserted (if enabled)
Processing continues

Descriptor Write.

After the BCRTM processes the message, a final DMA burst occurs to update the descriptor block, if necessary:

DMA arbitration (BURST)
Message Status Word write
Data List Pointer write
(incremented by word count)
Message Status List Pointer write
(incremented by 1)
Control Word write (index decremented)

Note the following exceptions:

Mode codes without data require no descriptor update.

Predefined mode codes (18 and 19) which do not require access to memory for the data word, do not involve updating the Data List Pointer.

Messages with errors prevent updates to the Data List Pointer.

If the message index was zero, neither the Message Status List Pointer nor the Data List Pointer is updated.

7.0 Bus Controller Architecture

The BCRTM's bus controller architecture is based on a Command Block structure and internal, host-programmable registers. Each message transacted over the MIL-STD-1553B bus has an associated Command Block, which the CPU sets up (see figures 10 and 11). The Command Block contains all the relevant message and RT status information as well as programmable function bits that allow the user to select functions and interrupts. This memory interface system is flexible due to a doubly-linked list data structure.

HEAD POINTER
CONTROL WORD
COMMAND WORD 1
COMMAND WORD 2 (RT-RT ONLY)
DATA LIST POINTER
STATUS WORD 1
STATUS WORD 2 (RT-RT ONLY)
TAIL POINTER

Figure 10. Command Block

In a doubly-linked Command Block structure, pointers delimit each Command Block to the previous and successive blocks (see figure 12). The linking feature eases multiple message processing tasks and supports message scheduling because of its ability to loop through a series of transfers at a predetermined cycle time. A data pointer in the command allows efficient space allocation because data blocks only have to be configured to the exact word count used in the message. Data pointers also provide flexibility in data-bank switching.

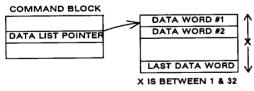


Figure 11. Data Placement

A control word with bit-programmable functions and a Message Error bit are in each Command Block. This allows selecting individual functions for each message and provides message validity information. The BCRTM's register set provides additional global parameters and address pointers.

A programmable auto retry function is selectable from the control word and Control Register.

The auto retry can be activated when any of the following occurs:

- · Busy Bit set in the status word
- Message Error (indicated by the RT status response)
- Response Time-Out
- Message Error detected by the Bus Controller

One to four retries are programmable on the same or opposite bus.

The Bus Controller also has a programmable intermessage delay timer that facilitates message transfer scheduling (see figures 13 and 14). This timer, programmed in the control word, automatically delays between the start of two successive commands.

A polling function is also provided. The Bus Controller, when programmed, compares incoming status words to a host-specified status word and generates an interrupt if the comparison indicates any matching bits. An Interrupt and Continue function facilitates the host subsystem's synchronization by generating an interrupt when the specified Command Block's message is executed.

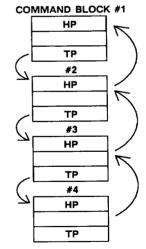


Figure 12. Command Block Chaining

7.1 BC Functional Operation

The Bus Controller off-loads the host computer of many functions needed to coordinate 1553B bus data transfers. Special architectural features provide message-by-message flexibility. In addition, a programmable interrupt scheme, programmable intermessage timing delays, and internal registers enhance the BCRTM's operation.

The host determines the first Command Block by setting the initial starting address in the current Command Block Register. Once set, the BCRTM updates the current Command Block register with the

next Command Block Address. The BCRTM then executes the sequential Command Blocks and counts out message delays (where programmed) until it encounters the last Command Block listed (indicated by the End of List bit in the control word). Interrupts are asserted when enabled events occur (see the Exception Handling and Interrupt Logging section, page 37).

The functions and their programming instructions are described below. The registers also contain many programmable functions and function parameters.

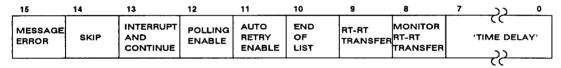


Figure 13. Control Word

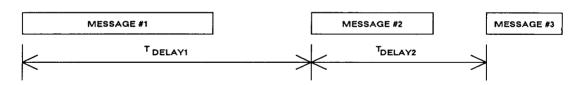


Figure 14. BC Timing Delays

BC Command Block Definition

Each Command Block contains (see figure 10):

- A. Head Pointer. Host-written, this location can contain the address of the previous Command Block's Head Pointer. The BCRTM does not access this location.
- B. Control Word. Host-written, the Control Word contains bit-selectable options and a Message Error bit the BCRTM provides (see figure 13). The bit definitions follow.

Bit Description

Number

- BIT 15 Message Error. The BCRTM sets this bit when it detects an invalid RT response as defined in MIL-STD-1553B.
- BIT 14 Skip. When set, this bit instructs the BCRTM to skip this Command Block and execute the next.
- BIT 13 Interrupt and Continue. If set, a Standard Interrupt is asserted when this block is addressed; operation, however, continues. Note that this interrupt must also be enabled by setting bit 0 of Register 9.
- BIT 12 Polling Enable. Enables the BCRTM's polling operation.
- BIT 11 Auto Retry Enable. When set, the Auto Retry function, governed by the global parameters in the Control Register, is enabled for this message.
- BIT 10 End of List. Set by the CPU, this bit indicates that the BCRTMP, upon completion of the current message, will halt and assert a High-Priority Interrupt. The interrupt must also be enabled in the High-Priority Interrupt Enable Register.
- BIT 9 RT-RT. Set by the CPU, this indicates that this Command Block transacts an RT-RT transfer.
- BIT 8 Monitor RT-RT Transfer. Set by the CPU, this function indicates that the BCRTM should receive and store the message beginning at the location indicated by the data pointer.
- BITs 7-0 Time Delay. The CPU sets this field, which causes the BCRTM to delay the specified time between sequential message starts (see figures 13 and 14). Regardless of the value in the Time Delay field (including zero), the BCRTM will at least meet the minimum 4 µs intermessage gap time as specified in MIL-STD-1553B. The timer is enabled by having a non-zero value in this bit field. When using this function, please note:
 - Timer resolution is 16 microseconds. As an example, if a given message requires 116 μs to complete (including the minimum 4 μs intermessage gap time) the value in the Time Delay field must be at least 00001000 (8 x 16 μs = 128 μs) to provide an intermessage gap greater than the 4 μs minimum requirement.
 - If the timer is enabled and the Skip bit is set, the timer provides the programmed delay before
 proceeding.
 - If the message duration exceeds the timer delay, the message is completed just as if the timer were not enabled.
 - If SKIP = 1 and EOL = 1, the HPINT is generated if enabled.
 - If SKIP = 1 and Interrupt and Continue = 1, the STDINT is generated if enabled.
- C. Command Word One. Initialized by the CPU, this location contains the first command word corresponding to the Command Block's message transfer.
- D. Command Word Two. Initialized by the CPU, this location is for the second (transmit) command word in RT-RT transfers. In messages involving only one RT, the location is unused.
- E. Data Pointer. Initialized by the CPU, this location contains the starting location in RAM for the command block's message (see figure 15).
- F. Status Word One. Stored by the BCRTM, this location contains the entire Remote Terminal status response.
- G. Status Word Two. Stored by the BCRTM, this location contains the receiving Remote Terminal status word. For transfers involving one Remote Terminal, the location is unused.
- H. Tail Pointer. Initialized by the host CPU, the Tail Pointer contains the next Command Block's starting address.

