## 32-BIT CMOS MICROPROGRAM MICROPROCESSOR

ADVANCE INFORMATION IDT49C404 IDT49C404A

MICROSLICE ™ PRODUCT

#### **FEATURES:**

- High speed CMOS
  - Microcycle Time: 80ns
- Three bi-directional 32-bit data I/O ports
  - DA, DB, Y
- 64-word x 32-bit expandable 7-port register file
  - 3 input ports and 4 output ports
  - Writes 3 operands and reads 4 operands in one cycle
- · 64-bit in, 32-bit out cascadable funnel shifter
  - Fast alignment to any bit boundary
- 32-bit high-speed ALU cascadable to 64 bits
  - Selects status flags from any bit boundary
- Flexible mask generator and merge logic
  - Selects bit-fields on any width, on any boundary
- Priority encoder
- Powerful orthogonal instruction set
- Built-in multiplication/division support
- Counter function
   Includes Serial Protocol Channel (SPC <sup>™</sup>)
  - Flexible on-chip diagnostics
  - Serially monitors all pin states
  - Reads and writes to Register File
- Single 5V supply
- · Available in 208-pin PGA
- Military product compliant to MIL-STD-883, Class B

## **DESCRIPTION:**

The IDT49C404 is a cascadable, microprogrammable, high-speed CMOS 32-bit microprocessor slice. This monolithic, highly parallel, 3-port device consists of a 7-port 64-word by 32-bit working RAM, 64 bits in/32 bits out cascadable funnel shifter, high-speed multi-function 32-bit ALU and 32-bit mask generation and merge logic.

The IDT49C404 uniquely incorporates shift, ALU and merge functions into a single cycle and utilizes an orthogonal instruction set to create a highly parallel architecture that achieves added performance.

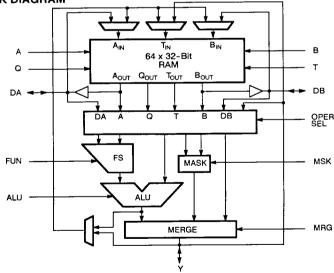
Supporting ultra-fast cycle times, the IDT49C404 offers a verylow-power CMOS alternative to existing bipolar counterparts.

This 32-bit device has been optimized, both architecturally and instruction set-wise, for use in all types of dedicated intelligent controllers such as high-speed graphics engines, array processors, fast disk and communication controllers, robotics, data base manipulation, design automation and Al.

Also featured on the IDT49C404 is an innovative diagnostics capability known as Serial Protocol Channel (SPC). This on-chip feature greatly simplifies the task of writing and debugging microcode, field maintenance debug and test, along with system testing during manufacturing.

The IDT49C404 is fabricated using CEMOS ™, IDT's advanced CMOS technology designed for high-performance and high-reliability. The device is packaged in a 208-lead pin grid array. Military grade product is manufactured in compliance with the latest revision of MIL-STD-883, Class B, making them ideally suited to military temperature applications demanding the highest level of performance and reliability.

## SIMPLIFIED BLOCK DIAGRAM



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## **MILITARY AND COMMERCIAL TEMPERATURE RANGES**

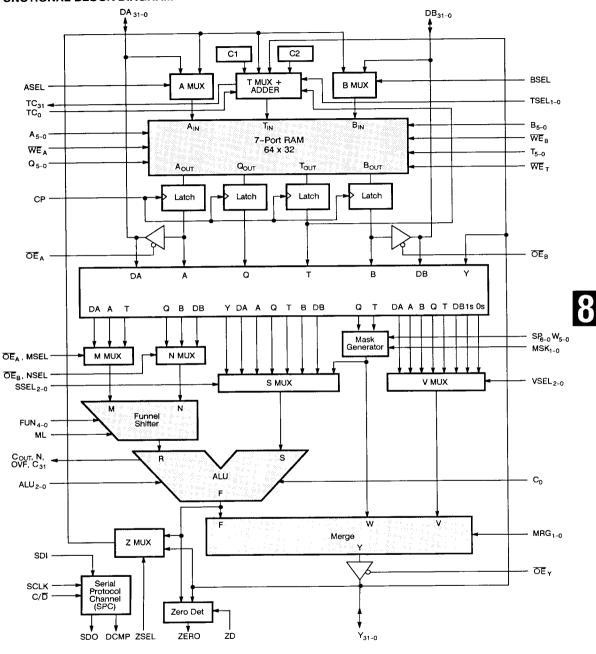
**DECEMBER 1987** 

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DSC-9013/-

## **FUNCTIONAL BLOCK DIAGRAM**



## **IDT49C404 PIN CONFIGURATION**

ı																		l
l	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	s
İ	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	R
١	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	Q
ı	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	Р
	0	0	0	0										0	0	0	0	N
I	0	0	0	0										0	0	0	0	М
1	0	0	0	0			-00	100			F 1986	1		0	0	0	0	L
	0	0	0	0										0	0	0	0	κ
	0	0	0	0			2.5							0	0	0	0	J
i	0	0	0	0			100							0	0	0	0	н
ı	0	0	0	0				*				1		0	0	0	0	G
ı	0	0	0	0										0	0	0	0	F
ı	0	0	0	0										0	0	0	0	Ε
	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	D
1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	С
	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	В
	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	Α
																		]