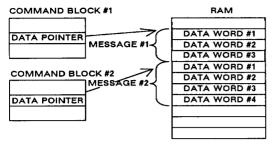


Figure 15. Contiguous Data Storage

7.2 Polling

During a typical polling scenario (see figure 16) the Bus Controller interrogates remote terminals by requesting them to transmit their status words. This feature can also alert the host if a bit is set in any RT status word response during normal message transactions. The BCRTM enables the host to initialize a chain of Command Blocks with the command word's Polling Enable bit. A programmable Polling Compare Register (PCR) is provided. In the polling mode, the Remote Terminal response is compared to the Polling Compare Register contents. Program the PCR by setting the PCR bits corresponding to the RT's 1553 status word bits to be compared. If they match (i.e., two 1's in the same bit position) then, if enabled in both the BC Command Block Control Word and in the Standard Interrupt Enable Register (Register 9), a polling comparison interrupt is generated.

Example 1. No bit match is present

PCR 00000000001

RT's 1553 status word response 00000100010

Result No Polling Comparison Interrupt

Example 2. Bit match is present

PCR 00100100000 RT's 1553 status word response 00000100000 Result Polling Comparison Interrupt

7.3 BC Error Detection

The Bus Controller checks for errors (see the Exception Handling and Interrupt Logging and the RT Error Detection sections, pages 37 and 28) on each message transaction. In addition, the BC compares the RT command word addresses to the incoming status word addresses. The BC monitors for response time-out and checks data and control words for proper format according to MIL-STD-1553B. Illogical commands include incorrectly formatted RT-RT Command Blocks.

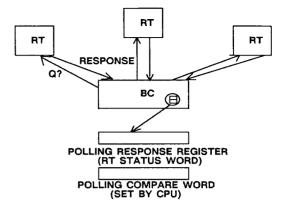


Figure 16. Polling Operation

7.4 Bus Controller Operational Sequence

The following is a general description of the typical behavior of the BCRTM as it processes a message in the BC mode.

The user starts BC operation by writing a "1" to Register 0, Bit 0.

 Command Block DMA - the following occurs immediately after Bus Controller startup:

> DMA arbitration (BURST) Control Word read Command Word 1 read (from third location of Command Block) Data List Pointer read

A. For BC-to-RT Command Blocks:

The BCRTM transmits the Command Word.

• Data Word DMA

DMA arbitration Data Word read (starting at Data List Pointer address, incremented for each successive word)

The BCRTM transmits the Data Word. Data Word DMAs and transmissions continue until all Data Words are transmitted.

• Status Word DMA

The BCRTM receives the RT Status Word.

DMA arbitration Status Word write (to sixth location of Command Block) B. For RT-to-BC Command Blocks:

The BCRTM transmits the Command Word.

Status Word DMA

The BCRTM receives the RT Status Word.

DMA arbitration Status Word write (to sixth location of Command Block)

The BCRTM receives the first Data Word.

Data Word DMA

DMA arbitration Data Word write (starting at Data List Pointer address, incremented for each successive word)

Data Word receptions and DMAs continue until all Data Words are received.

C. For RT(B)-to-RT(A) Command Blocks:

The BCRTM transmits Command Word 1 to RT(B).

• Command Word 2 DMA

DMA arbitration Command Word 2 read (from fourth location of Command Block)

The BCRTM transmits Command Word 2 to RT(A).

The BCRTM receives the RT Status Word from RT(A).

Status Word DMA for RT(A) Status Word

DMA arbitration Status Word write (to sixth location of Command Block)

The BCRTM receives the first Data Word

 Data Word DMA (only if the BCRTM is enabled to monitor the RT-to-RT message).

DMA arbitration
Data Word write (starting at Data List
Pointer address, incremented for each
successive word)

Data Word receptions and DMAs continue until all Data Words are received.

The BCRTM receives the RT Status Word from RT(B).

• Status Word DMA for RT(B) Status Word

DMA arbitration Status Word write (to seventh location of Command Block) Exception Handling.

If an interrupting condition occurs during the message, the following occurs:

For High-Priority Interrupts:

HPINT is asserted (if enabled in Register 7). For message errors, the BCRTM is put in a hold state until the interrupt is acknowledged (by writing a "1" to the appropriate bit in Register 8).

For Standard Interrupts:

DMA arbitration (BURST)
Interrupt Status Word write
Command Block Pointer write
Tail Pointer read (into Register 6)
STDINTP pulses low
STDINTL asserted (if enabled)
Processing continues

If Retries are enabled and a Retry condition occurs, the following DMA occurs:

DMA arbitration (BURST)
Control Word read
Command Word 1 read (from third
location of Command Block)
Data List Pointer read

The BCRTM proceeds from the current Command Block to the next successive Command Block.

 If no Message Error has occurred during the current Command Block, the following occurs:

DMA arbitration (BURST)
Command Block Tail Pointer read (to determine location of next Command Block. Note that this occurs only if no Retry.)
DMA hold cycle
Control Word read (next Command Block)
Command Word 1 read (next Command Block)
Data List Pointer read

 If the BCRTM detects a Message Error while processing the current Command Block, the following occurs:

DMA arbitration (BURST)
Control Word write
Command Block Tail Pointer read (to
determine location of next Command
Block. Note that this occurs only if no
Retry.)
DMA hold cycle
Control Word read (next Command
Block)
Command Word 1 read (next Command
Block)
Data List Pointer read

The BCRTM proceeds again from point A, B, or C as shown above.

7.5 BC Operational Example (see figure 22 on page 39)

The BCRTM is programmed initially to accomplish the following:

The first Command Block is for a four-word RT-RT transfer with the BCRTM monitoring the transfer and storing the data.

- Auto-retry is enabled on the opposite bus using only one retry attempt, if the incoming Status Word is received with the Message Error bit set.
- Wait for a time delay of 400 microseconds before proceeding to the next Command Block.
- The Data List Pointer contains the address

The second Command Block is for a BC-RT transfer of two words.

- The End of List bit is set in its Control Word.
- The Data List Pointer contains the address 0404H.
- The Polling Enable bit is set and the Polling Compare Register contains a one in the Subsystem Fail position (Bit 2).

Then:

- A. The CPU initializes all the appropriate registers and Command Blocks, and issues a Start Enable by writing a "1" to Register 0, Bit 0.
- B. The BCRTM, through executing a DMA cycle, reads the Control Word, Command Words, and the Data List Pointer. The delay timer starts and message execution begins by transmitting the receive and transmit commands stored in the Command Blocks. The BCRTM then waits to receive the Status Word back from the transmitting RT.
- C. The BCRTM receives the RT Status Word with all status bits low from the transmitting RT and stores the Status Word in Command Block 1. The incoming data words from the transmitting RT follow. The BCRTM stores them in memory locations 0400H - 0403H.

If the Status Word indicates that the message cannot be transmitted (Message Error), the response time-out clock counts to zero and the allotted message time runs out. An auto-retry can

- be initiated if programmed to do so. Nevertheless, the ME bit in the Control Word is set.
- D. The BCRTM receives the Status Word response from the receiving RT. The ME bit in the Status Word is set, indicating the message is invalid. The BCRTM initiates the auto retry function, (as programmed) on the alternate bus, re-transmits the Command Words, receives the correct Status Word, and stores the data again in locations 0400H 0403H. This time the Status Word response from the receiving RT indicates the message transfer is successful.
- E. The timer delay between the two successive transactions counts down another 135 microseconds before proceeding. This is determined as follows:

The message transaction time is approximately 130 microseconds (the only approximation is due to the range in status response and intermessage gap times specified by MIL-STD-1553B). Approximating that with the retry, the total duration for the two attempts is 265 microseconds.

- F. The BCRTM reads the Tail Pointer of Command Block 1 and places it in the Current Command Register. It also reads the Control Word, Command Word, and Data List Pointer, and the first data word in the second Command Block.
- G. Since this is a BC-RT transfer, the BCRTM transmits the receive command followed by two data words from locations 0404H 0405H in memory. The BCRTM reads the second data word from memory while transmitting the first.
- H. The BCRTM receives the status response from the RT. In this case, the Status Word indicates, by the ME bit being low, that the message is valid. The Status Word also has the Subsystem Fail bit set.
- I. The Status Word is stored in the Command Block. The BCRTM, having encountered the end of the list, halts message transactions and waits for another start signal.
- J. The BCRTM asserts a High-Priority Interrupt indicating the end of the command list. Due to the polling comparison match, the BCRTM also asserts a Standard Priority Interrupt and logs the event in the Interrupt Log List.

8.0 Bus Monitor Architecture

The BCRTM's bus monitor architecture is based on a Command Block structure and internal, host-programmable registers. Each message transacted over the MIL-STD-1553B bus (for a monitored RT address) has an associated Command Block, which the CPU sets up (see figures 17 and 18). The Command Block contains all the relevant message and RT status information as well as programmable function bits that allow the user to select functions and interrupts.

MC	ONITOR CONTROL/STATUS
15	53 COMMAND WORD 1
15	53 COMMAND WORD 2
DA	TA LIST POINTER
15	53 STATUS WORD 1
15	53 STATUS WORD 2
ΤA	IL POINTER

Figure 17. BCRTM Bus Monitor Command Block

In a linked list Command Block structure, pointers delimit each Command Block to the successive block (see figure 19). A data pointer in the Command Block allows efficient space allocation because data blocks do not have to be placed contiguously in memory.

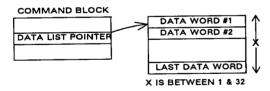


Figure 18. Data Placement

A Monitor control/status word with an eight-bit Time Tag, an Interrupt When Addressed bit, a Message Error bit, and a Command Block Activated bit are in each Monitor Command Block. The user can access these control/status words to determine which Monitor Command Block was last utilized. Also, by selecting the Interrupt When Addressed feature that is available for each Command Block, the host can

determine when particular remote terminal activity has occurred.

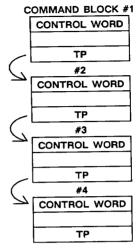


Figure 19. Monitor Command Block Tail Pointers

8.1 Monitor Functional Operation

The Bus Monitor function is a register-selectable mode of operation. The host uses Registers 14, 16, and 17 in conjunction with Register 0 to program the BCRTM to monitor any combination of remote terminals or all of the remote terminals.

The BCRTM's memory management scheme gives the host a great deal of flexibility for processing 1553 bus data and is compatible with many bus monitoring applications. The host CPU is responsible for initializing the Monitor control/status word, the Data Pointer, and the Tail Pointer in each of the Monitor Control Blocks. The Monitor structure is analogous to the Bus Controller Command Block scheme. The only real difference is the direction of information flow.

The number of Monitor Control Blocks that the host initializes depends on the data latency requirements for post-processing of 1553 commands. The linked list of Command Blocks could be connected in a loop fashion, with the BCRTM accessing the loop at one point and the host CPU processing the message behind that point. The bit positions of the BM control/status word are defined as shown in figure 20.