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17

PGA BOTTOM VIEW

PIN NO.	PIN NAME	PIN NO.	PIN NAME	PIN NO.	PIN NAME
G2		МЗ		Q11	
G3				012	
G4		M14		Q13	
G14		M15		Q14	
G15		M16		Q15	
G16		M17	256	Q16	
G17		N1	n grayî Xala ê rayî	Q17	
Н1		N2		R1	
H2		N3		R2	
нз		N4		R3	
		N14		R4	
H14		N15		R5	
H15		N16		R6	
H16		N17		R7	
H17		P1		R8	
J1		P2		R9	
J2		P3		R10	
J3		P4		R11	
J4		P5		R12	
J14		P6		R13	
J15		P7		R14	
J16		P8		R15	

PIN	PIN	PIN	PIN	PIN PIN	PIN	PIN	PIN	PIN	J17	P9	R16
NO.	NAME	NO.	NAME	NO. NAME	NO.	NAME	NO.	NAME	K1	P10	R17
A1		B1		C1	D1		E1		K2	P11	S1
A2		B2		C2	D2		E2		К3	P12	S2
А3		В3	930	СЗ	D3		E3		K4	P13	S3
A4		В4		C4	. D4		E4		K14	P14	S4
<b>A</b> 5		B5		C5	D5		E14		K15	P15	S5
A6		B6	1:	C6	D6		E15		K16	P16	S6
A7	i	B7		C7	D7		E16		K17	P17	S7
A8		B8		C8	D8		E17		L1	Q1	S8
A9	3.44	В9	4.5 TAS	C9	D9		F1		L2	Q2	S9
A10		B10		C10	D10		F2		L3	Q3	S10
A11		B11		C11	D11	:	F3		L4	Q4	S11
A12	1.84	B12		C12	D12		F4		L14	Q5	S12
A13		B13		C13	D13		F14		L15	Q6	S13
A14		B14	Š	C14	D14		F15		L16	Q7	S14
A15		B15		C15	D15		F16		L17	Q8	S15
A16		B16		C16	D16		F17		M1	<b>Q</b> 9	S16
A17		B17		C17	D17		G1		M2	Q10	S17

# <u>8</u>

## PIN DESCRIPTIONS

DESCRIPTION
Thirty-two-bit data input/output port is under control of the signal $\overline{OE}_A$ . When the $\overline{OE}_A$ is low, RAM output port A can be directly read on these lines. Data on these lines can be selected as the source for the ALU, funnel-shifter or loaded into port A of the working RAM.
Thirty-two-bit data input/output port is under control of the signal $\overline{OE}_B$ . When the $\overline{OE}_B$ is low, RAM output port B can be directly read on these lines. Data on these lines can be selected as the source for the ALU, funnel-shifter or loaded into port B of tworking RAM.
Thirty-two-bit data input/output port is under control of the signal $\overline{OE}_Y$ . When $\overline{OE}_Y$ is low, the merge output can be direct read on these lines. Data on the lines can be loaded into port T of the working RAM or selected as the source for the All when $\overline{OE}_Y$ is high.
A control input pin which, when low, enables the output of merge-logics on the lines $Y_{31-0}$ and, when high, disables the $Y_{31-0}$ three-state output buffers.
The write control signal for RAM input port A. If the signal $\overline{WE}_A$ is low, the data on the DA lines or Z bus is written into the RA (input port A) when the clock signal is low.
The write control signal for RAM input port B. If the signal $\overline{WE}_B$ is low, the data on the DB lines or Z bus is written into the RA (input port B) when the clock signal is low.
The write control signal for RAM input port T. If the signal $\overline{WE}_T$ is low, the data on the Z lines, Y lines, T + C1 or T + C2 is writt into the RAM (input port T) when clock signal is low.
A control input for data input/output port DA. When $\overline{OE}_A$ is low, RAM output port A is read on the DA line. When $\overline{OE}_A$ high, the data on the data lines can be selected as the source for the ALU or loaded into port A of the working RAM.
A control input for data input/output port DB. When $\overline{OE}_B$ is low, RAM output port B can be read on these lines. When $\overline{OE}_B$ high, the data on the DB lines can be selected as the source for the ALU or loaded into port T of the working RAM
The clock input to the IDT49C404. When clock is low, data is written in the seven-port RAM.
Used as carry input for the T counter.
Used as carry output for the T counter.
The input pin which can be used to load the external bit in order to fill in the vacant positions of a word in shift-linkage.
The carry input to the least significant bit of the ALU.
Indicates the carry-output
Indicates the sign N of the ALU operation.
Indicates the conventional two's complement overflow.
The carry output pin which is used to ripple the carry in the expansion mode (64-bit).
The open drain input/output pin which, when high, generally indicates that all outputs are low.
Instruction inputs are used to select the operations for the ALU.
Six RAM address inputs which contains the address of the RAM word appearing at RAM output port A and into which r data is written when WE <sub>a</sub> is low.
Six RAM address inputs which contains the address of the RAM word appearing at RAM output port B and into which r data is written when WE <sub>B</sub> is low.
Six RAM address inputs which contains the address of the RAM word appearing at output port T and into which new d is written under control of TSEL.
Defines what data RAM port A receives, either DA or Z bus.
Defines what data RAM port B receives, either DB or Z bus.
Six RAM address inputs which contain the address of the RAM word appearing at output port Q.
The seven pins are used to specify the start positions or the number of shift positions.
The six pins are used to specify the word width.
Selects the source of the Z bus between the output of the ALU (F) or the Y bus.
Taken together with OE <sub>A</sub> , selects the source of the M input into the funnel shifter.
Taken together with $\overline{\sf OE}_{\sf B}$ , selects the source of the N input into the funnel shifter.
Selects the source of the V bus used for merging with the output of the ALU.
Chooses zero detect of the ALU output (F) or the Y bus.
Selects the source of the S operand input to the ALU.
Controls the operation of the funnel shifter.
Selects the function of the mask generator.
Geracia de la locación de tria masor generator.

## PIN DESCRIPTIONS (Cont'd)

PIN NAME	DESCRIPTION							
TSEL <sub>1-0</sub>	Selects the source of the data to be written into the T port of the RAM.							
SDI	Serial data input to the SPC command and data registers for diagnostics.							
SDO	Serial data output from SPC command and data registers for diagnostics.							
SCLK	SHIFT clock for loading the SPC command and data registers for diagnostics.							
C/D	Command/data control input for SPC operation.							
DCMP	The open drain compare output for SPC diagnostics.							
VCC <sub>7-0</sub>	Eight pins for power supply 5 volt, all of which must be connected to 5 volts.							
GND <sub>16-0</sub>	Sixteen pins for ground, all of which must be connected to ground.							

## **DEVICE ARCHITECTURE**

The IDT49C404 is a high-speed 32-bit microprogrammable CMOS microprocessor slice which can be cascaded to 64-bits. It allows simple, yet high-speed, arithmetic and logic operations on subfields, shift, rotate, mask and merge.

In general, the IDT49C404 can be viewed as a 7-port working RAM feeding into a funnel shifter, then into an ALU and then into merge logic. The control of each of these blocks is orthogonal, allowing the user to select data from registers, shift it, operate on it with the ALU and then merge it in only one cycle. Optionally, the funnel shifter or ALU can be bypassed, allowing the user additional flexibility. In this way, the designer may avoid paying a performance penalty when a particular algorithm requires only one or the other. Thus, the cycle time can be tailored to match the processing requirements.

The IDT49C404 can be divided into the following functional segments:

- Three 32-bit bidirectional I/O ports
- Seven-port 64-word x 32-bit RAM
- 64 bits in/32 bits out cascadable funnel shifter
- 32-bit ALU
- Mask generator
- Merge logic
- Diagnostics circuitry

## THREE BUS ARCHITECTURE

The IDT49C404's 3-bus architecture consists of three bi-directional 32-bit ports (DA, DB and Y). The DA and DB bi-directional buses connect respectively to the A and B RAM outputs and A and B RAM inputs. Thus, data can be read out of the RAM on DA and DB or data can be brought in independently on DA and DB. This special feature allows for easy RAM expansion. Since data can be brought out on the DA and DB buses, other ALU elements can be connected externally which extend the overall ALU, funnel shifter, mask and merge capabilities.

The third 32-bit bus, Y, is the output of the merge logic and also the input back into the RAM ports A, B and T via the Z bus or internal Y bus. The Z MUX multiplexes between the ALU or the Y bus. By selecting the output of the ALU, the results of the ALU operation can be stored back into the RAM while data may be brought out through the merge path onto the Y bus. This results in an ALU operation in parallel with the extraction of data out of the register file. Additionally, there is an alternate data path which allows the Y bus to connect directly into the T MUX such that data can be written from the ALU back into the RAM while data is being brought in, at the same time, through the Y bus to the RAM.

This three bus approach allows for the easy data accessibility necessary when designing high-performance microprocessor-based systems.

#### SEVEN-PORT RAM

The IDT49C404 incorporates a 64-word by 32-bit RAM which has seven ports – four read ports and three write ports. The four read ports are A, B, Q and T. The A and B ports are considered the data path ports and can be used interchangeably. During most cycles, they supply data to the funnel shifter, ALU and merge logic. These ports can be considered to be similar to the A and B ports of the IDT39C203. The Q and T output ports are used mainly for controlling such things as start and width for the funnel shifter and mask generation for merge operations. Since the Q and T ports are outputs of the RAM, the start positions may be computed on previous cycles using the ALU, thus providing extensive programmer flexibility.

There are three write ports; A, B and T. The A and B ports are typically used for results from the current cycle and the T port is used for incrementing counter values in the RAM, as well as loading data from the Y bus in parallel with ALU operations. There are four address buses controlling A, B, Q and T. In one cycle, the seven-port RAM is capable of writing to three locations while reading from four locations. This feature highlights the IDT49C404's highly parallel architecture.

#### **64-BIT FUNNEL SHIFTER**

The funnel shifter accepts two 32-bit operands (A, B, Q, DA, DB or T) which are operated on as a 64-bit word. The output of the funnel shifter is the result of selecting any consecutive 32-bit word within the 64-bit operand. The 32-bit word can start on any bit boundary between 0 and 31. The M and N input muxes allow the user to swap the data as well as duplicate it, allowing for barrel shifting. The funnel shifter also has the capability of taking any 32-bit word as an input and extending the sign, as well as providing zero fill. Through special hooks in the architecture, the funnel shifter can be expanded along with the ALU/merge logic to perform 64-bit operations in a single cycle.