15	14	13	12	11	10	9	8
CBA	ME	IOA	BA/B	X	×	×	х
7	6	5	4	3	2	1	0
TT7	ТТ6	TT5	TT4	ттз	TT2	TT1	TTO

B0 - B7 Time Tag (64 μs resolution)

B8 - B11 Not used

B12 - Bus A/B:Defines which of the dual redundant 1553 buses on which the BCRTM received this message.

B13 - IWA: Interrupt When Addressed. The BCRTM issues a standard priority

interrupt when the host accesses this Bus Monitor Command Block. Message Error. The BCRTM sets this bit if a 1553 message error oc

curred while receiving this message.

B15 - CBA: Command Block Activated. The BCRTM sets this bit when this Bus

Monitor Command Block is accessed by the BCRTM.

Figure 20. Bus Monitor Control/Status Word

8.2 Monitor Error Detection

R14 - MF:

In the Monitor mode, the BCRTM checks all monitored messages for errors. Detectable errors include word count errors, long words, short words, Manchester errors (including zero crossing deviation), parity errors, and data contiguity.

Due to the nature of the 1553 protocol, it can be very difficult for any monitoring device to interpret some types of errors on the 1553 bus. For example, suppose an RT (whose RT Address the BCRTM is monitoring) incorrectly responds to a Broadcast command. The BCRTM, which is not receiving or transmitting the message, cannot distinguish between the erroneous status word and a new command sent from the Bus Controller, since the status sync is identical to the command sync. In this case, the BCRTM will put the extra status word in a new Monitor Command Block, and then report a Message Error due to the incorrect protocol on the erroneously interpreted status word the RT transmitted.

8.3 Monitor Operational Sequence

The following is a general description of the operation of the BCRTM as it processes a monitored BC-to-RT message in the Monitor mode. DMA operations will vary slightly depending on the type of message (i.e., RT-to-BC, RT-to-RT, etc.). It is assumed that the user has already written a "1" to Register 0, bit 0, and all other registers are in the appropriate state to enable Monitor operation.

Valid Command Received.

COMSTR goes active.

 DMA Monitor Command Block Read. After receiving a valid command, the BCRTM initiates a burst DMA: DMA arbitration (BURST) Control Word read Command Word write Data List Pointer read

Data Received.

• Data Word DMA

The BCRTM initiates a DMA cycle for each Data Word to store the data in memory, whether the command was a transmit or receive command to any valid monitored RT Address.

DMA arbitration
Data Word write (starting at the Data List
Pointer address, incremented for each
successive word)

Status Word Received.

Status Word DMA

DMA Arbitration Status Word write

Exception Handling.

If an interrupting condition occurs during the message, the following occurs:

For High-Priority Interrupts:

HPINT is asserted (if enabled in Register 7). For message errors, the BCRTM is put in a hold state until the interrupt is acknowledged (by writing a "1" to the appropriate bit in Register 8).

For Standard Interrupts:

DMA arbitration (BURST)
Interrupt Status Word write
Monitor Command Block Pointer write
Tail Pointer read (into Register 6)
STDINTP pulses low
STDINTL asserted (if enabled)
Processing continues

Message Completion.

Upon completion of the message, the BCRTM initiates a DMA cycle to update the status word and fetch the address of the next Monitor Command Block:

DMA Arbitration (BURST) Control/Status Word Write Tail Pointer Read

9.0 Exception Handling and Interrupt Logging

The exception handling scheme the BCRTM uses is based on an interrupt structure and provides a high degree of flexibility in:

- defining the events that cause an interrupt,
- selecting between High-Priority and Standard interrupts, and
- selecting the amount of interrupt history retained.

The interrupt structure consists of internal registers that enable interrupt generation, control bits in the RT and BC data structures (see the Remote Terminal Descriptor Definition section, page 23, and the Bus Controller Command Block definition, page 31), and an Interrupt Log List that sequentially stores an interrupt events record in system memory.

The BCRTM generates the Interrupt Log List (see figure 21) to allow the host CPU to view the Standard Interrupt occurrences in chronological order. Each Interrupt Log List entry contains three words. The first, the Interrupt Status Word, indicates the type of interrupt (entries are only for interrupts enabled). In the BC mode, the second word is a Command Block Pointer that refers to the corresponding Command Block. In the RT mode, the second word is a Descriptor Pointer that refers to the corresponding subaddress descriptor. The CPU-initialized third word, a Tail Pointer, is read by the BCRTM to determine the next Interrupt Log List address. The list length can be as long or as short as required. The configuration of the Tail Pointers determines the list length.

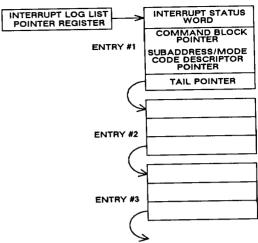


Figure 21. Interrupt Log List

The host CPU initializes the list by setting the tail pointers. This gives flexibility in the list capacity and the ability to link the list around noncontiguous blocks of memory. The host CPU sets the list's starting address using the Interrupt Log List Register. The BCRTM then updates this register with the address of the next list entry.

The internal High-Priority Interrupt Status/Reset Register indicates the cause of a High-Priority Interrupt. The High-Priority Interrupt signal is reset by writing a "1" to the set bits in this register.

The interrupt structure also uses three BCRTM-driven output signals to indicate when an interrupt event occurs:

STDINTL Standard Interrupt Level. This

signal is asserted when one or more of the events enabled in the Standard Interrupt Enable Register occurs. Clear the signal by resetting the Standard Interrupt bit in the High-Priority

Interrupt Status/Reset Register.

STDINTP Standard Interrupt Pulse. This signal is pulsed for each occurrence of an

event enabled in the Standard Interrupt

Enable Register.

HPINT High-Priority Interrupt. This signal is

asserted for each occurrence of an event enabled in the High-Priority Interrupt/Enable Register. Writing to the corresponding bit in the High-Priority

Status/Reset Register resets it.

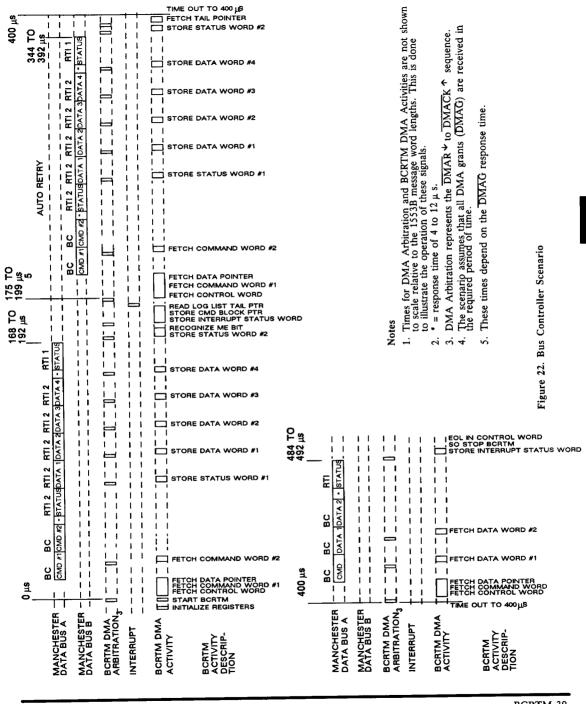
Interrupt Status Word Definition

All bits in the Interrupt Status Word are active high and have the following functions:

Bit	Description

Number

- BIT 15 Interrupt Status Word Accessed. The BCRTM always sets this bit during the DMA Write of the Interrupt Status Word. If the CPU resets this bit after reading the Interrupt Status Word, the bit can help the CPU determine which entries have been acknowledged.
- BIT 14 No Response Time-Out (Message Error condition). Further defines the Message Error condition to indicate that a Response Time-Out condition has occurred.
- BIT 13 (RT) Message Error (ME). Indicates the ME bit was set in the 1553 status word response.
- BITs 12-8 Reserved.
- BIT 7 (RT) Subaddress Event or Mode Code with Data Word Interrupt. Indicates a descriptor control word has been accessed with either an Interrupt Upon Valid Command Received bit set or an Interrupt when Index=0 bit set (and the Index is decremented to 0).
- BIT 6 (RT) Mode Code without Data Word Interrupt. Indicates a mode code has occurred with an Interrupt When Addressed interrupt enabled.
- BIT 5 (RT) Illegal Broadcast Command. Applies to receive commands only. This bit indicates that a received command, due to an illegal mode code or subaddress field, has been received in the broadcast mode. This does not include invalid commands.
- BIT 4 (RT) Illegal Command. This indicates that an illegal command has occurred due to an illegal mode code or subaddress and T/R field. This does not include invalid commands.
- BIT 3 (BC) Polling Comparison Match. Indicates a polling comparison interrupt.
- BIT 2 (BC) Retry Fail. Indicates all the programmed retries have failed.
- BIT 1 (BC,RT) Message Error. Indicates a Message Error has occurred.
- BIT 0 (BC) Interrupt and Continue. This corresponds to the interrupt and continue function described in the Command Block.



BCRTM-39

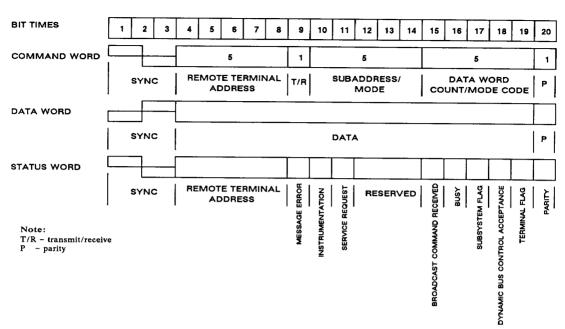


Figure 23. MIL-STD-1553B Word Formats

10.0 OPERATING CONDITIONS*

(Referenced to VSS)

SYMBOL	PARAMETER	LIMITS	UNIT
V_{DD}	DC supply voltage	-0.3 to +7.0	v
V _{I/O}	Voltage on any pin	-0.3 to V _{DD} +0.3	v
II	DC input current	±10	mA
T _{STG}	Storage temperature	-65 to +150	°C
T _{JMAX}	Maximum junction temperature	+175	°c
P_{D}	Average power dissipation (1)	300	mW
$\Theta_{ m JC}$	Thermal resistance, junction-to-case	10	°C/W

Note:

^{1.} Does not reflect the added PD due to an output short-circuited.

Stresses outside the listed absolute maximum ratings may cause permanent damage to the device. This is a stress rating only, and functional operation of the device at these or any other conditions beyond limits indicated in the operational sections of this specification is not recommended. Exposure to absolute maximum rating conditions for extended periods may affect device reliability.

RECOMMENDED OPERATING CONDITIONS

SYMBOL	PARAMETER	LIMITS	UNIT
V _{DD}	DC supply voltage	4.5 to 5.5	V
T _C	Temperature range	-55 to +125	°C
Fo	Operating frequency	12 ±.01%	MHz

11.0 DC ELECTRICAL CHARACTERISTICS

 $(V_{DD} = 5.0 \text{ V} \pm 10\%; -55^{\circ}\text{C} < T_{C} < + 125^{\circ}\text{C})$