#### A 1 1

The output of the funnel shifter feeds the 32-bit ALU. The ALU can perform conventional binary operations such as logic, addition, subtraction, as well as multiplication and division. Also, the sum of the start and the width information can be used to select the bit boundary from which the carry, sign and overflow flags will output as status. This allows for true arbitrary subfield operations. The other ALU inputs are selected from A, B, Q, T or mask generator.

## MASK GENERATION AND MERGE LOGIC

The mask generation and merge logic allows for field manipulation within the 32-bit resulting word. The mask generator, which determines how the bits will be merged between V and F, is controlled by start and width input pins. The start and width can also come from Q or T. T is used for start and Q is used for width, thus start/width can be calculated, stored in the register file and used in the mask generator. An alternate to the mask generator is a mask

which comes directly from the Q or T outputs of the RAM, allowing for totally arbitrary masks.

The V input of the merge logic comes from a multiplexer which can select any output of the RAM, DA, DB, all 1s or all 0s. The F input is connected to the output of the ALU. The mask is used to merge the V and F input on a bit-by-bit basis, which results in the Y output.

Included in the merge logic is a priority detect circuit. It is used to produce a binary weighted code to indicate the location of the highest order one on its input.

#### SERIAL DIAGNOSTICS

The Serial Protocol Channel (SPC) is a set of pins by which data can be entered into and extracted from a device (such as the IDT49C404) through a serial data input and output port. SPC can be used at many points in the life of a product for diagnostic purposes such as: system level design debug and development, system test during manufacturing and field maintenance debug and test. SPC is of significant benefit as board level packing densities increase. This is because access to test and debug points becomes difficult. This is particularly true in double-sided surface mount technologies.

As companies like IDT continue to integrate more onto each device and put each device into smaller and smaller packages such as surface mount devices, the board level testing becomes more complex for the designer and the manufacturing divisions of companies. To help this situation, a serial diagnostics scheme was developed. It allows for observation of critical signals deep within the system. During system test, when an error is observed, these signals may be modified in order to zero in on the fault in the system.

SPC is primarily a scheme utilizing only four pins SDI, SDO, SCLK, C/D to examine and alter the internal state of a system, for the purpose of monitoring and diagnosing system faults. The SPC has been defined in such a way that it can be implemented with a small number of gates. In many cases, SPC can be added by utilizing less than 5% of the total logic gates. As more gates are added to each device and the number of pins increase, the overhead for diagnostics decreases.

In the following block diagram of a typical application, the Serial Protocol Channel is shown being used with a writable control store in a microprogrammed design. The control store can be initialized through the SPC path. A register with SPC is used for the instruction register going into the IDT49C410 (16-bit microprogram sequencer), as well as data registers around the IDT49C404. In this way, the designer may use the Serial Protocol Channel to observe and modify the microcode coming out of the writable control store, as well as observing and being able to modify data and instructions in the overall machine.

The block diagram of the diagnostics ring in Figure 1 shows how the devices with diagnostics are hooked together in a serial ring via the SDI and SDO signals. The diagnostics signals may be generated through registers which are hooked up to a microprocessor. This microprocessor could conceivably be an IBM PC.

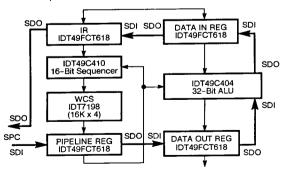


Figure 1. Typical Microprogram Application With SPC

The IDT49C404 accommodates a variety of diagnostics operations. It not only includes the standard Serial Protocol Channel but also the ability to scan data out of the I/O pad cells (as shown in Figure 2) which are connected to the pins of the device. In this way, the state of external connections can be observed, thus telling a lot about the system surrounding the IDT49C404. The scan path through the I/O pad cells is in series with the serial data register.

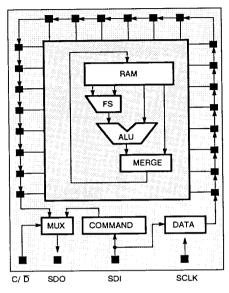


Figure 2. Conceptual Diagram of IDT49C404 Die Incorporating SPC

#### **CONTROL INPUTS**

The control inputs of the IDT49C404 which make up the instruction set are highly orthogonal and provide the user with the highest degree of control over each individual functional unit. Each major unit in the IDT49C404 has its own set of control lines. The following diagrams show the microprogram word layout of the individual fields, as well as the opcodes and functions for each of the fields.

In order to maintain simplicity and orthogonality, the instruction combinations which are used infrequently and require special extended control (divide, multiply, etc.) are grouped together and labled as special instructions. To use these instructions, a special instruction trap door mechanism was employed in the ALU control field (opcode = 101). In the case of special instructions, the T source select and merge control define the particular instructions to be performed. Some special instructions require immediate operands which are provided by other fields such as start, width, funnel shift control, etc.