SYMBOL	$V \pm 10\%$; -55°C < T _C < + 125°C	CONDITION	MINIMUM	MAXIMUM	UNIT
V _{IL}	Low-level input voltage TTL inputs			0.8	v
V _{IH} Non-RAD	High-level input voltage TTL inputs		2.0		
I _{IN}	Input leakage current TTL inputs Inputs with pull-up resistors Inputs with pull-up resistors	$V_{IN} = V_{DD}$ or V_{SS} $V_{IN} = V_{DD}$ $V_{IN} = V_{SS}$	-1 -1 -550	1 1 -80	μ ^Α μΑ μΑ
V _{IH} RAD-HARD	High-level input voltage (6) TTL inputs (7) CMOS inputs		2.2 .7 V _{DD}		v v
I _{IN}	Input leakage current TTL (7), CMOS inputs Inputs with pull-up resistors Inputs with pull-up resistors	$V_{IN} = V_{DD}$ or V_{SS} $V_{IN} = V_{DD}$ $V_{IN} = V_{SS}$	-10 +150 -900	10 +900 -150	μA μΑ μΑ
V _{OL}	Low-level output voltage TTL outputs CMOS outputs	I _{OL} = 3.2 mA I _{OL} = 50 μA		0.4 V SS +.1	v v
V _{OH}	High-level output voltage TTL outputs CMOS outputs	I _{OH} = -400 μA I _{OH} = 50 μA	2.4 VDD1		v v
I _{OZ}	Three-state output leakage current TTL outputs	$V_{O} = V_{DD}$ or V_{SS}	-10	10	μА
I _{OS}	Short-circuit output current (1,2)	$V_{DD} = 5.5 \text{ V}, V_{O} = V_{DD}$ $V_{DD} = 5.5 \text{ V}, V_{O} = 0 \text{ V}$	-100	100	mA mA
C _{IN}	Input capacitance (3)	F = 1 MHz @ 0 V		10	pF
C _{OUT}	Output capacitance (3)	F = 1 MHz @ 0 V		15	pF
C _{IO}	Bidirect I/O capacitance (3)	F = 1 MHz @ 0 V		20	pF
I _{DD}	Average operating current (1,4)	$F = 12 \text{ MHz}, C_L = 50 \text{ pl}$	F	50	mA
QI _{DD}	Quiescent current	See Note 5		3	mA

- Notes:

 1. Supplied as a design limit, but not guaranteed or tested.

 2. Not more than one output may be shorted at a time for a maximum duration of one second.

 2. Not more than one output may be shorted at a time for a maximum duration of one second.

 3. Measured only for initial qualification, and after process or design changes which may affect input/output capacitance.

 4. Includes current through input pull-up. Instantaneous surge currents on the order of 1 ampere can occur during output switching. Voltage supply should be adequately sized and decoupled to handle a large current surge.

 5. All inputs with internal pull-ups should be left floating. All other inputs should be tied high or low.

 6. Functional tests are conducted in accordance with MIL-STD-883 with the following input test conditions:

 V_{IH} = V_{IH}(min) +20%, -0%; V_{IL} = V_{IL}(max) +0%, -50%, as specified herein, for TTL compatible inputs.

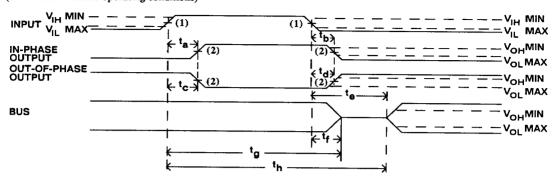
 Devices may be tested using input voltage within the above specified range, but are guaranteed to V_{IH}(min) and V_{IL}(max).

 7. To 1 x 10⁶ total dose; above this level CMOS I/Os required.

BCRTM-41

12.0 AC ELECTRICAL CHARACTERISTICS

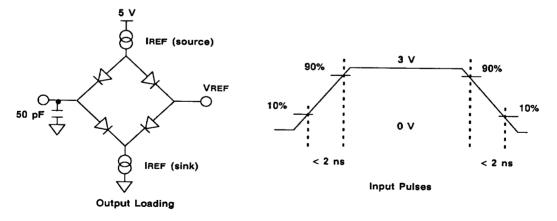
(Over recommended operating conditions)



SYMBOL	PARAMETER
^t a	INPUT ↑ to response ↑
t _b	INPUT ↓ to response ↓
t _c	INPUT ↑ to response ↓
t _d	INPUT ↓ to response ↑
te	INPUT ↓ to data valid
tę	INPUT ↓ to high Z
ta	INPUT ↑ to high Z
th	INPUT ↑ to data valid

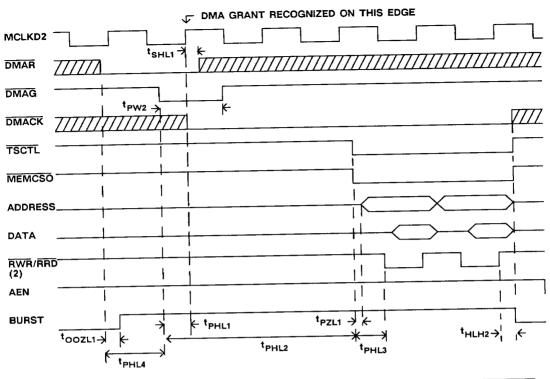
- 1. Timing measurements made at $(V_{IL} MIN + V_{IL} MAX)/2$. 2. Timing measurements made at $(V_{OL} MAX + V_{OH} MIN)/2$.
- 3. Based on 50 pF load.
- 4. Unless otherwise noted, all AC electrical characteristics are guaranteed by design or characterization.

Figure 24. Typical Timing Measurements



50 pF including scope probe and test socket

Figure 25. AC Test Loads and Input Waveforms



SYMBOL	PARAMETER	MIN	MAX	UNITS
t _{SHL1}	DMACK to DMAR High Impedance	0	10	กร
t _{PHL1} (1)	DMAG √to DMACK ✓ See note 3	0	45	ns
t _{PHL2}	DMAG √to TSCTL ↓	2xMCLK	4xMCLK	ns
t _{PZL1}	TSCTL V to ADDRESS valid	0	40	ns
t _{HLH2}	RWR/RRD 10 DMACK 1	THMC1-10	THMC1+10	ns
t _{PHL3}	TSCTL V to RWR/RRD V	MCLK-20	MCLK+20	ns
	DMAG Vto DMAG ↑	MCLK	6xMCLK	ns
t _{PW2}	DMAR Vto BURST 1	0	10	ns
	DMAR √to DMAG ↓ See note 4	0	3.5 (1.9)	μs
t _{PHL4}	DMAR √to DMAG ↓ See note 5	0	1.9 (0.8)	μs

Notes:

1. Guaranteed by test.

2. See figures 27 & 28 for detailed DMA read and write timing.

3. DMAG must be asserted at least 45 ns prior to the rising edge of MCLKD2 in order to be recognized for the next MCLKD2 cycle. If DMAG is not asserted at least 45 ns prior to the rising edge of MCLKD2, DMAG is not recognized until the following MCLKD2 cycle.