## **IDT49C404 INSTRUCTION FIELDS**

44	43 42	41 40	39 38	37	36	35	34 30	29	24	23 17	16	14	13	12	11	10	9	7	6 5	4 3	2 0
ŌĒY		MSEL NSEL	SSEL	MAS GEN	- 1	ML	FUNNEL SHIFTER	WIC	тн	START	TSI		AL	J	MF	iG.	vs	SEL	ZSEL ZD	ASEL BSEL	WEA WEB WET
							16-	16-BIT IMMEDIATE FIELD 7-BIT IMMEDIATE FIELD					SPE	CIA	L INS		CTI	ON			

## **INSTRUCTION SET SUMMARY**

M SOURCE SELECTION										
MNEMONIC OE MSEL MSOURCE										
AOE	0	0	Α							
Т	0	1	т							
Α	1	0	Α							
DA	1	1	DA							

MASK SOURCE										
MNEMONIC	М	SK	SOURCE							
EXT	0	0	Start and Width from Instruction							
INT	0	1	T & Q Supply Start and Width							
T32	1	0	T as a 32-Bit Mask							
Q32	1	1	Q as a 32-Bit Mask							

			ALI	J
MNEMONIC		ALU		FUNCTION
ADD	0	0	0	R + S + C <sub>0</sub>
SUBR	0	0	1	S-R-1+C0
SUBS	0	1	0	R - S - 1 + C <sub>0</sub>
OR	0	1	1	R or S
AND	1	0	0	R and S
_	1	0	1	Special Instruction
EXOR	1	1 1		R exor S
EXNOR	1	1	1	R exnor S

	V SOURCE									
MNEMONIC	C VSEL			SOURCE						
DA	0	0	0	DA						
Α	0	0	1	Α						
O.	0	1	0	O						
Т	0	1	1	Т						
В	1	0	0	В						
DB	1	0	1	DB						
ZEROS	1	1	0	0's						
ONEs	1	1	1	1's						

	N SOURCE SELECTION										
MNEMONIC	OE <sub>B</sub>	NSEL	N SOURCE								
BOE	0	0	В								
Q	0	1	Q								
В	1	0	В								
DB	1	1	DB								

-	S SOURCE											
MNEMONIC		SSEL	-	SOURCE								
DA	0	0	0	DA								
Α	0	0	1	Α								
O	0	1	0	Q								
Т	0	1	1	Т								
В	1	0	0	В								
DB	1	0	1	DB								
Y	1	1	0	Y								
MASK	1	1	1	MASK								

		ME	RGE CONTROL
MNEMONIC	М	RG	FUNCTION
F	0	0	Pass F
٧	0	1	Pass V
F to V	1	0	Merge F/V
V to F	1	1	Merge V/F

· · · · · · · · · · · · · · · · · · ·				S	PECIAL	INSTRUCTIONS (ALU = 101)	
MNEMONIC	M F	iG.	VSEL			FUNCTION	OPERANDS
UMLT	0	0	0	0	0	Unsigned Multiply	A, B, T
TMLT	0	0	0	0	1	Two's Complement Multiply	A, B, T
TMLTL	0	0	0	1	0	Two's Complement Multiply Last Cycle	A, B, T
DIVF	0	0	0	1	1	First Divide	A, B, T
DIV	0	0	1	0	0	Second Divide	A, B, T
DIVL	0	0	1	0	1	Last Divide	A, B, T
PRF	0	0	1	1	0	Prioritize First Cycle (32 Bits)	S, Mask
PRS	0	0	1	1	1	Prioritize Second Cycle (64 Bits)	S
INC	0	1	0	0	0	S + Imm (7-Bit) + C <sub>0</sub>	S, Imm
DEC	0	1	0	0	1	S - Imm (7-Bit) - 1 + C <sub>0</sub>	S, Imm
LDI	0	1	0	1	0	Load T with Imm (16-Bit)	16-Bit Imm
LDC1	0	1	0	1	1	Load C1 from Z bus	S
LDC2	0	1	1	0	0	Load C2 from Z bus	S
EXCHG	0	1	1	0	1	Exchange RAM Locations	DA, DB
LDAB	0	1	1	1	0	Load DA into B address	DA
LDBA	0	1	1	1	1	Load DB into A address	DB
SMAGT	1	0	0	0	0	Sign Magnitude/Two's Complement Conversion	s
PROGS	1	0	0	0	1	Program Slice	_

T SOURCE								
MNEMONIC	TS	EL	SOURCE					
z	0	0	Z Bus					
Y	0	1	Y Bus					
TC1	1	0	T + C1 + TC <sub>0</sub>					
TC2	1	1	T + C2 + TC <sub>0</sub>					

A RAM DEST					
MNEMONIC	ASEL	SOURCE			
DA	0	DA Bus			
Z	1	Z Bus			

B RAM DEST						
MNEMONIC	BSEL	SOURCE				
z	0	Z Bus				
DB	1	DB Bus				

## **Z BUS CONTROL**

	Z BUS SO	URCE
MNEMONIC	ZSEL	SOURCE
F	0	F Bus
Y	1	Y Bus

ZERO DETECT SOURCE							
MNEMONIC	ZD	SOURCE					
F	0	F Bus					
Y	1	Y Bus					

		<del></del> -				FUNNEL SHIFT OPERATIONS	
MNEMONIC			FUN			FUNCTION	OPERANDS
SMLZ	0	0	0	0	0	Shift M and fill with 0	O, M
SNLZ	0	0	0	0	1	Shift N and fill with 0	0, N
SMLM	0	0	0	1	0	Shift M and fill with ML	ML, M
SNLM	0	0	0	1	1	Shift N and fill with ML	ML, N
XNM	0	0	1	0	0	Extract field from	N, M
XMN	0	0	1	0	1	Extract field from	M, N
SMAZ	0	0	1	1	0	Shift M arithmetic and fill 0	Sign, M, 0
SNAZ	0	0	1	1	1	Shift N arithmetic and fill 0	Sign, N, 0
SMAM	0	1	0	0	0	Shift M arithmetic and fill ML	Sign, M, ML
SNAM	0	1	0	0	1	Shift N arithmetic and fill ML	Sign, N, ML
ВМ	0	1	0	1	0	Barrel shift M	М
BN	0	1	0	1	1	Barrel shift N	N
РМ	0	1	1	0	0	Pass M	M
PN	0	1	1	0	1	Pass N	N
PZ	0	1	1	1	0	Pass all 0s	0
PO	0	1	1	1	1	Pass all 1s	1
SMLZBA	1	0	0	0	0	Shift M and fill with 0, Bypass ALU	0, M
SNLZBA	1	0	0	0	1	Shift N and fill with 0, Bypass ALU	0, N
SMLMBA	1	0	0	1	0	Shift M and fill with ML, Bypass ALU	ML, M
SNLMBA	1	0	0	1	1	Shift N and fill with ML, Bypass ALU	ML, N
XNMBA	1	0	1	0	0	Extract field from N & M, Bypass ALU	N, M
XMNBA	1	0	1	0	1	Extract field from M & N, Bypass ALU	M, N
SMAZBA	1	0	1	1	0	Shift M arith. and fill 0, Bypass ALU	Sign, M, 0
SNAZBA	1	0	1	1	1	Shift N arith. and fill 0, Bypass ALU	Sign, N, 0
SMAMBA	1	1	0	0	0	Shift M arith. and fill ML, Bypass ALU	Sign, M, ML
SNAMBA	1	1	0	0	1	Shift N arith. and fill ML, Bypass ALU	Sign, N, ML
BMBA	1	1	0	1	0	Barrel shift M, Bypass ALU	М
BNBA	1	1	0	1	1	Barrel shift N, Bypass ALU	N
POCM	1	1	1	0	0	Pass 1s Complement of M	М
POCN	1	1	1	0	1	Pass 1s Complement of N	N
PMFM	1	1	1	1	0	Pass M and fill ML bit from Bit0 to SP	М
PNFM	1	1	1	1	1	Pass N and fill ML from Bit0 to SP	N

# 8

## ABSOLUTE MAXIMUM RATINGS (1)

ABSOLUTE MAXIMUM RATINGS										
SYMBOL	RATING	COMMERCIAL	MILITARY	UNIT						
V <sub>TERM</sub>	Terminal Voltage with Respect to GND	-0.5 to +7.0	-0.5 to +7.0	٧						
TA	Operating Temperature	0 to +70	-55 to +125	°C						
T <sub>BIAS</sub>	Temperature Under Bias	-55 to +125	-65 to +135	°C						
T <sub>STG</sub>	Storage Temperature	-55 to +125	-65 to +150	°C						
PT	Power Dissipation	1.8	1.8	w						
Тоит	DC Output Current	50	50	mA						

#### NOTE:

 Stresses greater than those listed under ABSOLUTE MAXIMUM RAT-INGS may cause permanent damage to the device. This is a stress rating only and functional operation of the device at these or any other conditions above those indicated in the operational sections of this specification is not implied. Exposure to absolute maximum rating conditions for extended periods may affect reliability.

### CAPACITANCE ( $T_A = +25^{\circ}C$ , f = 1.0MHz)

OAI AO	CAI ACITATOE (A									
SYMBOL	PARAMETER(1)	CONDITIONS	TYP.	UNIT						
C <sub>IN</sub>	Input Capacitance	V <sub>IN</sub> = 0V	10	pF						
C <sub>1/O</sub> (2)	I/O Capacitance	V <sub>OUT</sub> = 0V	15	pF						

#### NOTE:

- 1. This parameter is sampled and not 100% tested.
- 2. Includes only output pins.

## DC ELECTRICAL CHARACTERISTICS

 $T_A = 0$ °C to +70°C

 $V_{CC} = 5.0V \pm 5\%$  (Commercial)

 $T_A = -55^{\circ}\text{C to} + 125^{\circ}\text{C}$   $V_{CC} = 5.0\text{V} \pm 10\% \text{ (Military)}$ 