4. Provided MCLK = 12 MHz. Number in parentheses indicates the longest DMAR to DMAG alllowed during worst-case bus switching conditions in order to meet MIL-STD-1553B RT Response Time. The number not in parentheses applies to all other circumstances.

Provided MCLK = 6 MHz. Number in parentheses indicates the longest \overline{DMAR} ψ to \overline{DMAG} ψ allowed during worst-case bus switching conditions in order to meet MIL-STD-1553B RT Response Time. The number not in parentheses applies to all other circumstances.

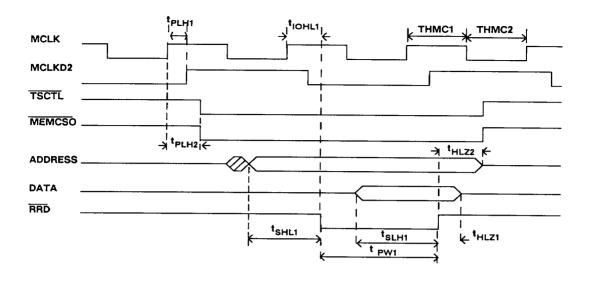
MCLK = period of the memory clock cycle.

BURST signal is for multiple-word DMA accesses.

THMC1 is equivalent to the positive phase of MCLK (see figure 27).

Figure 26. BURST DMA Timing

BCRTM-43

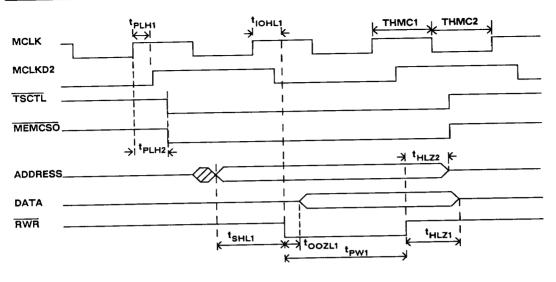


SYMBOL	PARAMETER		MIN	MAX	UNITS
^t sHL1	ADDRESS valid to RRD ↓	(ADDRESS setup)	THMC2-20	THMC2	ns
t _{PW1}	RRD √ to RRD ↑		MCLK-5	MCLK+5	ns
^t HLZ2	RRD 个to ADDRESS High Impedance	(ADDRESS hold)	THMC1-10	THMC1	ns
t _{HLZ1}	RRD 个to DATA High Impedance	(DATA hold)	5	_	ns
^t sLH1	DATA valid to RRD 个	(DATA setup)	40	_	ns
^t PLH1 (1)	MCLK 个to MCLKD2 个		0	40	ns
t _{PLH2}	MCLK ↑ to TSCTL/MEMCSO ↓	***	0	40	ns
t _{IOHL1} (1)	MCLK ↑to RRD ↓	······································	0	60	ns

Note:

Figure 27. BCRTM DMA Read Timing (One-Word Read)

^{1.} Guaranteed by test.

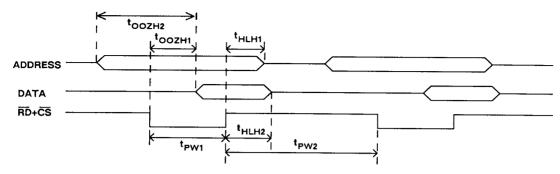


SYMBOL	PARAMETER		MIN	MAX	UNITS
	ADDRESS valid to RWR ↓	(ADDRESS setup)	THMC2-20	THMC2	ns
	RWR vo DATA valid		0	30	ns
t _{HLZ1}	RWR↑to DATA High Impedance	(DATA hold)	THMC1-20	THMC1	ns
tHLZ2	RWR个to ADDRESS High Impedance	(ADDRESS hold)	THMC1-20	THMC1	ns
t _{PW1}	RWR ↓to RWR ↑		MCLK-5	MCLK+5	ns
			0	40	ns
t _{PLH2}	MCLK ↑to TSCTL/MEMCSO ↓		0	40	ns
	MCLK ↑to RWR ↓		0	60	ns

Note:

Figure 28. BCRTM DMA Write Timing (One-Word Write)

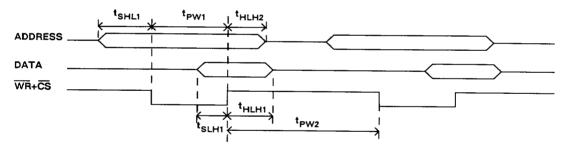
^{1.} Guaranteed by test.



SYMBOL	PARAMETER		MIN	MAX	UNITS
t _{OOZH2}	ADDRESS valid to DATA valid		-	80	ns
t _{HLH2}	RD+CS ↑to DATA High Impedance	(DATA hold)	5	50	ns
^t oozh1 (1)	RD+CS √to DATA valid	(DATA access)	- 1	60	ns
t _{HLH1}	RD+CS 1to ADDRESS High Impedance	(ADDRESS hold)	5	_	ns
t _{PW1}	RD+CS Vto RD+CS ↑		60	-	ns
t _{PW2}	RD+CS ↑to RD+CS ↓		80	-	ns

Notes:

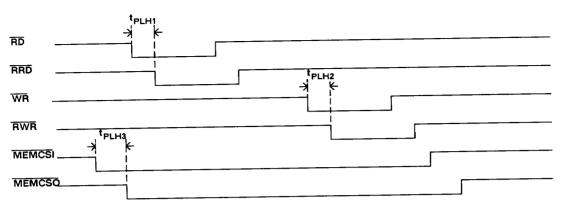
Figure 29. BCRTM Register Read Timing



SYMBOL	PARAMETER		MIN	MAX	UNITS
t _{SHL1}	ADDRESS valid to WR+CS ↓	(ADDRESS setup)	60	_	ns
t _{SLH1}	DATA valid to WR+CS ↑	(DATA setup)	60		ns
t _{PW1}	WR+CSV to WR+CS ↑		60	_	ns
t _{HLH1}	WR+CS↑ to DATA High Impedance	(DATA hold)	10	_	ns
t _{HLH2}	WR+CS↑ to ADDRESS High Impedance		10	_	กร
t _{PW2}	WR+CS↑ to WR+CS V		80	_	ns

Figure 30. BCRTM Register Write Timing

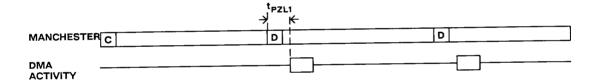
^{1.} Guaranteed by test.
2. User must adhere to both 100ZH1 and 100ZH2 timing constraints to ensure valid data.