SYMBOL	PARAMETER	TEST C	CONDITIONS (1)	MIN.	TYP. <sup>(2)</sup>	MAX.	UNIT
V <sub>iH</sub>	Input HIGH Level (4)	V <sub>CC</sub> = Max.	2.0		-	V	
V <sub>IL</sub>	Input LOW Level (4)	V <sub>CC</sub> = Min.		-		0.8	٧
- III	Input HIGH Current	V <sub>CC</sub> = Max., V <sub>IN</sub> = V <sub>CC</sub>		_	-	5	μA
	Input LOW Current	V <sub>CC</sub> = Max., V <sub>IN</sub> = 0V				-5	μA
-112	,		i <sub>OH</sub> = -300μA	V <sub>HC</sub>	V <sub>cc</sub>	_	
V <sub>OH</sub>	Output HIGH Voltage	$V_{CC} = Min.$ $V_{IN} = V_{IH} \text{ or } V_{IL}$	I <sub>OH</sub> = -6mA MIL.	2.4	4.3	-	٧
			I <sub>OH</sub> = -8mA COM'L.	2.4	4.3		
			I <sub>OL</sub> = 300µA	_	GND	V <sub>LC</sub>	
14	Output LOW Voltage	V <sub>CC</sub> = Min.	loL = 12mA MIL.	-	0.3	0.5	V
V <sub>OL</sub>	Supur Low Voluge	$V_{IN} = V_{IH} \text{ or } V_{IL}$	I <sub>OL</sub> = 15mA COM'L.	_	0.3	0.5	
	Off Co. 1. (1) in the large decree (		V <sub>O</sub> = 0V		-0.1	-10	
loz	Off State (High Impedance) Output Current	V <sub>CC</sub> = Max.	$V_O = V_{CC}$ (Max.)		0.1	10	μΑ
los	Short Circuit Current	V <sub>CC</sub> = Min., V <sub>OUT</sub> = 0V (S	3)	-15	Ţ -	_	m/

#### NOTES:

- 1. For conditions shown as max. or min. use appropriate value specified under Electrical Characteristics.
- 2. Typical values are at  $V_{CC} = 5.0V$ ,  $+25^{\circ}C$  ambient and maximum loading.
- 3. Not more than one output should be shorted at one time. Duration of the short circuit test should not exceed one second.
- 4. These input levels provide zero noise immunity and should only be static tested in a noise-free environment.

## DC ELECTRICAL CHARACTERISTICS (Cont'd)

 $T_A = 0$ °C to +70°C

 $V_{CC} = 5.0V \pm 5\%$  (Commercial)  $V_{CC} = 5.0V \pm 10\%$  (Military)

 $T_A = -55$ °C to + 125°C  $V_{LC} = 0.2V$ 

 $V_{HC} = V_{CC} - 0.2V$ 

SYMBOL	PARAMETER	TEST CONDITIONS (1)		MIN.	TYP. <sup>(2)</sup>	MAX.	UNIT
Іссан	Quiescent Power Supply Current CP = H	$\begin{array}{l} V_{CC} = \text{Max.} \\ V_{1H} \geq V_{HC}, V_{1L} \leq V_{LC} \\ f_{CP} = 0, CP = V_{CC} \end{array}$		_	_	_	mA
Iccal.	Quiescent Power Supply Current CP = L	$\begin{array}{c} V_{CC} = Max. \\ V_{IH} \geq V_{HC}, V_{IL} \leq V_{LC} \\ I_{CP} = 0, CP = 0V \end{array}$		_	_	_	mA
Гост	Quiescent Input Power Supply (5) Current (per Input @ TTL High)	$V_{CC} = Max. V_{IL} = 3.4V, f_{CP} = 0$		-	-	-	mA
lccD Dynamic Power Supply Cur		V <sub>CC</sub> = Max.	M1L.	-	-		mA/ MHz
	Dynamic Power Supply Current	$V_{HC} \le V_{IH}, V_{IL} \le \underline{V}_{LC}$ Outputs Open, $\overline{OE} = 0V$	COM'L.	_	-	-	
		50% Duty cycle	MIL.	_	-	-	mA
	Total Power Supply Current <sup>(6)</sup>		COM'L.	_	-	-	
lcc		V <sub>CC</sub> = Max., f <sub>CP</sub> = 10MHz Outputs Open, OE = 0V 50% Duty cycle V <sub>IH</sub> = 3.4V, V <sub>IL</sub> = 0.4V	MIL.	-	300	400	
			COM'L.	_	250	350	

#### NOTES:

5. Iccoτ is derived by measuring the total current with all the inputs tied together at 3.4V, subtracting out Iccoch, then dividing by the total number of inputs.

Total Supply Current is the sum of the Quiescent current and the Dynamic current (at either CMOS or TTL input levels). For all conditions, the Total Supply Current can be calculated by using the following equation:

$$I_{CC} = I_{CCQH}(CD_H) + I_{CCQL} (1 - CD_H) + I_{CCT} (N_T \times D_H) + I_{CCD} (f_{CP})$$

CDH = Clock duty cycle high period

 $D_H$  = Data duty cycle TTL high period ( $V_{IN}$  = 3.4V)

N<sub>T</sub> = Number of dynamic inputs driven at TTL levels

f<sub>CP</sub> = Clock Input frequency

## **CMOS TESTING CONSIDERATIONS**

Special test board considerations must be taken into account when applying high-speed CMOS products to the automatic test environment. Large output currents are being switched in very short periods and proper testing demands that test set-ups have minimized inductance and guaranteed zero voltage grounds. The techniques listed below will assist the user in obtaining accurate testing results:

- 1) All input pins should be connected to a voltage potential during testing. If left floating, the device may oscillate, causing improper device operation and possible latchup.
- 2) Placement and value of decoupling capacitors is critical. Each physical set-up has different electrical characteristics and it is recommended that various decoupling capacitor sizes be experimented with. Capacitors should be positioned using the minimum lead lengths. They should also be distributed to decouple power supply lines and be placed as close as possible to the DUT power pins.
- 3) Device grounding is extremely critical for proper device testing. The use of multi-layer performance boards with radial decoupling between power and ground planes is necessary. The ground plane must be sustained from the performance board to the DUT interface board and wiring unused interconnect pins to the ground plane is recommended. Heavy gauge stranded wire should be used for power wiring, with twisted pairs being recommended for minimized inductance.
- 4) To guarantee data sheet compliance, the input thresholds should be tested per input pin, in a static environment. To allow for testing and hardware-induced noise, IDT recommends using  $V_{\rm IL} \leq 0V$  and  $V_{\rm HI} \geq 3V$  for AC tests.