SYMBOL	PARAMETER	MIN	MAX	UNITS
t _{PLH1} (1)	RD √to RRD ↓	0	30	ns
t _{PLH2} (1)	WR V to RWR	0	30	ns
t _{PLH3} (1)	MEMCSI ↓to MEMCSO ↓	0	30	ns

Note:

Figure 31. BCRTM Dual-Port Interface Timing Delays



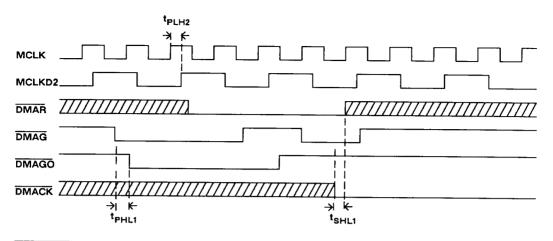
SYMBOL	PARAMETER	MIN	MAX	UNITS
t _{PZL1}	Data word to DMA activity	0	4	μs

This diagram indicates the relationship between the incoming Manchester code and DMA activity (i.e., $\overline{DMAR} \downarrow to \overline{DMACK} \uparrow to \overline{DMACK} \downarrow to \overline{DMACK} \uparrow t$

The pulsewidth = (11 μ s - t_{DMA} - t_{PZL1}) where t_{DMA} is the time to complete DMA activity (i.e., \overline{DMAR} to \overline{DMACK} ?).

Figure 32. DMA Activity (RT Mode)

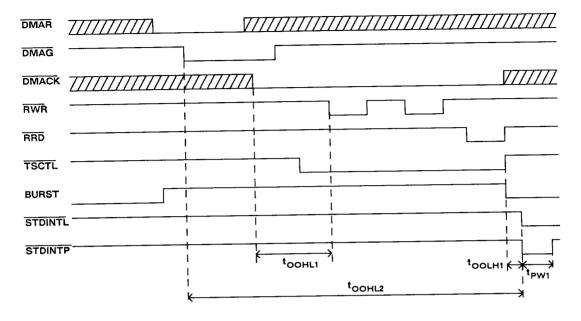
^{1.} Guaranteed by test.



SYMBOL	PARAMETER	MIN	MAX	UNITS
t _{PHL1}	DMAG Vto DMAGO V	0	30	ns
t SHL1	DMACK √to DMAR High Impedance	0	10	ns
t _{PLH2} (1)	MCLK Tto MCLKD2 T	0	40	ns

Figure 33. BCRTM Arbitration when DMAG is Asserted before Arbitration

Notes:
1. Guaranteed by test.
2. When DMAG is asserted before DMAR, the DMAG signal passes through the BCRTM as DMAGO.



SYMBOL	PARAMETER	MIN	MAX	UNITS
	TSCTL↑ to STDINTP/STDINTL ↓	_	1	μS
t _{PW1}	STDINTP Vto STDINTP 1	320	340	ns
toohL1	DMACK ↓to RWR ↓	3xMCLK-10	5xMCLK	ns
t _{OOHL2}	DMAG Vto STDINTL V	8xMCLK+0.5	10xMCLK+1	ns

Note: Address and data bus relationships (not shown) are identical to figure 26.

Figure 34. BCRTM Interrupt Log List Entry Operation Timing

13.0 PACKAGE OUTLINE DRAWINGS

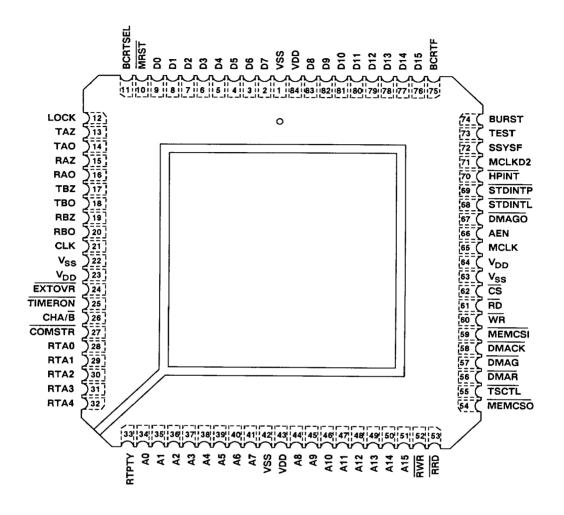


Figure 35a. BCRTM Flatpack and LCC Pin Identification (Top View)
(Flatpack Leads Omitted for Clarity)

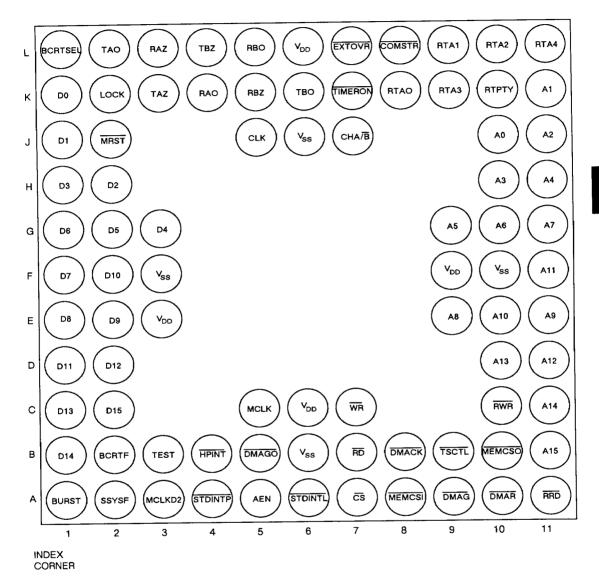


Figure 35b. BCRTM Pingrid Array Pin Identification (Bottom View)