## IDT49C404 PROPAGATION DELAYS (1)

 $T_A = 0$ °C to +70°C,  $V_{CC} = +5V \pm 5\%$ 

					1	O OUTPL	JΤ					
	Y			ZERO		DCMP			STATUS FLAGS			
FROM INPUT	ALU	FS only	ALU & FS	ALU only	FS only	ALU & FS	ALU only	FS only	ALU & FS	ALU only	ALU & FS	TINU
A, B, Q, T	55	55	75	58	58	78	_	_		40	54	ns
DA, DB	37	37	56	40	40	60	_	-	- 11	40	34	ns
NSEL, MSEL, MS, NS, OEA, OEB	_	-	_	_	_	-	_	- 50 20	\$- <u>-</u>		3 -	ns
MSK, FS, STR, W		_			_	_	- 1	<del>-</del>				ns
ALU				_			ंक	-	_	-	_	ns
C <sub>0</sub>		<del></del>	36			-,		-	-	_	_	ns
MRG, ZD			<del>                                     </del>				-1	-	_	-	-	ns
Who, 2D			<del>  _</del> _		-, 20 %	: _	-	* is —	_	_	_	ns

## NOTE:

SET-UP AND HOLD TIMES RELATIVE TO CLOCK (CP INPUT)

CP:						
INPUT	SET-UP TIME BEFORE H-L	HOLD TIME AFTER H-L	SET-UP TIME BEFORE L-H	HOLD TIME AFTER L-H	UNIT	
A, B, T, Q Address (Source or Destination)	18	0	_	1	ns	
DA, DB	-		_	0	ns	
Co, ML	_	-	-	11	ns	
LINES	_	_	_	0	ns	
Y	-		-	0	ns	

On any given cycle, an arithmetic operation without a shift operation can be performed (ALU only) or a shift operation without an arithmetic operation can be performed (FS only) or, finally, both operations in series in one single cycle (ALU + FS).

## IDT49C404A PROPAGATION DELAYS (1)

 $T_A = 0$ °C to +70°C,  $V_{CC} = +5V \pm 5\%$ 

						TO OUTPU	JT					
	Y			ZERO		DCMP			STATUS FLAGS			
FROM INPUT	ALU only	FS only	ALU & FS	ALU only	FS only	ALU & FS	ALU only	FS only	ALU & FS	ALU only	ALU & FS	UNIT
A, B, Q, T	_	_	_	_	_	_	-	_			-	ns
DA, DB	_			_	-	_	_			-	-	ns
NSEL, MSEL, MS, NS, OEA, OEB	_	_	_	-	_		_	-	-	-	_	ns
MSK, FS, STR, W			_		_	-	- 💖	- 1	-	-	_	ns
ALU	_	_	-	_			÷.	-	_	-	_	ns
Co	_	_	_	_	_	6	-	-	1000	_	-	ns
MRG, ZD	_	_	_	_	_		- "	-	_	_	_	ns
Υ	_	_	_	-	- 80		-		_	_		ns

#### NOTE:

SET-UP AND HOLD TIMES RELATIVE TO CLOCK (CP INPUT)

	CP:	<u> </u>			
INPUT	SET-UP TIME BEFORE H-L	HOLD TIME AFTER H-L	SET-UP TIME BEFORE L-H	HOLD TIME AFTER L-H	UNIT
A, B, T, Q Address (Source or Destination)	-	_	_	_	ns
DA, DB	_	<del>-</del>			ns
C <sub>0</sub> , M <sub>L</sub>	_	_			ns
LINES		_	_		ns
Υ		_	_	-	ns

On any given cycle, an arithmetic operation without a shift operation can be performed (ALU only) or a shift operation without an arithmetic operation can be performed (FS only) or, finally, both operations in series in one single cycle (ALU + FS).

## **TEST LOAD CIRCUIT**

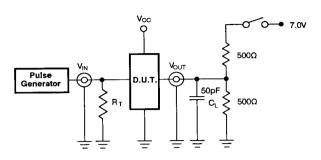


Figure 1. Switching Test Circuits (All Outputs)

#### **SWITCH POSITION**

TEST	SWITCH			
Open Drain Disable Low Enable Low	Closed			
All Other Outputs	Open			

#### DEFINITIONS

 $C_L$  = Load capacitance: includes jig and probe capacitance  $R_T$  = Termination resistance: should be equal to  $Z_{OUT}$  of the Pulse Generator

## **AC TEST CONDITIONS**

Input Pulse Levels	GND to 3.0V
Input Rise/Fall Times	1V/ns
Input Timing Reference Levels	1.5V
Output Reference Levels	1.5V
Output Load	See Figure 1

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## ORDERING INFORMATION

